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**Announcing the**

**SECURE HASH STANDARD**

Federal Information Processing Standards Publications (FIPS PUBS) are issued by the National Institute of Standards and Technology (NIST) after approval by the Secretary of Commerce pursuant to Section 5131 of the Information Technology Management Reform Act of 1996 (Public Law 104-106), and the Computer Security Act of 1987 (Public Law 100-235).

- 1. Name of Standard:** Secure Hash Signature Standard (SHS) (FIPS PUB 180-2).
- 2. Category of Standard:** Computer Security Standard, Cryptography.
- 3. Explanation:** This Standard specifies four secure hash algorithms - SHA-1, SHA-256, SHA-384, and SHA-512 - for computing a condensed representation of electronic data (message). When a message of any length  $< 2^{64}$  bits (for SHA-1 and SHA-256) or  $< 2^{128}$  bits (for SHA-384 and SHA-512) is input to an algorithm, the result is an output called a message digest. The message digests range in length from 160 to 512 bits, depending on the algorithm. Secure hash algorithms are typically used with other cryptographic algorithms, such as digital signature algorithms and keyed-hash message authentication codes, or in the generation of random numbers (bits).

The four hash algorithms specified in this standard are called secure because, for a given algorithm, it is computationally infeasible 1) to find a message that corresponds to a given message digest, or 2) to find two different messages that produce the same message digest. Any change to a message will, with a very high probability, result in a different message digest. This will result in a verification failure when the secure hash algorithm is used with a digital signature algorithm or a keyed-hash message authentication algorithm.

This standard supersedes FIPS 180-1, adding three algorithms that are capable of producing larger message digests. The SHA-1 algorithm specified herein is the same algorithm that was specified previously in FIPS 180-1, although some of the notation has been modified to be consistent with the notation used in the SHA-256, SHA-384, and SHA-512 algorithms.

- 4. Approving Authority:** Secretary of Commerce.
- 5. Maintenance Agency:** U.S. Department of Commerce, National Institute of Standards and Technology (NIST), Information Technology Laboratory (ITL).

**6. Applicability:** This standard is applicable to all Federal departments and agencies for the protection of sensitive unclassified information that is not subject to section 2315 of Title 10, United States Code, or section 3502(2) of Title 44, United States Code. This standard shall be implemented whenever a secure hash algorithm is required for Federal applications, including use by other cryptographic algorithms and protocols. The adoption and use of this standard is available to private and commercial organizations.

**7. Specifications:** Federal Information Processing Standard (FIPS) 180-2, Secure Hash Standard (SHS) (affixed).

**8. Implementations:** The secure hash algorithms specified herein may be implemented in software, firmware, hardware or any combination thereof. Only algorithm implementations that are validated by NIST will be considered as complying with this standard. Information about the planned validation program can be obtained at <http://csrc.nist.gov/cryptval/> or from the National Institute of Standards and Technology, Information Technology Laboratory, Attn: SHS Validation, 100 Bureau Drive Stop 8930, Gaithersburg, MD 20899-8930.

**9. Implementation Schedule:** This standard becomes effective on February 1, 2003.

**10. Patents:** Implementations of the secure hash algorithms in this standard may be covered by U.S. or foreign patents.

**11. Export Control:** Certain cryptographic devices and technical data regarding them are subject to Federal export controls. Exports of cryptographic modules implementing this standard and technical data regarding them must comply with these Federal regulations and be licensed by the Bureau of Export Administration of the U.S. Department of Commerce. Applicable Federal government export controls are specified in Title 15, Code of Federal Regulations (CFR) Part 740.17; Title 15, CFR Part 742; and Title 15, CFR Part 774, Category 5, Part 2.

**12. Qualifications:** While it is the intent of this standard to specify general security requirements for generating a message digest, conformance to this standard does not assure that a particular implementation is secure. The responsible authority in each agency or department shall assure that an overall implementation provides an acceptable level of security. This standard will be reviewed every five years in order to assess its adequacy.

**13. Waiver Procedure.** Under certain exceptional circumstances, the heads of Federal agencies, or their delegates, may approve waivers to Federal Information Processing Standards (FIPS). The heads of such agencies may redelegate such authority only to a senior official designated pursuant to Section 3506(b) of Title 44, U.S. Code. Waivers shall be granted only when compliance with this standard would

- a. adversely affect the accomplishment of the mission of an operator of a Federal computer system or
- b. cause a major adverse financial impact on the operator that is not offset by government-wide savings.

Agency heads may act upon a written waiver request containing the information detailed above. Agency heads may also act without a written waiver request when they determine that conditions for meeting the standard cannot be met. Agency heads may approve waivers only by a written decision that explains the basis on which the agency head made the required finding(s). A copy of each such decision, with procurement sensitive or classified portions clearly identified, shall be sent to: National Institute of Standards and Technology; ATTN: FIPS Waiver Decision, Information Technology Laboratory, 100 Bureau Drive, Stop 8900, Gaithersburg, MD 20899-8900.

In addition, a notice of each waiver granted and each delegation of authority to approve waivers shall be sent promptly to the Committee on Government Operations of the House of Representatives and the Committee on Government Affairs of the Senate and shall be published promptly in the Federal Register.

When the determination on a waiver applies to the procurement of equipment and/or services, a notice of the waiver determination must be published in the Commerce Business Daily as a part of the notice of solicitation for offers of an acquisition or, if the waiver determination is made after that notice is published, by amendment to such notice.

A copy of the waiver, any supporting documents, the document approving the waiver and any supporting and accompanying documents, with such deletions as the agency is authorized and decides to make under Section 552(b) of Title 5, U.S. Code, shall be part of the procurement documentation and retained by the agency.

**14. Where to Obtain Copies of the Standard:** This publication is available electronically by accessing <http://csrc.nist.gov/publications/>. A list of other available computer security publications, including ordering information, can be obtained from NIST Publications List 91, which is available at the same web site. Alternatively, copies of NIST computer security publications are available from: National Technical Information Service (NTIS), 5285 Port Royal Road, Springfield, VA 22161.



**Federal Information**  
**Processing Standards Publication 180-2**  
  
**2002 August 1**  
  
**Specifications for the**  
  
**SECURE HASH STANDARD**

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## 1. INTRODUCTION

This standard specifies four secure hash algorithms, SHA-1<sup>1</sup>, SHA-256, SHA-384, and SHA-512. All four of the algorithms are iterative, one-way hash functions that can process a message to produce a condensed representation called a *message digest*. These algorithms enable the determination of a message's integrity: any change to the message will, with a very high probability, result in a different message digest. This property is useful in the generation and verification of digital signatures and message authentication codes, and in the generation of random numbers (bits).

Each algorithm can be described in two stages: preprocessing and hash computation. Preprocessing involves padding a message, parsing the padded message into  $m$ -bit blocks, and setting initialization values to be used in the hash computation. The hash computation generates a *message schedule* from the padded message and uses that schedule, along with functions, constants, and word operations to iteratively generate a series of hash values. The final hash value generated by the hash computation is used to determine the message digest.

The four algorithms differ most significantly in the number of bits of security that are provided for the data being hashed – this is directly related to the message digest length. When a secure hash algorithm is used in conjunction with another algorithm, there may be requirements specified elsewhere that require the use of a secure hash algorithm with a certain number of bits of security. For example, if a message is being signed with a digital signature algorithm that provides 128 bits of security, then that signature algorithm may require the use of a secure hash algorithm that also provides 128 bits of security (e.g., SHA-256).

Additionally, the four algorithms differ in terms of the size of the blocks and words of data that are used during hashing. Figure 1 presents the basic properties of all four secure hash algorithms.

Algorithm	Message Size (bits)	Block Size (bits)	Word Size (bits)	Message Digest Size (bits)	Security <sup>2</sup> (bits)
SHA-1	$< 2^{64}$	512	32	160	80
SHA-256	$< 2^{64}$	512	32	256	128
SHA-384	$< 2^{128}$	1024	64	384	192
SHA-512	$< 2^{128}$	1024	64	512	256

**Figure 1: Secure Hash Algorithm Properties**

<sup>1</sup> The SHA-1 algorithm specified in this document is identical to the SHA-1 algorithm specified in FIPS 180-1 [180-1]. However, this specification, FIPS 180-2, uses  $ROTL^n(X)$  instead of  $S^n(X)$  [180-1] to denote “circular left shift by  $n$  bits” (i.e., “left rotation by  $n$  bits”). This is described in Sec. 3.2. Some other notational changes have been made in order to be consistent with the specifications for SHA-256, SHA-384, and SHA-512.

<sup>2</sup> In this context, “security” refers to the fact that a birthday attack [HAC] on a message digest of size  $n$  produces a collision with a workfactor of approximately  $2^{n/2}$ .

## 2. DEFINITIONS

### 2.1 Glossary of Terms and Acronyms

Bit	A binary digit having a value of 0 or 1.
Byte	A group of eight bits.
FIPS	Federal Information Processing Standard.
Word	A group of either 32 bits (4 bytes) or 64 bits (8 bytes), depending on the secure hash algorithm.

### 2.2 Algorithm Parameters, Symbols, and Terms

#### 2.2.1 Parameters

The following parameters are used in the secure hash algorithm specifications in this standard.

$a, b, c, \dots, h$	Working variables that are the $w$ -bit words used in the computation of the hash values, $H^{(i)}$ .
$H^{(i)}$	The $i^{\text{th}}$ hash value. $H^{(0)}$ is the <i>initial</i> hash value; $H^{(N)}$ is the <i>final</i> hash value and is used to determine the message digest.
$H_j^{(i)}$	The $j^{\text{th}}$ word of the $i^{\text{th}}$ hash value, where $H_0^{(i)}$ is the left-most word of hash value $i$ .
$K_t$	Constant value to be used for iteration $t$ of the hash computation.
$k$	Number of zeroes appended to a message during the padding step.
$\ell$	Length of the message, $M$ , in bits.
$m$	Number of bits in a message block, $M^{(i)}$ .
$M$	Message to be hashed.
$M^{(i)}$	Message block $i$ , with a size of $m$ bits.
$M_j^{(i)}$	The $j^{\text{th}}$ word of the $i^{\text{th}}$ message block, where $M_0^{(i)}$ is the left-most word of message block $i$ .



$n$	Number of bits to be rotated or shifted when a word is operated upon.
$N$	Number of blocks in the padded message.
$T$	Temporary $w$ -bit word used in the hash computation.
$w$	Number of bits in a word.
$W_t$	The $t^{\text{th}}$ $w$ -bit word of the message schedule.

### 2.2.2 Symbols

The following symbols are used in the secure hash algorithm specifications, and each operates on  $w$ -bit words.

$\wedge$	Bitwise AND operation.
$\vee$	Bitwise OR (“inclusive-OR”) operation.
$\oplus$	Bitwise XOR (“exclusive-OR”) operation.
$\neg$	Bitwise complement operation.
$+$	Addition modulo $2^w$ .
$\ll$	Left-shift operation, where $x \ll n$ is obtained by discarding the left-most $n$ bits of the word $x$ and then padding the result with $n$ zeroes on the right.
$\gg$	Right-shift operation, where $x \gg n$ is obtained by discarding the right-most $n$ bits of the word $x$ and then padding the result with $n$ zeroes on the left.

### 3. NOTATION AND CONVENTIONS

#### 3.1 Bit Strings and Integers

The following terminology related to bit strings and integers will be used.

1. A *hex digit* is an element of the set  $\{0, 1, \dots, 9, a, \dots, f\}$ . A hex digit is the representation of a 4-bit string. For example, the hex digit “7” represents the 4-bit string “0111”, and the hex digit “a” represents the 4-bit string “1010”.
2. A *word* is a  $w$ -bit string that may be represented as a sequence of hex digits. To convert a word to hex digits, each 4-bit string is converted to its hex digit equivalent, as described in (1) above. For example, the 32-bit string

```
1010 0001 0000 0011 1111 1110 0010 0011
```

can be expressed as “a103fe23”, and the 64-bit string

```
1010 0001 0000 0011 1111 1110 0010 0011
0011 0010 1110 1111 0011 0000 0001 1010
```

can be expressed as “a103fe2332ef301a”.

*Throughout this specification, the “big-endian” convention is used when expressing both 32- and 64-bit words, so that within each word, the most significant bit is stored in the left-most bit position.*

3. An *integer* may be represented as a word or pair of words. A word representation of the message length,  $\ell$ , in bits, is required for the padding techniques of Sec. 5.1.

An integer between 0 and  $2^{32}-1$  *inclusive* may be represented as a 32-bit word. The least significant four bits of the integer are represented by the right-most hex digit of the word representation. For example, the integer  $291 = 2^8 + 2^5 + 2^1 + 2^0 = 256+32+2+1$  is represented by the hex word 00000123.

The same holds true for an integer between 0 and  $2^{64}-1$  *inclusive*, which may be represented as a 64-bit word.

If  $Z$  is an integer,  $0 \leq Z < 2^{64}$ , then  $Z = 2^{32}X + Y$ , where  $0 \leq X < 2^{32}$  and  $0 \leq Y < 2^{32}$ . Since  $X$  and  $Y$  can be represented as 32-bit words  $x$  and  $y$ , respectively, the integer  $Z$  can be represented as the pair of words  $(x, y)$ . This property is used for SHA-1 and SHA-256.

If  $Z$  is an integer,  $0 \leq Z < 2^{128}$ , then  $Z = 2^{64}X + Y$ , where  $0 \leq X < 2^{64}$  and  $0 \leq Y < 2^{64}$ . Since  $X$  and  $Y$  can be represented as 64-bit words  $x$  and  $y$ , respectively, the integer  $Z$  can be represented as the pair of words  $(x, y)$ . This property is used for SHA-384 and SHA-512.

4. For the secure hash algorithms, the size of the *message block* -  $m$  bits - depends on the algorithm.
  - a) For **SHA-1** and **SHA-256**, each message block has **512 bits**, which are represented as a sequence of sixteen **32-bit words**.
  - b) For **SHA-384** and **SHA-512**, each message block has **1024 bits**, which are represented as a sequence of sixteen **64-bit words**.

### 3.2 Operations on Words

The following operations are applied to  $w$ -bit words in all four secure hash algorithms. SHA-1 and SHA-256 operate on 32-bit words ( $w = 32$ ), and SHA-384 and SHA-512 operate on 64-bit words ( $w = 64$ ).

1. Bitwise *logical* word operations:  $\wedge$ ,  $\vee$ ,  $\oplus$ , and  $\neg$  (see Sec. 2.2.2).
2. Addition modulo  $2^w$ .

The operation  $x + y$  is defined as follows. The words  $x$  and  $y$  represent integers  $X$  and  $Y$ , where  $0 \leq X < 2^w$  and  $0 \leq Y < 2^w$ . For positive integers  $U$  and  $V$ , let  $U \bmod V$  be the remainder upon dividing  $U$  by  $V$ . Compute

$$Z = (X + Y) \bmod 2^w.$$

Then  $0 \leq Z < 2^w$ . Convert the integer  $Z$  to a word,  $z$ , and define  $z = x + y$ .

3. The *right shift* operation **SHR**  $^n(x)$ , where  $x$  is a  $w$ -bit word and  $n$  is an integer with  $0 \leq n < w$ , is defined by

$$\text{SHR}^n(x) = x \gg n.$$

This operation is used in the SHA-256, SHA-384, and SHA-512 algorithms.

4. The *rotate right* (circular right shift) operation **ROTR**  $^n(x)$ , where  $x$  is a  $w$ -bit word and  $n$  is an integer with  $0 \leq n < w$ , is defined by

$$\text{ROTR}^n(x) = (x \gg n) \vee (x \ll w - n).$$

Thus, **ROTR**  $^n(x)$  is equivalent to a circular shift (rotation) of  $x$  by  $n$  positions to the right.

This operation is used by the SHA-256, SHA-384, and SHA-512 algorithms.

5. The *rotate left* (circular left shift) operation,  $ROTL^n(x)$ , where  $x$  is a  $w$ -bit word and  $n$  is an integer with  $0 \leq n < w$ , is defined by

$$ROTL^n(x) = (x \ll n) \vee (x \gg w - n).$$

Thus,  $ROTL^n(x)$  is equivalent to a circular shift (rotation) of  $x$  by  $n$  positions to the left.

This operation is used only in the SHA-1 algorithm. Note that in Ref. [180-1] this operation was referred to as “ $S^n(X)$ ”; however, the notation has been modified for clarity and consistency with the notation used for operations in the other secure hash algorithms.

6. Note the following equivalence relationships, where  $w$  is fixed in each relationship:

$$ROTL^n(x) \approx ROTR^{w-n}(x)$$

$$ROTR^n(x) \approx ROTL^{w-n}(x).$$

## 4. FUNCTIONS AND CONSTANTS

### 4.1 Functions

This section defines the functions that are used by each of the algorithms. Although the SHA-256, SHA-384, and SHA-512 algorithms all use similar functions, their descriptions are separated into sections for SHA-256 (Sec. 4.1.2) and for SHA-384 and SHA-512 (Sec. 4.1.3), since the input and output for these functions are words of different sizes. Each of the algorithms include  $Ch(x, y, z)$  and  $Maj(x, y, z)$  functions; the exclusive-OR operation ( $\oplus$ ) in these functions may be replaced by a bitwise OR operation ( $\vee$ ) and produce identical results.

#### 4.1.1 SHA-1 Functions

SHA-1 uses a sequence of logical functions,  $f_0, f_1, \dots, f_{79}$ . Each function  $f_t$ , where  $0 \leq t < 79$ , operates on three 32-bit words,  $x$ ,  $y$ , and  $z$ , and produces a 32-bit word as output. The function  $f_t(x, y, z)$  is defined as follows:

$$f_t(x, y, z) = \begin{cases} Ch(x, y, z) = (x \wedge y) \oplus (\neg x \wedge z) & 0 \leq t \leq 19 \\ Parity(x, y, z) = x \oplus y \oplus z & 20 \leq t \leq 39 \\ Maj(x, y, z) = (x \wedge y) \oplus (x \wedge z) \oplus (y \wedge z) & 40 \leq t \leq 59 \\ Parity(x, y, z) = x \oplus y \oplus z & 60 \leq t \leq 79. \end{cases} \quad (4.1)$$

#### 4.1.2 SHA-256 Functions

SHA-256 uses six logical functions, where *each function operates on 32-bit words*, which are represented as  $x$ ,  $y$ , and  $z$ . The result of each function is a new 32-bit word.

$$Ch(x, y, z) = (x \wedge y) \oplus (\neg x \wedge z) \quad (4.2)$$

$$Maj(x, y, z) = (x \wedge y) \oplus (x \wedge z) \oplus (y \wedge z) \quad (4.3)$$

$$\sum_0^{[256]}(x) = ROTR^2(x) \oplus ROTR^{13}(x) \oplus ROTR^{22}(x) \quad (4.4)$$

$$\sum_1^{[256]}(x) = ROTR^6(x) \oplus ROTR^{11}(x) \oplus ROTR^{25}(x) \quad (4.5)$$

$$\mathbf{s}_0^{[256]}(x) = ROTR^7(x) \oplus ROTR^{18}(x) \oplus SHR^3(x) \quad (4.6)$$

$$\mathbf{s}_1^{[256]}(x) = ROTR^{17}(x) \oplus ROTR^{19}(x) \oplus SHR^{10}(x) \quad (4.7)$$

#### 4.1.3 SHA-384 and SHA-512 Functions

SHA-384 and SHA-512 each use six logical functions, where *each function operates on 64-bit words*, which are represented as  $x$ ,  $y$ , and  $z$ . The result of each function is a new 64-bit word.

$$Ch(x, y, z) = (x \wedge y) \oplus (\neg x \wedge z) \quad (4.8)$$

$$Maj(x, y, z) = (x \wedge y) \oplus (x \wedge z) \oplus (y \wedge z) \quad (4.9)$$

$$\sum_0^{[512]}(x) = ROTR^{28}(x) \oplus ROTR^{34}(x) \oplus ROTR^{39}(x) \quad (4.10)$$

$$\sum_1^{[512]}(x) = ROTR^{14}(x) \oplus ROTR^{18}(x) \oplus ROTR^{41}(x) \quad (4.11)$$

$$s_0^{[512]}(x) = ROTR^1(x) \oplus ROTR^8(x) \oplus SHR^7(x) \quad (4.12)$$

$$s_1^{[512]}(x) = ROTR^{19}(x) \oplus ROTR^{61}(x) \oplus SHR^6(x) \quad (4.13)$$

## 4.2 Constants

### 4.2.1 SHA-1 Constants

SHA-1 uses a sequence of eighty constant 32-bit words,  $K_0, K_1, \dots, K_{79}$ , which are given by

$$K_t = \begin{cases} 5a827999 & 0 \leq t \leq 19 \\ 6ed9eba1 & 20 \leq t \leq 39 \\ 8f1bbcdc & 40 \leq t \leq 59 \\ ca62c1d6 & 60 \leq t \leq 79. \end{cases} \quad (4.14)$$

### 4.2.2 SHA-256 Constants

SHA-256 uses a sequence of sixty-four constant 32-bit words,  $K_0^{[256]}, K_1^{[256]}, \dots, K_{63}^{[256]}$ . These words represent the first thirty-two bits of the fractional parts of the cube roots of the first sixty-four prime numbers. In hex, these constant words are (from left to right)

```
428a2f98 71374491 b5c0fbcf e9b5dba5 3956c25b 59f111f1 923f82a4 ab1c5ed5
d807aa98 12835b01 243185be 550c7dc3 72be5d74 80deb1fe 9bdc06a7 c19bf174
e49b69c1 efbe4786 0fc19dc6 240calcc 2de92c6f 4a7484aa 5cb0a9dc 76f988da
983e5152 a831c66d b00327c8 bf597fc7 c6e00bf3 d5a79147 06ca6351 14292967
27b70a85 2e1b2138 4d2c6dfc 53380d13 650a7354 766a0abb 81c2c92e 92722c85
a2bfe8a1 a81a664b c24b8b70 c76c51a3 d192e819 d6990624 f40e3585 106aa070
19a4c116 1e376c08 2748774c 34b0bcb5 391c0cb3 4ed8aa4a 5b9cca4f 682e6ff3
748f82ee 78a5636f 84c87814 8cc70208 90bffffffa a4506ceb bef9a3f7 c67178f2.
```

### 4.2.3 SHA-384 and SHA-512 Constants

SHA-384 and SHA-512 use the same sequence of eighty constant 64-bit words,  $K_0^{[512]}, K_1^{[512]}, \dots, K_{79}^{[512]}$ . These words represent the first sixty-four bits of the fractional parts of the cube roots of the first eighty prime numbers. In hex, these constant words are (from left to right)

```
428a2f98d728ae22 7137449123ef65cd b5c0fbcfec4d3b2f e9b5dba58189dbbc
3956c25bf348b538 59f111f1b605d019 923f82a4af194f9b ab1c5ed5da6d8118
d807aa98a3030242 12835b0145706f6e 243185be4ee4b28c 550c7dc3d5ffb4e2
```

72be5d74f27b896f 80deb1fe3b1696b1 9bdc06a725c71235 c19bf174cf692694  
e49b69c19ef14ad2 efbe4786384f25e3 0fc19dc68b8cd5b5 240calcc77ac9c65  
2de92c6f592b0275 4a7484aa6ea6e483 5cb0a9dcbd41fbd4 76f988da831153b5  
983e5152ee66dfab a831c66d2db43210 b00327c898fb213f bf597fc7beef0ee4  
c6e00bf33da88fc2 d5a79147930aa725 06ca6351e003826f 142929670a0e6e70  
27b70a8546d22ffc 2e1b21385c26c926 4d2c6dfc5ac42aed 53380d139d95b3df  
650a73548baf63de 766a0abb3c77b2a8 81c2c92e47edaee6 92722c851482353b  
a2bfe8a14cf10364 a81a664bbc423001 c24b8b70d0f89791 c76c51a30654be30  
d192e819d6ef5218 d69906245565a910 f40e35855771202a 106aa07032bbd1b8  
19a4c116b8d2d0c8 1e376c085141ab53 2748774cdf8eeb99 34b0bcb5e19b48a8  
391c0cb3c5c95a63 4ed8aa4ae3418acb 5b9cca4f7763e373 682e6ff3d6b2b8a3  
748f82ee5defb2fc 78a5636f43172f60 84c87814a1f0ab72 8cc702081a6439ec  
90beffffa23631e28 a4506cebde82bde9 bef9a3f7b2c67915 c67178f2e372532b  
ca273eceeaa26619c d186b8c721c0c207 eada7dd6cde0eb1e f57d4f7fee6ed178  
06f067aa72176fba 0a637dc5a2c898a6 113f9804bef90dae 1b710b35131c471b  
28db77f523047d84 32caab7b40c72493 3c9ebe0a15c9bebc 431d67c49c100d4c  
4cc5d4becb3e42b6 597f299cfc657e2a 5fcb6fab3ad6faec 6c44198c4a475817.

## 5. PREPROCESSING

Preprocessing shall take place before hash computation begins. This preprocessing consists of three steps: padding the message,  $M$  (Sec. 5.1), parsing the padded message into message blocks (Sec. 5.2), and setting the initial hash value,  $H^{(0)}$  (Sec. 5.3).

### 5.1 Padding the Message

The message,  $M$ , shall be padded before hash computation begins. The purpose of this padding is to ensure that the padded message is a multiple of 512 or 1024 bits, depending on the algorithm.

#### 5.1.1 SHA-1 and SHA-256

Suppose that the length of the message,  $M$ , is  $\ell$  bits. Append the bit “1” to the end of the message, followed by  $k$  zero bits, where  $k$  is the smallest, non-negative solution to the equation  $\ell + 1 + k \equiv 448 \pmod{512}$ . Then append the 64-bit block that is equal to the number  $\ell$  expressed using a binary representation. For example, the (8-bit ASCII) message “abc” has length  $8 \times 3 = 24$ , so the message is padded with a one bit, then  $448 - (24 + 1) = 423$  zero bits, and then the message length, to become the 512-bit padded message

$$\begin{array}{ccccccc} & & & & 423 & & 64 \\ & & & & \overbrace{\hspace{1.5cm}} & & \overbrace{\hspace{1.5cm}} \\ 01100001 & 01100010 & 01100011 & 1 & 00\dots00 & 00\dots0111000 & . \\ \underbrace{\hspace{1.5cm}} & \underbrace{\hspace{1.5cm}} & \underbrace{\hspace{1.5cm}} & & & \underbrace{\hspace{1.5cm}} & \\ \text{“a”} & \text{“b”} & \text{“c”} & & & \ell = 24 & \end{array}$$

The length of the padded message should now be a multiple of 512 bits.

#### 5.1.2 SHA-384 and SHA-512

Suppose the length of the message  $M$ , in bits, is  $\ell$  bits. Append the bit “1” to the end of the message, followed by  $k$  zero bits, where  $k$  is the smallest non-negative solution to the equation  $\ell + 1 + k \equiv 896 \pmod{1024}$ . Then append the 128-bit block that is equal to the number  $\ell$  expressed using a binary representation. For example, the (8-bit ASCII) message “abc” has length  $8 \times 3 = 24$ , so the message is padded with a one bit, then  $896 - (24 + 1) = 871$  zero bits, and then the message length, to become the 1024-bit padded message

$$\begin{array}{ccccccc} & & & & 871 & & 128 \\ & & & & \overbrace{\hspace{1.5cm}} & & \overbrace{\hspace{1.5cm}} \\ 01100001 & 01100010 & 01100011 & 1 & 00\dots00 & 00\dots0111000 & . \\ \underbrace{\hspace{1.5cm}} & \underbrace{\hspace{1.5cm}} & \underbrace{\hspace{1.5cm}} & & & \underbrace{\hspace{1.5cm}} & \\ \text{“a”} & \text{“b”} & \text{“c”} & & & \ell = 24 & \end{array}$$

The length of the padded message should now be a multiple of 1024 bits.



## 5.2 Parsing the Padded Message

After a message has been padded, it must be parsed into  $N$   $m$ -bit blocks before the hash computation can begin.

### 5.2.1 SHA-1 and SHA-256

For SHA-1 and SHA-256, the padded message is parsed into  $N$  512-bit blocks,  $M^{(1)}, M^{(2)}, \dots, M^{(N)}$ . Since the 512 bits of the input block may be expressed as sixteen 32-bit words, the first 32 bits of message block  $i$  are denoted  $M_0^{(i)}$ , the next 32 bits are  $M_1^{(i)}$ , and so on up to  $M_{15}^{(i)}$ .

### 5.2.2 SHA-384 and SHA-512

For SHA-384 and SHA-512, the padded message is parsed into  $N$  1024-bit blocks,  $M^{(1)}, M^{(2)}, \dots, M^{(N)}$ . Since the 1024 bits of the input block may be expressed as sixteen 64-bit words, the first 64 bits of message block  $i$  are denoted  $M_0^{(i)}$ , the next 64 bits are  $M_1^{(i)}$ , and so on up to  $M_{15}^{(i)}$ .

## 5.3 Setting the Initial Hash Value ( $H^{(0)}$ )

Before hash computation begins for each of the secure hash algorithms, the initial hash value,  $H^{(0)}$ , must be set. The size and number of words in  $H^{(0)}$  depends on the message digest size.

### 5.3.1 SHA-1

For SHA-1, the initial hash value,  $H^{(0)}$ , shall consist of the following five 32-bit words, in hex:

$$\begin{aligned} H_0^{(0)} &= 67452301 \\ H_1^{(0)} &= \text{efcdab89} \\ H_2^{(0)} &= 98badcfe \\ H_3^{(0)} &= 10325476 \\ H_4^{(0)} &= \text{c3d2e1f0}. \end{aligned}$$

### 5.3.2 SHA-256

For SHA-256, the initial hash value,  $H^{(0)}$ , shall consist of the following eight 32-bit words, in hex:

$$\begin{aligned} H_0^{(0)} &= 6a09e667 \\ H_1^{(0)} &= \text{bb67ae85} \\ H_2^{(0)} &= 3c6ef372 \\ H_3^{(0)} &= \text{a54ff53a} \\ H_4^{(0)} &= 510e527f \\ H_5^{(0)} &= 9b05688c \\ H_6^{(0)} &= 1f83d9ab \\ H_7^{(0)} &= 5be0cd19. \end{aligned}$$

These words were obtained by taking the first thirty-two bits of the fractional parts of the square roots of the first eight prime numbers.

### 5.3.3 SHA-384

For SHA-384, the initial hash value,  $H^{(0)}$ , shall consist of the following eight 64-bit words, in hex:

$$\begin{aligned} H_0^{(0)} &= \text{cbbb9d5dc1059ed8} \\ H_1^{(0)} &= \text{629a292a367cd507} \\ H_2^{(0)} &= \text{9159015a3070dd17} \\ H_3^{(0)} &= \text{152fec8d8f70e5939} \\ H_4^{(0)} &= \text{67332667ffc00b31} \\ H_5^{(0)} &= \text{8eb44a8768581511} \\ H_6^{(0)} &= \text{db0c2e0d64f98fa7} \\ H_7^{(0)} &= \text{47b5481dbefa4fa4.} \end{aligned}$$

These words were obtained by taking the first sixty-four bits of the fractional parts of the square roots of the ninth through sixteenth prime numbers.

### 5.3.4 SHA-512

For SHA-512, the initial hash value,  $H^{(0)}$ , shall consist of the following eight 64-bit words, in hex:

$$\begin{aligned} H_0^{(0)} &= \text{6a09e667f3bcc908} \\ H_1^{(0)} &= \text{bb67ae8584caa73b} \\ H_2^{(0)} &= \text{3c6ef372fe94f82b} \\ H_3^{(0)} &= \text{a54ff53a5f1d36f1} \\ H_4^{(0)} &= \text{510e527fade682d1} \\ H_5^{(0)} &= \text{9b05688c2b3e6c1f} \\ H_6^{(0)} &= \text{1f83d9abfb41bd6b} \\ H_7^{(0)} &= \text{5be0cd19137e2179.} \end{aligned}$$

These words were obtained by taking the first sixty-four bits of the fractional parts of the square roots of the first eight prime numbers.

## 6. SECURE HASH ALGORITHMS

In the following sections, SHA-512 is described before SHA-384. That is because the SHA-384 algorithm is identical to SHA-512, with the exception of using a different initial hash value and truncating the final hash value to 384 bits.

For each of the secure hash algorithms, there may exist alternate computation methods that yield identical results; one example is the alternative SHA-1 computation described in Sec. 6.1.3. Such alternate methods may be implemented in conformance to this standard.

### 6.1 SHA-1

SHA-1 may be used to hash a message,  $M$ , having a length of  $\ell$  bits, where  $0 \leq \ell < 2^{64}$ . The algorithm uses 1) a message schedule of eighty 32-bit words, 2) five working variables of 32 bits each, and 3) a hash value of five 32-bit words. The final result of SHA-1 is a 160-bit message digest.

The words of the message schedule are labeled  $W_0, W_1, \dots, W_{79}$ . The five working variables are labeled  $a, b, c, d$ , and  $e$ . The words of the hash value are labeled  $H_0^{(i)}, H_1^{(i)}, \dots, H_4^{(i)}$ , which will hold the initial hash value,  $H^{(0)}$ , replaced by each successive intermediate hash value (after each message block is processed),  $H^{(i)}$ , and ending with the final hash value,  $H^{(N)}$ . SHA-1 also uses a single temporary word,  $T$ .

Appendix A gives several detailed examples of SHA-1.

#### 6.1.1 SHA-1 Preprocessing

1. Pad the message,  $M$ , according to Sec. 5.1.1;
2. Parse the padded message into  $N$  512-bit message blocks,  $M^{(1)}, M^{(2)}, \dots, M^{(N)}$ , according to Sec. 5.2.1; and
3. Set the initial hash value,  $H^{(0)}$ , as specified in Sec. 5.3.1.

#### 6.1.2 SHA-1 Hash Computation

The SHA-1 hash computation uses functions and constants previously defined in Sec. 4.1.1 and Sec. 4.2.1, respectively. Addition (+) is performed modulo  $2^{32}$ .

After preprocessing is completed, each message block,  $M^{(1)}, M^{(2)}, \dots, M^{(N)}$ , is processed in order, using the following steps:

For  $i = 1$  to  $N$ :

{

1. Prepare the message schedule,  $\{W_t\}$ :

$$W_t = \begin{cases} M_t^{(i)} & 0 \leq t \leq 15 \\ ROTL^1(W_{t-3} \oplus W_{t-8} \oplus W_{t-14} \oplus W_{t-16}) & 16 \leq t \leq 79 \end{cases}$$

2. Initialize the five working variables,  $a$ ,  $b$ ,  $c$ ,  $d$ , and  $e$ , with the  $(i-1)^{\text{st}}$  hash value:

$$a = H_0^{(i-1)}$$

$$b = H_1^{(i-1)}$$

$$c = H_2^{(i-1)}$$

$$d = H_3^{(i-1)}$$

$$e = H_4^{(i-1)}$$

3. For  $t = 0$  to 79:

$$\left\{ \begin{array}{l} T = ROTL^5(a) + f_t(b, c, d) + e + K_t + W_t \\ e = d \\ d = c \\ c = ROTL^{30}(b) \\ b = a \\ a = T \end{array} \right\}$$

4. Compute the  $i^{\text{th}}$  intermediate hash value  $H^{(i)}$ :

$$H_0^{(i)} = a + H_0^{(i-1)}$$

$$H_1^{(i)} = b + H_1^{(i-1)}$$

$$H_2^{(i)} = c + H_2^{(i-1)}$$

$$H_3^{(i)} = d + H_3^{(i-1)}$$

$$H_4^{(i)} = e + H_4^{(i-1)}$$

}

After repeating steps one through four a total of  $N$  times (i.e., after processing  $M^{(N)}$ ), the resulting 160-bit message digest of the message,  $M$ , is

$$H_0^{(N)} \parallel H_1^{(N)} \parallel H_2^{(N)} \parallel H_3^{(N)} \parallel H_4^{(N)}.$$

### 6.1.3 Alternate Method for Computing a SHA-1 Message Digest

The SHA-1 hash computation method described in Sec. 6.1.2 assumes that the message schedule  $W_0, W_1, \dots, W_{79}$  is implemented as an array of eighty 32-bit words. This is efficient from the standpoint of the minimization of execution time, since the addresses of  $W_{t-3}, \dots, W_{t-16}$  in step (2) of Sec. 6.1.2 are easily computed.

However, if memory is limited, an alternative is to regard  $\{W_t\}$  as a circular queue that may be implemented using an array of sixteen 32-bit words,  $W_0, W_1, \dots, W_{15}$ . The alternate method that is described in this section yields the same message digest as the SHA-1 computation method described in Sec. 6.1.2. Although this alternate method saves sixty-four 32-bit words of storage, it is likely to lengthen the execution time due to the increased complexity of the address computations for the  $\{W_t\}$  in step (3).

For this alternate SHA-1 method, let  $MASK = 0000000f$  (in hex). As in Sec. 6.1.1, addition is performed modulo  $2^{32}$ . Assuming that the preprocessing as described in Sec. 6.1.1 has been performed, the processing of  $M^{(i)}$  is as follows:

For  $i = 1$  to  $N$ :

{

1. For  $t = 0$  to 15:

{

$$W_t = M_t^{(i)}$$

}

2. Initialize the five working variables,  $a, b, c, d$ , and  $e$ , with the  $(i-1)^{\text{st}}$  hash value:

$$a = H_0^{(i-1)}$$

$$b = H_1^{(i-1)}$$

$$c = H_2^{(i-1)}$$

$$d = H_3^{(i-1)}$$

$$e = H_4^{(i-1)}$$

3. For  $t = 0$  to 79:

{

$$s = t \wedge MASK$$

If  $t \geq 16$  then

{

$$W_s = ROTL^1(W_{(s+13) \wedge MASK} \oplus W_{(s+8) \wedge MASK} \oplus W_{(s+2) \wedge MASK} \oplus W_s)$$

}

$$\begin{aligned}
T &= ROTL^5(a) + f_i(b, c, d) + e + K_t + W_s \\
e &= d \\
d &= c \\
c &= ROTL^{30}(b) \\
b &= a \\
a &= T \\
&\}
\end{aligned}$$

4. Compute the  $i^{\text{th}}$  intermediate hash value  $H^{(i)}$ :

$$\begin{aligned}
H_0^{(i)} &= a + H_0^{(i-1)} \\
H_1^{(i)} &= b + H_1^{(i-1)} \\
H_2^{(i)} &= c + H_2^{(i-1)} \\
H_3^{(i)} &= d + H_3^{(i-1)} \\
H_4^{(i)} &= e + H_4^{(i-1)} \\
&\}
\end{aligned}$$

After repeating steps one through four a total of  $N$  times (i.e., after processing  $M^{(N)}$ ), the resulting 160-bit message digest of the message,  $M$ , is

$$H_0^{(N)} \parallel H_1^{(N)} \parallel H_2^{(N)} \parallel H_3^{(N)} \parallel H_4^{(N)}.$$

## 6.2 SHA-256

SHA-256 may be used to hash a message,  $M$ , having a length of  $\ell$  bits, where  $0 \leq \ell < 2^{64}$ . The algorithm uses 1) a message schedule of sixty-four 32-bit words, 2) eight working variables of 32 bits each, and 3) a hash value of eight 32-bit words. The final result of SHA-256 is a 256-bit message digest.

The words of the message schedule are labeled  $W_0, W_1, \dots, W_{63}$ . The eight working variables are labeled  $a, b, c, d, e, f, g$ , and  $h$ . The words of the hash value are labeled  $H_0^{(i)}, H_1^{(i)}, \dots, H_7^{(i)}$ , which will hold the initial hash value,  $H^{(0)}$ , replaced by each successive intermediate hash value (after each message block is processed),  $H^{(i)}$ , and ending with the final hash value,  $H^{(N)}$ . SHA-256 also uses two temporary words,  $T_1$  and  $T_2$ .

Appendix B gives several detailed examples of SHA-256.

### 6.2.1 SHA-256 Preprocessing

1. Pad the message,  $M$ , according to Sec. 5.1.1;
2. Parse the padded message into  $N$  512-bit message blocks,  $M^{(1)}, M^{(2)}, \dots, M^{(N)}$ , according to Sec. 5.2.1; and
3. Set the initial hash value,  $H^{(0)}$ , as specified in Sec. 5.3.2.

### 6.2.2 SHA-256 Hash Computation

The SHA-256 hash computation uses functions and constants previously defined in Sec. 4.1.2 and Sec. 4.2.2, respectively. Addition (+) is performed modulo  $2^{32}$ .

After preprocessing is completed, each message block,  $M^{(1)}, M^{(2)}, \dots, M^{(N)}$ , is processed in order, using the following steps:

For  $i = 1$  to  $N$ :

{

1. Prepare the message schedule,  $\{W_t\}$ :

$$W_t = \begin{cases} M_t^{(i)} & 0 \leq t \leq 15 \\ \mathbf{s}_1^{\{256\}}(W_{t-2}) + W_{t-7} + \mathbf{s}_0^{\{256\}}(W_{t-15}) + W_{t-16} & 16 \leq t \leq 63 \end{cases}$$

2. Initialize the eight working variables,  $a, b, c, d, e, f, g$ , and  $h$ , with the  $(i-1)^{\text{st}}$  hash value:

$$a = H_0^{(i-1)}$$

$$b = H_1^{(i-1)}$$

$$c = H_2^{(i-1)}$$

$$d = H_3^{(i-1)}$$

$$e = H_4^{(i-1)}$$

$$f = H_5^{(i-1)}$$

$$g = H_6^{(i-1)}$$

$$h = H_7^{(i-1)}$$

3. For  $t = 0$  to 63:

{

$$T_1 = h + \sum_1^{\{256\}}(e) + Ch(e, f, g) + K_t^{\{256\}} + W_t$$

$$T_2 = \sum_0^{\{256\}}(a) + Maj(a, b, c)$$

$$h = g$$

$$g = f$$

$$f = e$$

$$e = d + T_1$$

$$d = c$$

$$c = b$$

$$b = a$$

$$a = T_1 + T_2$$

}

4. Compute the  $i^{\text{th}}$  intermediate hash value  $H^{(i)}$ :

$$H_0^{(i)} = a + H_0^{(i-1)}$$

$$H_1^{(i)} = b + H_1^{(i-1)}$$

$$H_2^{(i)} = c + H_2^{(i-1)}$$

$$H_3^{(i)} = d + H_3^{(i-1)}$$

$$H_4^{(i)} = e + H_4^{(i-1)}$$

$$H_5^{(i)} = f + H_5^{(i-1)}$$

$$H_6^{(i)} = g + H_6^{(i-1)}$$

$$H_7^{(i)} = h + H_7^{(i-1)}$$

}

After repeating steps one through four a total of  $N$  times (i.e., after processing  $M^{(N)}$ ), the resulting 256-bit message digest of the message,  $M$ , is

$$H_0^{(N)} \| H_1^{(N)} \| H_2^{(N)} \| H_3^{(N)} \| H_4^{(N)} \| H_5^{(N)} \| H_6^{(N)} \| H_7^{(N)} .$$

### 6.3 SHA-512

SHA-512 may be used to hash a message,  $M$ , having a length of  $\ell$  bits, where  $0 \leq \ell < 2^{128}$ . The algorithm uses 1) a message schedule of eighty 64-bit words, 2) eight working variables of 64 bits each, and 3) a hash value of eight 64-bit words. The final result of SHA-512 is a 512-bit message digest.

The words of the message schedule are labeled  $W_0, W_1, \dots, W_{79}$ . The eight working variables are labeled  $a, b, c, d, e, f, g,$  and  $h$ . The words of the hash value are labeled  $H_0^{(i)}, H_1^{(i)}, \dots, H_7^{(i)}$ , which will hold the initial hash value,  $H^{(0)}$ , replaced by each successive intermediate hash value



(after each message block is processed),  $H^{(i)}$ , and ending with the final hash value,  $H^{(N)}$ . SHA-512 also uses two temporary words,  $T_1$  and  $T_2$ .

Appendix C gives several detailed examples of SHA-512.

### 6.3.1 SHA-512 Preprocessing

1. Pad the message,  $M$ , according to Sec. 5.1.2;
2. Parse the padded message into  $N$  1024-bit message blocks,  $M^{(1)}$ ,  $M^{(2)}$ , ...,  $M^{(N)}$ , according to Sec. 5.2.2; and
3. Set the initial hash value,  $H^{(0)}$ , as specified in Sec. 5.3.4.

### 6.3.2 SHA-512 Hash Computation

The SHA-512 hash computation uses functions and constants previously defined in Sec. 4.1.3 and Sec. 4.2.3, respectively. Addition (+) is performed modulo  $2^{64}$ .

After preprocessing is completed, each message block,  $M^{(1)}$ ,  $M^{(2)}$ , ...,  $M^{(N)}$ , is processed in order, using the following steps:

For  $i = 1$  to  $N$ :

{

1. Prepare the message schedule,  $\{W_t\}$ :

$$W_t = \begin{cases} M_t^{(i)} & 0 \leq t \leq 15 \\ \mathbf{s}_1^{(512)}(W_{t-2}) + W_{t-7} + \mathbf{s}_0^{(512)}(W_{t-15}) + W_{t-16} & 16 \leq t \leq 79 \end{cases}$$

2. Initialize the eight working variables,  $a$ ,  $b$ ,  $c$ ,  $d$ ,  $e$ ,  $f$ ,  $g$ , and  $h$ , with the  $(i-1)^{\text{st}}$  hash value:

$$a = H_0^{(i-1)}$$

$$b = H_1^{(i-1)}$$

$$c = H_2^{(i-1)}$$

$$d = H_3^{(i-1)}$$

$$e = H_4^{(i-1)}$$

$$f = H_5^{(i-1)}$$

$$g = H_6^{(i-1)}$$

$$h = H_7^{(i-1)}$$

3. For  $t = 0$  to 79:

$$\left\{
\begin{aligned}
T_1 &= h + \sum_1^{(512)}(e) + Ch(e, f, g) + K_t^{(512)} + W_t \\
T_2 &= \sum_0^{(512)}(a) + Maj(a, b, c) \\
h &= g \\
g &= f \\
f &= e \\
e &= d + T_1 \\
d &= c \\
c &= b \\
b &= a \\
a &= T_1 + T_2
\end{aligned}
\right.$$

4. Compute the  $i^{\text{th}}$  intermediate hash value  $H^{(i)}$ :

$$\left.
\begin{aligned}
H_0^{(i)} &= a + H_0^{(i-1)} \\
H_1^{(i)} &= b + H_1^{(i-1)} \\
H_2^{(i)} &= c + H_2^{(i-1)} \\
H_3^{(i)} &= d + H_3^{(i-1)} \\
H_4^{(i)} &= e + H_4^{(i-1)} \\
H_5^{(i)} &= f + H_5^{(i-1)} \\
H_6^{(i)} &= g + H_6^{(i-1)} \\
H_7^{(i)} &= h + H_7^{(i-1)}
\end{aligned}
\right\}$$

After repeating steps one through four a total of  $N$  times (i.e., after processing  $M^{(N)}$ ), the resulting 512-bit message digest of the message,  $M$ , is

$$H_0^{(N)} \| H_1^{(N)} \| H_2^{(N)} \| H_3^{(N)} \| H_4^{(N)} \| H_5^{(N)} \| H_6^{(N)} \| H_7^{(N)} .$$

## 6.4 SHA-384

SHA-384 may be used to hash a message,  $M$ , having a length of  $\ell$  bits, where  $0 \leq \ell < 2^{128}$ . The algorithm is defined in the exact same manner as SHA-512 (Sec. 6.3), with the following two exceptions:

1. The initial hash value,  $H^{(0)}$ , shall be set as specified in Sec. 5.3.3; and

2. The 384-bit message digest is obtained by truncating the final hash value,  $H^{(N)}$ , to its left-most 384 bits:

$$H_0^{(N)} \parallel H_1^{(N)} \parallel H_2^{(N)} \parallel H_3^{(N)} \parallel H_4^{(N)} \parallel H_5^{(N)} .$$

Appendix D gives several detailed examples of SHA-384.



## APPENDIX A: SHA-1 EXAMPLES

This appendix is for informational purposes only and is not required to meet the standard.

### A.1 SHA-1 Example (One-Block Message)

Let the message,  $M$ , be the 24-bit ( $\ell = 24$ ) ASCII string "abc", which is equivalent to the following binary string:

01100001 01100010 01100011.

The message is padded by appending a "1" bit, followed by 423 "0" bits, and ending with the hex value 00000000 00000018 (the two 32-bit word representation of the length, 24). Thus, the final padded message consists of one block ( $N = 1$ ).

For SHA-1, the initial hash value,  $H^{(0)}$ , is

$$H_0^{(0)} = 67452301$$

$$H_1^{(0)} = \text{efcdab89}$$

$$H_2^{(0)} = 98badcfe$$

$$H_3^{(0)} = 10325476$$

$$H_4^{(0)} = \text{c3d2e1f0}.$$

The words of the padded message block are then assigned to the words  $W_0, \dots, W_{15}$  of the message schedule:

$W_0 = 61626380$	$W_8 = 00000000$
$W_1 = 00000000$	$W_9 = 00000000$
$W_2 = 00000000$	$W_{10} = 00000000$
$W_3 = 00000000$	$W_{11} = 00000000$
$W_4 = 00000000$	$W_{12} = 00000000$
$W_5 = 00000000$	$W_{13} = 00000000$
$W_6 = 00000000$	$W_{14} = 00000000$
$W_7 = 00000000$	$W_{15} = 00000018.$

The following schedule shows the hex values for  $a$ ,  $b$ ,  $c$ ,  $d$ , and  $e$  after pass  $t$  of the "for  $t = 0$  to 79" loop described in Sec. 6.1.2, step 4.

	$a$	$b$	$c$	$d$	$e$
$t = 0 :$	0116fc33	67452301	7bf36ae2	98badcfe	10325476
$t = 1 :$	8990536d	0116fc33	59d148c0	7bf36ae2	98badcfe
$t = 2 :$	a1390f08	8990536d	c045bf0c	59d148c0	7bf36ae2

t = 3 :	cdd8e11b	a1390f08	626414db	c045bf0c	59d148c0
t = 4 :	cf499de	cdd8e11b	284e43c2	626414db	c045bf0c
t = 5 :	3fc7ca40	cf499de	f3763846	284e43c2	626414db
t = 6 :	993e30c1	3fc7ca40	b3f52677	f3763846	284e43c2
t = 7 :	9e8c07d4	993e30c1	0ff1f290	b3f52677	f3763846
t = 8 :	4b6ae328	9e8c07d4	664f8c30	0ff1f290	b3f52677
t = 9 :	8351f929	4b6ae328	27a301f5	664f8c30	0ff1f290
t = 10 :	fbda9e89	8351f929	12dab8ca	27a301f5	664f8c30
t = 11 :	63188fe4	fbda9e89	60d47e4a	12dab8ca	27a301f5
t = 12 :	4607b664	63188fe4	7ef6a7a2	60d47e4a	12dab8ca
t = 13 :	9128f695	4607b664	18c623f9	7ef6a7a2	60d47e4a
t = 14 :	196bee77	9128f695	1181ed99	18c623f9	7ef6a7a2
t = 15 :	20bdd62f	196bee77	644a3da5	1181ed99	18c623f9
t = 16 :	4e925823	20bdd62f	c65afb9d	644a3da5	1181ed99
t = 17 :	82aa6728	4e925823	c82f758b	c65afb9d	644a3da5
t = 18 :	dc64901d	82aa6728	d3a49608	c82f758b	c65afb9d
t = 19 :	fd9e1d7d	dc64901d	20aa99ca	d3a49608	c82f758b
t = 20 :	1a37b0ca	fd9e1d7d	77192407	20aa99ca	d3a49608
t = 21 :	33a23bfc	1a37b0ca	7f67875f	77192407	20aa99ca
t = 22 :	21283486	33a23bfc	868dec32	7f67875f	77192407
t = 23 :	d541f12d	21283486	0ce88eff	868dec32	7f67875f
t = 24 :	c7567dc6	d541f12d	884a0d21	0ce88eff	868dec32
t = 25 :	48413ba4	c7567dc6	75507c4b	884a0d21	0ce88eff
t = 26 :	be35fbd5	48413ba4	b1d59f71	75507c4b	884a0d21
t = 27 :	4aa84d97	be35fbd5	12104ee9	b1d59f71	75507c4b
t = 28 :	8370b52e	4aa84d97	6f8d7ef5	12104ee9	b1d59f71
t = 29 :	c5fbaf5d	8370b52e	d2aa1365	6f8d7ef5	12104ee9
t = 30 :	1267b407	c5fbaf5d	a0dc2d4b	d2aa1365	6f8d7ef5
t = 31 :	3b845d33	1267b407	717eebd7	a0dc2d4b	d2aa1365
t = 32 :	046faa0a	3b845d33	c499ed01	717eebd7	a0dc2d4b
t = 33 :	2c0ebc11	046faa0a	cee1174c	c499ed01	717eebd7
t = 34 :	21796ad4	2c0ebc11	811bea82	cee1174c	c499ed01
t = 35 :	dcbbb0cb	21796ad4	4b03af04	811bea82	cee1174c
t = 36 :	0f511fd8	dcbbb0cb	085e5ab5	4b03af04	811bea82
t = 37 :	dc63973f	0f511fd8	f72eec32	085e5ab5	4b03af04
t = 38 :	4c986405	dc63973f	03d447f6	f72eec32	085e5ab5
t = 39 :	32de1cba	4c986405	f718e5cf	03d447f6	f72eec32
t = 40 :	fc87dedf	32de1cba	53261901	f718e5cf	03d447f6
t = 41 :	970a0d5c	fc87dedf	8cb7872e	53261901	f718e5cf
t = 42 :	7f193dc5	970a0d5c	ff21f7b7	8cb7872e	53261901
t = 43 :	ee1blaaf	7f193dc5	25c28357	ff21f7b7	8cb7872e
t = 44 :	40f28e09	ee1blaaf	5fc64f71	25c28357	ff21f7b7
t = 45 :	1c51e1f2	40f28e09	fb86c6ab	5fc64f71	25c28357
t = 46 :	a01b846c	1c51e1f2	503ca382	fb86c6ab	5fc64f71
t = 47 :	bead02ca	a01b846c	8714787c	503ca382	fb86c6ab
t = 48 :	baf39337	bead02ca	2806e11b	8714787c	503ca382
t = 49 :	120731c5	baf39337	afab40b2	2806e11b	8714787c
t = 50 :	641db2ce	120731c5	eebce4cd	afab40b2	2806e11b
t = 51 :	3847ad66	641db2ce	4481cc71	eebce4cd	afab40b2
t = 52 :	e490436d	3847ad66	99076cb3	4481cc71	eebce4cd
t = 53 :	27e9f1d8	e490436d	8e11eb59	99076cb3	4481cc71
t = 54 :	7b71f76d	27e9f1d8	792410db	8e11eb59	99076cb3
t = 55 :	5e6456af	7b71f76d	09fa7c76	792410db	8e11eb59
t = 56 :	c846093f	5e6456af	5edc7ddb	09fa7c76	792410db
t = 57 :	d262ff50	c846093f	d79915ab	5edc7ddb	09fa7c76
t = 58 :	09d785fd	d262ff50	f211824f	d79915ab	5edc7ddb

$t = 59 :$	3f52de5a	09d785fd	3498bfd4	f211824f	d79915ab
$t = 60 :$	d756c147	3f52de5a	4275e17f	3498bfd4	f211824f
$t = 61 :$	548c9cb2	d756c147	8fd4b796	4275e17f	3498bfd4
$t = 62 :$	b66c020b	548c9cb2	f5d5b051	8fd4b796	4275e17f
$t = 63 :$	6b61c9e1	b66c020b	9523272c	f5d5b051	8fd4b796
$t = 64 :$	19dfa7ac	6b61c9e1	ed9b0082	9523272c	f5d5b051
$t = 65 :$	101655f9	19dfa7ac	5ad87278	ed9b0082	9523272c
$t = 66 :$	0c3df2b4	101655f9	0677e9eb	5ad87278	ed9b0082
$t = 67 :$	78dd4d2b	0c3df2b4	4405957e	0677e9eb	5ad87278
$t = 68 :$	497093c0	78dd4d2b	030f7cad	4405957e	0677e9eb
$t = 69 :$	3f2588c2	497093c0	de37534a	030f7cad	4405957e
$t = 70 :$	c199f8c7	3f2588c2	125c24f0	de37534a	030f7cad
$t = 71 :$	39859de7	c199f8c7	8fc96230	125c24f0	de37534a
$t = 72 :$	edb42de4	39859de7	f0667e31	8fc96230	125c24f0
$t = 73 :$	11793f6f	edb42de4	ce616779	f0667e31	8fc96230
$t = 74 :$	5ee76897	11793f6f	3b6d0b79	ce616779	f0667e31
$t = 75 :$	63f7dab7	5ee76897	c45e4fdb	3b6d0b79	ce616779
$t = 76 :$	a079b7d9	63f7dab7	d7b9da25	c45e4fdb	3b6d0b79
$t = 77 :$	860d21cc	a079b7d9	d8fdf6ad	d7b9da25	c45e4fdb
$t = 78 :$	5738d5e1	860d21cc	681e6df6	d8fdf6ad	d7b9da25
$t = 79 :$	42541b35	5738d5e1	21834873	681e6df6	d8fdf6ad

That completes the processing of the first and only message block,  $M^{(1)}$ . The final hash value,  $H^{(1)}$ , is calculated to be

$$\begin{aligned}
 H_0^{(1)} &= 67452301 + 42541b35 = a9993e36 \\
 H_1^{(1)} &= efcdab89 + 5738d5e1 = 4706816a \\
 H_2^{(1)} &= 98badcfe + 21834873 = ba3e2571 \\
 H_3^{(1)} &= 10325476 + 681e6df6 = 7850c26c \\
 H_4^{(1)} &= c3d2e1f0 + d8fdf6ad = 9cd0d89d.
 \end{aligned}$$

The resulting 160-bit message digest is

$$a9993e36 \ 4706816a \ ba3e2571 \ 7850c26c \ 9cd0d89d.$$

## A.2 SHA-1 Example (Multi-Block Message)

Let the message,  $M$ , be the 448-bit ( $\ell = 448$ ) ASCII string

"abcdbcdecdefdefgefghfghighijhijkijklklmklmnlmnomnopnopq".

The message is padded by appending a "1" bit, followed by 511 "0" bits, and ending with the hex value 00000000 000001c0 (the two 32-bit word representation of the length, 448). Thus, the final padded message consists of two blocks ( $N = 2$ ).

For SHA-1, the initial hash value,  $H^{(0)}$ , is

$$H_0^{(0)} = 67452301$$

$$H_1^{(0)} = \text{efcdab89}$$

$$H_2^{(0)} = 98badcfe$$

$$H_3^{(0)} = 10325476$$

$$H_4^{(0)} = \text{c3d2e1f0}.$$

The words of the first padded message block,  $M^{(1)}$ , are then assigned to the words  $W_0, \dots, W_{15}$  of the message schedule:

$W_0 = 61626364$	$W_8 = 696a6b6c$
$W_1 = 62636465$	$W_9 = 6a6b6c6d$
$W_2 = 63646566$	$W_{10} = 6b6c6d6e$
$W_3 = 64656667$	$W_{11} = 6c6d6e6f$
$W_4 = 65666768$	$W_{12} = 6d6e6f70$
$W_5 = 66676869$	$W_{13} = 6e6f7071$
$W_6 = 6768696a$	$W_{14} = 80000000$
$W_7 = 68696a6b$	$W_{15} = 00000000.$

The following schedule shows the hex values for  $a$ ,  $b$ ,  $c$ ,  $d$ , and  $e$  after pass  $t$  of the “for  $t = 0$  to 79” loop described in Sec. 6.1.2, step 4.

	$a$	$b$	$c$	$d$	$e$
$t = 0 :$	0116fc17	67452301	7bf36ae2	98badcfe	10325476
$t = 1 :$	ebf3b452	0116fc17	59d148c0	7bf36ae2	98badcfe
$t = 2 :$	5109913a	ebf3b452	c045bf05	59d148c0	7bf36ae2
$t = 3 :$	2c4f6eac	5109913a	bafced14	c045bf05	59d148c0
$t = 4 :$	33f4ae5b	2c4f6eac	9442644e	bafced14	c045bf05
$t = 5 :$	96b85189	33f4ae5b	0b13dbab	9442644e	bafced14
$t = 6 :$	db04cb58	96b85189	ccfd2b96	0b13dbab	9442644e
$t = 7 :$	45833f0f	db04cb58	65ae1462	ccfd2b96	0b13dbab
$t = 8 :$	c565c35e	45833f0f	36c132d6	65ae1462	ccfd2b96
$t = 9 :$	6350afda	c565c35e	d160cfc3	36c132d6	65ae1462
$t = 10 :$	8993ea77	6350afda	b15970d7	d160cfc3	36c132d6
$t = 11 :$	e19ecaa2	8993ea77	98d42bf6	b15970d7	d160cfc3
$t = 12 :$	8603481e	e19ecaa2	e264fa9d	98d42bf6	b15970d7
$t = 13 :$	32f94a85	8603481e	b867b2a8	e264fa9d	98d42bf6
$t = 14 :$	b2e7a8be	32f94a85	a180d207	b867b2a8	e264fa9d
$t = 15 :$	42637e39	b2e7a8be	4cbe52a1	a180d207	b867b2a8
$t = 16 :$	6b068048	42637e39	acb9ea2f	4cbe52a1	a180d207
$t = 17 :$	426b9c35	6b068048	5098df8e	acb9ea2f	4cbe52a1
$t = 18 :$	944b1bd1	426b9c35	1ac1a012	5098df8e	acb9ea2f
$t = 19 :$	6c445652	944b1bd1	509ae70d	1ac1a012	5098df8e
$t = 20 :$	95836da5	6c445652	6512c6f4	509ae70d	1ac1a012
$t = 21 :$	09511177	95836da5	9b111594	6512c6f4	509ae70d
$t = 22 :$	e2b92dc4	09511177	6560db69	9b111594	6512c6f4
$t = 23 :$	fd224575	e2b92dc4	c254445d	6560db69	9b111594
$t = 24 :$	eeb82d9a	fd224575	38ae4b71	c254445d	6560db69
$t = 25 :$	5a142c1a	eeb82d9a	7f48915d	38ae4b71	c254445d



t = 26 :	2972f7c7	5a142c1a	bbae0b66	7f48915d	38ae4b71
t = 27 :	d526a644	2972f7c7	96850b06	bbae0b66	7f48915d
t = 28 :	e1122421	d526a644	ca5cbdf1	96850b06	bbae0b66
t = 29 :	05b457b2	e1122421	3549a991	ca5cbdf1	96850b06
t = 30 :	a9c84bec	05b457b2	78448908	3549a991	ca5cbdf1
t = 31 :	52e31f60	a9c84bec	816d15ec	78448908	3549a991
t = 32 :	5af3242c	52e31f60	2a7212fb	816d15ec	78448908
t = 33 :	31c756a9	5af3242c	14b8c7d8	2a7212fb	816d15ec
t = 34 :	e9ac987c	31c756a9	16bcc90b	14b8c7d8	2a7212fb
t = 35 :	ab7c32ee	e9ac987c	4c71d5aa	16bcc90b	14b8c7d8
t = 36 :	5933fc99	ab7c32ee	3a6b261f	4c71d5aa	16bcc90b
t = 37 :	43f87ae9	5933fc99	aadf0cbb	3a6b261f	4c71d5aa
t = 38 :	24957f22	43f87ae9	564cff26	aadf0cbb	3a6b261f
t = 39 :	adeb7478	24957f22	50fe1eba	564cff26	aadf0cbb
t = 40 :	d70e5010	adeb7478	89255fc8	50fe1eba	564cff26
t = 41 :	79bcfb08	d70e5010	2b7add1e	89255fc8	50fe1eba
t = 42 :	f9bcb8de	79bcfb08	35c39404	2b7add1e	89255fc8
t = 43 :	633e9561	f9bcb8de	1e6f3ec2	35c39404	2b7add1e
t = 44 :	98clea64	633e9561	be6f2e37	1e6f3ec2	35c39404
t = 45 :	c6ea241e	98clea64	58cfa558	be6f2e37	1e6f3ec2
t = 46 :	a2ad4f02	c6ea241e	26307a99	58cfa558	be6f2e37
t = 47 :	c8a69090	a2ad4f02	b1ba8907	26307a99	58cfa558
t = 48 :	88341600	c8a69090	a8ab53c0	b1ba8907	26307a99
t = 49 :	7e846f58	88341600	3229a424	a8ab53c0	b1ba8907
t = 50 :	86e358ba	7e846f58	220d0580	3229a424	a8ab53c0
t = 51 :	8d2e76c8	86e358ba	1fa11bd6	220d0580	3229a424
t = 52 :	ce892e10	8d2e76c8	alb8d62e	1fa11bd6	220d0580
t = 53 :	edea95b1	ce892e10	234b9db2	alb8d62e	1fa11bd6
t = 54 :	36d1230a	edea95b1	33a24b84	234b9db2	alb8d62e
t = 55 :	776c3910	36d1230a	7b7aa56c	33a24b84	234b9db2
t = 56 :	a681b723	776c3910	8db448c2	7b7aa56c	33a24b84
t = 57 :	ac0a794f	a681b723	1ddb0e44	8db448c2	7b7aa56c
t = 58 :	f03d3782	ac0a794f	e9a06dc8	1ddb0e44	8db448c2
t = 59 :	9ef775c3	f03d3782	eb029e53	e9a06dc8	1ddb0e44
t = 60 :	36254b13	9ef775c3	bc0f4de0	eb029e53	e9a06dc8
t = 61 :	4080d4dc	36254b13	e7bddd70	bc0f4de0	eb029e53
t = 62 :	2bfaf7a8	4080d4dc	cd8952c4	e7bddd70	bc0f4de0
t = 63 :	513f9ca0	2bfaf7a8	10203537	cd8952c4	e7bddd70
t = 64 :	e5895c81	513f9ca0	0afebdea	10203537	cd8952c4
t = 65 :	1037d2d5	e5895c81	144fe728	0afebdea	10203537
t = 66 :	14a82da9	1037d2d5	79625720	144fe728	0afebdea
t = 67 :	6d17c9fd	14a82da9	440df4b5	79625720	144fe728
t = 68 :	2c7b07bd	6d17c9fd	452a0b6a	440df4b5	79625720
t = 69 :	fdf6efff	2c7b07bd	5b45f27f	452a0b6a	440df4b5
t = 70 :	112b96e3	fdf6efff	4b1ec1ef	5b45f27f	452a0b6a
t = 71 :	84065712	112b96e3	ff7dbbfff	4b1ec1ef	5b45f27f
t = 72 :	ab89fb71	84065712	c44ae5b8	ff7dbbfff	4b1ec1ef
t = 73 :	c5210e35	ab89fb71	a10195c4	c44ae5b8	ff7dbbfff
t = 74 :	352d9f4b	c5210e35	6ae27edc	a10195c4	c44ae5b8
t = 75 :	1a0e0e0a	352d9f4b	7148438d	6ae27edc	a10195c4
t = 76 :	d0d47349	1a0e0e0a	cd4b67d2	7148438d	6ae27edc
t = 77 :	ad38620d	d0d47349	86838382	cd4b67d2	7148438d
t = 78 :	d3ad7c25	ad38620d	74351cd2	86838382	cd4b67d2
t = 79 :	8ce34517	d3ad7c25	6b4e1883	74351cd2	86838382

That completes the processing of the first message block,  $M^{(1)}$ . The first intermediate hash value,  $H^{(1)}$ , is calculated to be

$$\begin{aligned} H_0^{(1)} &= 67452301 + 8ce34517 = f4286818 \\ H_1^{(1)} &= efcdab89 + d3ad7c25 = c37b27ae \\ H_2^{(1)} &= 98badcfe + 6b4e1883 = 0408f581 \\ H_3^{(1)} &= 10325476 + 74351cd2 = 84677148 \\ H_4^{(1)} &= c3d2e1f0 + 86838382 = 4a566572. \end{aligned}$$

The words of the *second* padded message block,  $M^{(2)}$ , are then assigned to the words  $W_0, \dots, W_{15}$  of the message schedule:

$$\begin{array}{ll} W_0 = 00000000 & W_8 = 00000000 \\ W_1 = 00000000 & W_9 = 00000000 \\ W_2 = 00000000 & W_{10} = 00000000 \\ W_3 = 00000000 & W_{11} = 00000000 \\ W_4 = 00000000 & W_{12} = 00000000 \\ W_5 = 00000000 & W_{13} = 00000000 \\ W_6 = 00000000 & W_{14} = 00000000 \\ W_7 = 00000000 & W_{15} = 000001c0. \end{array}$$

The following schedule shows the hex values for  $a$ ,  $b$ ,  $c$ ,  $d$ , and  $e$  after pass  $t$  of the “for  $t = 0$  to 79” loop described in Sec. 6.1.2, step 4.

	$a$	$b$	$c$	$d$	$e$
$t = 0 :$	2df257e9	f4286818	b0dec9eb	0408f581	84677148
$t = 1 :$	4d3dc58f	2df257e9	3d0a1a06	b0dec9eb	0408f581
$t = 2 :$	c352bb05	4d3dc58f	4b7c95fa	3d0a1a06	b0dec9eb
$t = 3 :$	eef743c6	c352bb05	d34f7163	4b7c95fa	3d0a1a06
$t = 4 :$	41e34277	eef743c6	70d4aec1	d34f7163	4b7c95fa
$t = 5 :$	5443915c	41e34277	bbbdd0f1	70d4aec1	d34f7163
$t = 6 :$	e7fa0377	5443915c	d078d09d	bbbdd0f1	70d4aec1
$t = 7 :$	c6946813	e7fa0377	1510e457	d078d09d	bbbdd0f1
$t = 8 :$	fdde1de1	c6946813	f9fe80dd	1510e457	d078d09d
$t = 9 :$	b8538aca	fdde1de1	f1a51a04	f9fe80dd	1510e457
$t = 10 :$	6ba94f63	b8538aca	7f778778	f1a51a04	f9fe80dd
$t = 11 :$	43a2792f	6ba94f63	ae14e2b2	7f778778	f1a51a04
$t = 12 :$	fec7bbf	43a2792f	daea53d8	ae14e2b2	7f778778
$t = 13 :$	a2604ca8	fec7bbf	d0e89e4b	daea53d8	ae14e2b2
$t = 14 :$	258b0baa	a2604ca8	ffb35eef	d0e89e4b	daea53d8
$t = 15 :$	d9772360	258b0baa	2898132a	ffb35eef	d0e89e4b
$t = 16 :$	5507db6e	d9772360	8962c2ea	2898132a	ffb35eef
$t = 17 :$	a51b58bc	5507db6e	365dc8d8	8962c2ea	2898132a
$t = 18 :$	c2eb709f	a51b58bc	9541f6db	365dc8d8	8962c2ea
$t = 19 :$	d8992153	c2eb709f	2946d62f	9541f6db	365dc8d8
$t = 20 :$	37482f5f	d8992153	f0badc27	2946d62f	9541f6db
$t = 21 :$	ee8700bd	37482f5f	f6264854	f0badc27	2946d62f

$t = 22 :$	9ad594b9	ee8700bd	cdd20bd7	f6264854	f0badc27
$t = 23 :$	8fb5aa5b9	9ad594b9	7ba1c02f	cdd20bd7	f6264854
$t = 24 :$	88fb5867	8fb5aa5b9	66b5652e	7ba1c02f	cdd20bd7
$t = 25 :$	eec50521	88fb5867	63eea96e	66b5652e	7ba1c02f
$t = 26 :$	50bce434	eec50521	e23ed619	63eea96e	66b5652e
$t = 27 :$	5c416daf	50bce434	7bb14148	e23ed619	63eea96e
$t = 28 :$	2429be5f	5c416daf	142f390d	7bb14148	e23ed619
$t = 29 :$	0a2fb108	2429be5f	d7105b6b	142f390d	7bb14148
$t = 30 :$	17986223	0a2fb108	c90a6f97	d7105b6b	142f390d
$t = 31 :$	8a4af384	17986223	028bec42	c90a6f97	d7105b6b
$t = 32 :$	6b629993	8a4af384	c5e61888	028bec42	c90a6f97
$t = 33 :$	f15f04f3	6b629993	2292bce1	c5e61888	028bec42
$t = 34 :$	295cc25b	f15f04f3	dad8a664	2292bce1	c5e61888
$t = 35 :$	696da404	295cc25b	fc57c13c	dad8a664	2292bce1
$t = 36 :$	cef5ae12	696da404	ca573096	fc57c13c	dad8a664
$t = 37 :$	87d5b80c	cef5ae12	1a5b6901	ca573096	fc57c13c
$t = 38 :$	84e2a5f2	87d5b80c	b3bd6b84	1a5b6901	ca573096
$t = 39 :$	03bb6310	84e2a5f2	21f56e03	b3bd6b84	1a5b6901
$t = 40 :$	c2d8f75f	03bb6310	a138a97c	21f56e03	b3bd6b84
$t = 41 :$	bfb25768	c2d8f75f	00eed8c4	a138a97c	21f56e03
$t = 42 :$	28589152	bfb25768	f0b63dd7	00eed8c4	a138a97c
$t = 43 :$	ec1d3d61	28589152	2fec95da	f0b63dd7	00eed8c4
$t = 44 :$	3caed7af	ec1d3d61	8a162454	2fec95da	f0b63dd7
$t = 45 :$	c3d033ea	3caed7af	7b074f58	8a162454	2fec95da
$t = 46 :$	7316056a	c3d033ea	cf2bb5eb	7b074f58	8a162454
$t = 47 :$	46f93b68	7316056a	b0f40cfa	cf2bb5eb	7b074f58
$t = 48 :$	dc8e7f26	46f93b68	9cc5815a	b0f40cfa	cf2bb5eb
$t = 49 :$	850d411c	dc8e7f26	11be4eda	9cc5815a	b0f40cfa
$t = 50 :$	7e4672c0	850d411c	b7239fc9	11be4eda	9cc5815a
$t = 51 :$	89fbd41d	7e4672c0	21435047	b7239fc9	11be4eda
$t = 52 :$	1797e228	89fbd41d	1f919cb0	21435047	b7239fc9
$t = 53 :$	431d65bc	1797e228	627ef507	1f919cb0	21435047
$t = 54 :$	2bdbb8cb	431d65bc	05e5f88a	627ef507	1f919cb0
$t = 55 :$	6da72e7f	2bdbb8cb	10c7596f	05e5f88a	627ef507
$t = 56 :$	a8495a9b	6da72e7f	caf6ee32	10c7596f	05e5f88a
$t = 57 :$	e785655a	a8495a9b	db69cb9f	caf6ee32	10c7596f
$t = 58 :$	5b086c42	e785655a	ea1256a6	db69cb9f	caf6ee32
$t = 59 :$	a65818f7	5b086c42	b9e15956	ea1256a6	db69cb9f
$t = 60 :$	7aab101b	a65818f7	96c21b10	b9e15956	ea1256a6
$t = 61 :$	93614c9c	7aab101b	e996063d	96c21b10	b9e15956
$t = 62 :$	f66d9bf4	93614c9c	deaac406	e996063d	96c21b10
$t = 63 :$	d504902b	f66d9bf4	24d85327	deaac406	e996063d
$t = 64 :$	60a9da62	d504902b	3d9b66fd	24d85327	deaac406
$t = 65 :$	8b687819	60a9da62	f541240a	3d9b66fd	24d85327
$t = 66 :$	083e90c3	8b687819	982a7698	f541240a	3d9b66fd
$t = 67 :$	f6226bbf	083e90c3	62dale06	982a7698	f541240a
$t = 68 :$	76c0563b	f6226bbf	c20fa430	62dale06	982a7698
$t = 69 :$	989dd165	76c0563b	fd889aef	c20fa430	62dale06
$t = 70 :$	8b2c7573	989dd165	ddb0158e	fd889aef	c20fa430
$t = 71 :$	aelb8e7b	8b2c7573	66277459	ddb0158e	fd889aef
$t = 72 :$	ca1840de	aelb8e7b	e2cb1d5c	66277459	ddb0158e
$t = 73 :$	16f3babb	ca1840de	eb86e39e	e2cb1d5c	66277459
$t = 74 :$	d28d83ad	16f3babb	b2861037	eb86e39e	e2cb1d5c
$t = 75 :$	6bc02dfe	d28d83ad	c5bceae	b2861037	eb86e39e
$t = 76 :$	d3a6e275	6bc02dfe	74a360eb	c5bceae	b2861037
$t = 77 :$	da955482	d3a6e275	9af00b7f	74a360eb	c5bceae

$t = 78 :$	58c0aac0	da955482	74e9b89d	9af00b7f	74a360eb
$t = 79 :$	906fd62c	58c0aac0	b6a55520	74e9b89d	9af00b7f

That completes the processing of the second and final message block,  $M^{(2)}$ . The final hash value,  $H^{(2)}$ , is calculated to be

$$\begin{aligned}
 H_0^{(1)} &= \text{f4286818} + \text{906fd62c} = \text{84983e44} \\
 H_1^{(1)} &= \text{c37b27ae} + \text{58c0aac0} = \text{1c3bd26e} \\
 H_2^{(1)} &= \text{0408f581} + \text{b6a55520} = \text{baae4aa1} \\
 H_3^{(1)} &= \text{84677148} + \text{74e9b89d} = \text{f95129e5} \\
 H_4^{(1)} &= \text{4a566572} + \text{9af00b7f} = \text{e54670f1}.
 \end{aligned}$$

The resulting 160-bit message digest is

84983e44 1c3bd26e baae4aa1 f95129e5 e54670f1.

### A.3 SHA-1 Example (Long Message)

Let the message  $M$  be the binary-coded form of the ASCII string which consists of 1,000,000 repetitions of the character “a”. The resulting SHA-1 message digest is

34aa973c d4c4daa4 f61eeb2b dbad2731 6534016f.

## APPENDIX B: SHA-256 EXAMPLES

This appendix is for informational purposes only and is not required to meet the standard.

### B.1 SHA-256 Example (One-Block Message)

Let the message,  $M$ , be the 24-bit ( $\ell = 24$ ) ASCII string "abc", which is equivalent to the following binary string:

```
01100001 01100010 01100011.
```

The message is padded by appending a "1" bit, followed by 423 "0" bits, and ending with the hex value 00000000 00000018 (the two 32-bit word representation of the length, 24). Thus, the final padded message consists of one block ( $N = 1$ ).

For SHA-256, the initial hash value,  $H^{(0)}$ , is

```

 $H_0^{(0)} = 6a09e667$ 
 $H_1^{(0)} = bb67ae85$ 
 $H_2^{(0)} = 3c6ef372$ 
 $H_3^{(0)} = a54ff53a$ 
 $H_4^{(0)} = 510e527f$ 
 $H_5^{(0)} = 9b05688c$ 
 $H_6^{(0)} = 1f83d9ab$ 
 $H_7^{(0)} = 5be0cd19.$ 

```

The words of the padded message block are then assigned to the words  $W_0, \dots, W_{15}$  of the message schedule:

```

 $W_0 = 61626380$ 
 $W_1 = 00000000$ 
 $W_2 = 00000000$ 
 $W_3 = 00000000$ 
 $W_4 = 00000000$ 
 $W_5 = 00000000$ 
 $W_6 = 00000000$ 
 $W_7 = 00000000$ 
 $W_8 = 00000000$ 
 $W_9 = 00000000$ 
 $W_{10} = 00000000$ 
 $W_{11} = 00000000$ 
 $W_{12} = 00000000$ 
 $W_{13} = 00000000$ 
 $W_{14} = 00000000$ 
 $W_{15} = 00000018.$ 

```

The following schedule shows the hex values for  $a, b, c, d, e, f, g,$  and  $h$  after pass  $t$  of the "for  $t = 0$  to 63" loop described in Sec. 6.2.2, step 4.

	<i>a</i>	<i>b</i>	<i>c</i>	<i>d</i>	<i>e</i>	<i>f</i>	<i>g</i>	<i>h</i>
<i>t</i> = 0 :	5d6aebcd	6a09e667	bb67ae85	3c6ef372	fa2a4622	510e527f	9b05688c	1f83d9ab
<i>t</i> = 1 :	5a6ad9ad	5d6aebcd	6a09e667	bb67ae85	78ce7989	fa2a4622	510e527f	9b05688c
<i>t</i> = 2 :	c8c347a7	5a6ad9ad	5d6aebcd	6a09e667	f92939eb	78ce7989	fa2a4622	510e527f
<i>t</i> = 3 :	d550f666	c8c347a7	5a6ad9ad	5d6aebcd	24e00850	f92939eb	78ce7989	fa2a4622
<i>t</i> = 4 :	04409a6a	d550f666	c8c347a7	5a6ad9ad	43ada245	24e00850	f92939eb	78ce7989
<i>t</i> = 5 :	2b4209f5	04409a6a	d550f666	c8c347a7	714260ad	43ada245	24e00850	f92939eb
<i>t</i> = 6 :	e5030380	2b4209f5	04409a6a	d550f666	9b27a401	714260ad	43ada245	24e00850
<i>t</i> = 7 :	85a07b5f	e5030380	2b4209f5	04409a6a	0c657a79	9b27a401	714260ad	43ada245
<i>t</i> = 8 :	8e04ecb9	85a07b5f	e5030380	2b4209f5	32ca2d8c	0c657a79	9b27a401	714260ad
<i>t</i> = 9 :	8c87346b	8e04ecb9	85a07b5f	e5030380	1cc92596	32ca2d8c	0c657a79	9b27a401
<i>t</i> = 10 :	4798a3f4	8c87346b	8e04ecb9	85a07b5f	436b23e8	1cc92596	32ca2d8c	0c657a79
<i>t</i> = 11 :	f71fc5a9	4798a3f4	8c87346b	8e04ecb9	816fd6e9	436b23e8	1cc92596	32ca2d8c
<i>t</i> = 12 :	87912990	f71fc5a9	4798a3f4	8c87346b	1e578218	816fd6e9	436b23e8	1cc92596
<i>t</i> = 13 :	d932eb16	87912990	f71fc5a9	4798a3f4	745a48de	1e578218	816fd6e9	436b23e8
<i>t</i> = 14 :	c0645fde	d932eb16	87912990	f71fc5a9	0b92f20c	745a48de	1e578218	816fd6e9
<i>t</i> = 15 :	b0fa238e	c0645fde	d932eb16	87912990	07590dcd	0b92f20c	745a48de	1e578218
<i>t</i> = 16 :	21da9a9b	b0fa238e	c0645fde	d932eb16	8034229c	07590dcd	0b92f20c	745a48de
<i>t</i> = 17 :	c2fbd9d1	21da9a9b	b0fa238e	c0645fde	846ee454	8034229c	07590dcd	0b92f20c
<i>t</i> = 18 :	fe777bbf	c2fbd9d1	21da9a9b	b0fa238e	cc899961	846ee454	8034229c	07590dcd
<i>t</i> = 19 :	e1f20c33	fe777bbf	c2fbd9d1	21da9a9b	b0638179	cc899961	846ee454	8034229c
<i>t</i> = 20 :	9dc68b63	e1f20c33	fe777bbf	c2fbd9d1	8ada8930	b0638179	cc899961	846ee454
<i>t</i> = 21 :	c2606d6d	9dc68b63	e1f20c33	fe777bbf	e1257970	8ada8930	b0638179	cc899961
<i>t</i> = 22 :	a7a3623f	c2606d6d	9dc68b63	e1f20c33	49f5114a	e1257970	8ada8930	b0638179
<i>t</i> = 23 :	c5d53d8d	a7a3623f	c2606d6d	9dc68b63	aa47c347	49f5114a	e1257970	8ada8930
<i>t</i> = 24 :	1c2c2838	c5d53d8d	a7a3623f	c2606d6d	2823ef91	aa47c347	49f5114a	e1257970
<i>t</i> = 25 :	cde8037d	1c2c2838	c5d53d8d	a7a3623f	14383d8e	2823ef91	aa47c347	49f5114a
<i>t</i> = 26 :	b62ec4bc	cde8037d	1c2c2838	c5d53d8d	c74c6516	14383d8e	2823ef91	aa47c347
<i>t</i> = 27 :	77d37528	b62ec4bc	cde8037d	1c2c2838	edffbf8	c74c6516	14383d8e	2823ef91
<i>t</i> = 28 :	363482c9	77d37528	b62ec4bc	cde8037d	6112a3b7	edffbf8	c74c6516	14383d8e
<i>t</i> = 29 :	a0060b30	363482c9	77d37528	b62ec4bc	ade79437	6112a3b7	edffbf8	c74c6516
<i>t</i> = 30 :	ea992a22	a0060b30	363482c9	77d37528	0109ab3a	ade79437	6112a3b7	edffbf8
<i>t</i> = 31 :	73b33bf5	ea992a22	a0060b30	363482c9	ba591112	0109ab3a	ade79437	6112a3b7
<i>t</i> = 32 :	98e12507	73b33bf5	ea992a22	a0060b30	9cd9f5f6	ba591112	0109ab3a	ade79437
<i>t</i> = 33 :	fe604df5	98e12507	73b33bf5	ea992a22	59249dd3	9cd9f5f6	ba591112	0109ab3a
<i>t</i> = 34 :	a9a7738c	fe604df5	98e12507	73b33bf5	085f3833	59249dd3	9cd9f5f6	ba591112
<i>t</i> = 35 :	65a0cfe4	a9a7738c	fe604df5	98e12507	f4b002d6	085f3833	59249dd3	9cd9f5f6
<i>t</i> = 36 :	41a65cb1	65a0cfe4	a9a7738c	fe604df5	0772a26b	f4b002d6	085f3833	59249dd3
<i>t</i> = 37 :	34df1604	41a65cb1	65a0cfe4	a9a7738c	a507a53d	0772a26b	f4b002d6	085f3833
<i>t</i> = 38 :	6dc57a8a	34df1604	41a65cb1	65a0cfe4	f0781bc8	a507a53d	0772a26b	f4b002d6
<i>t</i> = 39 :	79ea687a	6dc57a8a	34df1604	41a65cb1	1efbc0a0	f0781bc8	a507a53d	0772a26b
<i>t</i> = 40 :	d6670766	79ea687a	6dc57a8a	34df1604	26352d63	1efbc0a0	f0781bc8	a507a53d
<i>t</i> = 41 :	df46652f	d6670766	79ea687a	6dc57a8a	838b2711	26352d63	1efbc0a0	f0781bc8
<i>t</i> = 42 :	17aa0dfe	df46652f	d6670766	79ea687a	decd4715	838b2711	26352d63	1efbc0a0
<i>t</i> = 43 :	9d4baf93	17aa0dfe	df46652f	d6670766	fda24c2e	decd4715	838b2711	26352d63
<i>t</i> = 44 :	26628815	9d4baf93	17aa0dfe	df46652f	a80f11f0	fda24c2e	decd4715	838b2711
<i>t</i> = 45 :	72ab4b91	26628815	9d4baf93	17aa0dfe	b7755da1	a80f11f0	fda24c2e	decd4715
<i>t</i> = 46 :	a14c14b0	72ab4b91	26628815	9d4baf93	d57b94a9	b7755da1	a80f11f0	fda24c2e
<i>t</i> = 47 :	4172328d	a14c14b0	72ab4b91	26628815	fecf0bc6	d57b94a9	b7755da1	a80f11f0
<i>t</i> = 48 :	05757ceb	4172328d	a14c14b0	72ab4b91	bd714038	fecf0bc6	d57b94a9	b7755da1
<i>t</i> = 49 :	f11bfaa8	05757ceb	4172328d	a14c14b0	6e5c390c	bd714038	fecf0bc6	d57b94a9
<i>t</i> = 50 :	7a0508a1	f11bfaa8	05757ceb	4172328d	52f1ccf7	6e5c390c	bd714038	fecf0bc6
<i>t</i> = 51 :	886e7a22	7a0508a1	f11bfaa8	05757ceb	49231c1e	52f1ccf7	6e5c390c	bd714038

```

t = 52 : 101fd28f 886e7a22 7a0508a1 f11bfaa8 529e7d00 49231c1e 52f1ccf7 6e5c390c
t = 53 : f5702fdb 101fd28f 886e7a22 7a0508a1 9f4787c3 529e7d00 49231c1e 52f1ccf7
t = 54 : 3ec45cdb f5702fdb 101fd28f 886e7a22 e50e1b4f 9f4787c3 529e7d00 49231c1e
t = 55 : 38cc9913 3ec45cdb f5702fdb 101fd28f 54cb266b e50e1b4f 9f4787c3 529e7d00
t = 56 : fcd1887b 38cc9913 3ec45cdb f5702fdb 9b5e906c 54cb266b e50e1b4f 9f4787c3
t = 57 : c062d46f fcd1887b 38cc9913 3ec45cdb 7e44008e 9b5e906c 54cb266b e50e1b4f
t = 58 : ffb70472 c062d46f fcd1887b 38cc9913 6d83bfc6 7e44008e 9b5e906c 54cb266b
t = 59 : b6ae8fff ffb70472 c062d46f fcd1887b b21bad3d 6d83bfc6 7e44008e 9b5e906c
t = 60 : b85e2ce9 b6ae8fff ffb70472 c062d46f 961f4894 b21bad3d 6d83bfc6 7e44008e
t = 61 : 04d24d6c b85e2ce9 b6ae8fff ffb70472 948d25b6 961f4894 b21bad3d 6d83bfc6
t = 62 : d39a2165 04d24d6c b85e2ce9 b6ae8fff fb121210 948d25b6 961f4894 b21bad3d
t = 63 : 506e3058 d39a2165 04d24d6c b85e2ce9 5ef50f24 fb121210 948d25b6 961f4894

```

That completes the processing of the first and only message block,  $M^{(1)}$ . The final hash value,  $H^{(1)}$ , is calculated to be

$$\begin{aligned}
H_0^{(1)} &= 6a09e667 + 506e3058 = ba7816bf \\
H_1^{(1)} &= bb67ae85 + d39a2165 = 8f01cfea \\
H_2^{(1)} &= 3c6ef372 + 04d24d6c = 414140de \\
H_3^{(1)} &= a54ff53a + b85e2ce9 = 5dae2223 \\
H_4^{(1)} &= 510e527f + 5ef50f24 = b00361a3 \\
H_5^{(1)} &= 9b05688c + fb121210 = 96177a9c \\
H_6^{(1)} &= 1f83d9ab + 948d25b6 = b410ff61 \\
H_7^{(1)} &= 5be0cd19 + 961f4894 = f20015ad.
\end{aligned}$$

The resulting 256-bit message digest is

```
ba7816bf 8f01cfea 414140de 5dae2223 b00361a3 96177a9c b410ff61 f20015ad.
```

## B.2 SHA-256 Example (Multi-Block Message)

Let the message,  $M$ , be the 448-bit ( $\ell = 448$ ) ASCII string

**"abcdbcdecdefdefgefghfghighijhijkijkljklmklmnlmnomnopnopq".**

The message is padded by appending a "1" bit, followed by 511 "0" bits, and ending with the hex value 00000000 000001c0 (the two 32-bit word representation of the length, 448). Thus, the final padded message consists of two blocks ( $N = 2$ ).

For SHA-256, the initial hash value,  $H^{(0)}$ , is

$$\begin{aligned}
H_0^{(0)} &= 6a09e667 \\
H_1^{(0)} &= bb67ae85 \\
H_2^{(0)} &= 3c6ef372
\end{aligned}$$

$$\begin{aligned}
 H_3^{(0)} &= \text{a54ff53a} \\
 H_4^{(0)} &= \text{510e527f} \\
 H_5^{(0)} &= \text{9b05688c} \\
 H_6^{(0)} &= \text{1f83d9ab} \\
 H_7^{(0)} &= \text{5be0cd19}.
 \end{aligned}$$

The words of the first padded message block,  $M^{(1)}$ , are then assigned to the words  $W_0, \dots, W_{15}$  of the message schedule:

$$\begin{array}{ll}
 W_0 = 61626364 & W_8 = 696a6b6c \\
 W_1 = 62636465 & W_9 = 6a6b6c6d \\
 W_2 = 63646566 & W_{10} = 6b6c6d6e \\
 W_3 = 64656667 & W_{11} = 6c6d6e6f \\
 W_4 = 65666768 & W_{12} = 6d6e6f70 \\
 W_5 = 66676869 & W_{13} = 6e6f7071 \\
 W_6 = 6768696a & W_{14} = 80000000 \\
 W_7 = 68696a6b & W_{15} = 00000000.
 \end{array}$$

The following schedule shows the hex values for  $a, b, c, d, e, f, g,$  and  $h$  after pass  $t$  of the “for  $t = 0$  to 63” loop described in Sec. 6.2.2, step 4.

	$a$	$b$	$c$	$d$	$e$	$f$	$g$	$h$
$t = 0$ :	5d6aebb1	6a09e667	bb67ae85	3c6ef372	fa2a4606	510e527f	9b05688c	1f83d9ab
$t = 1$ :	2f2d5fcf	5d6aebb1	6a09e667	bb67ae85	4eblcfce	fa2a4606	510e527f	9b05688c
$t = 2$ :	97651825	2f2d5fcf	5d6aebb1	6a09e667	62d5c49e	4eblcfce	fa2a4606	510e527f
$t = 3$ :	4a8d64d5	97651825	2f2d5fcf	5d6aebb1	6494841b	62d5c49e	4eblcfce	fa2a4606
$t = 4$ :	f921c212	4a8d64d5	97651825	2f2d5fcf	05c4f88a	6494841b	62d5c49e	4eblcfce
$t = 5$ :	55c8ef48	f921c212	4a8d64d5	97651825	7ff91c94	05c4f88a	6494841b	62d5c49e
$t = 6$ :	485835b7	55c8ef48	f921c212	4a8d64d5	39a5b2ca	7ff91c94	05c4f88a	6494841b
$t = 7$ :	d237e6db	485835b7	55c8ef48	f921c212	a401d211	39a5b2ca	7ff91c94	05c4f88a
$t = 8$ :	359f2bce	d237e6db	485835b7	55c8ef48	c09ffec4	a401d211	39a5b2ca	7ff91c94
$t = 9$ :	3a474b2b	359f2bce	d237e6db	485835b7	9037b3b8	c09ffec4	a401d211	39a5b2ca
$t = 10$ :	b8e2b4cb	3a474b2b	359f2bce	d237e6db	443ed29e	9037b3b8	c09ffec4	a401d211
$t = 11$ :	1762215c	b8e2b4cb	3a474b2b	359f2bce	ee1c97a8	443ed29e	9037b3b8	c09ffec4
$t = 12$ :	101a4861	1762215c	b8e2b4cb	3a474b2b	839a0fc9	ee1c97a8	443ed29e	9037b3b8
$t = 13$ :	d68e6457	101a4861	1762215c	b8e2b4cb	9243f8af	839a0fc9	ee1c97a8	443ed29e
$t = 14$ :	ddl6cbb3	d68e6457	101a4861	1762215c	9162aded	9243f8af	839a0fc9	ee1c97a8
$t = 15$ :	c3486194	ddl6cbb3	d68e6457	101a4861	1496a54f	9162aded	9243f8af	839a0fc9
$t = 16$ :	b9dcacb1	c3486194	ddl6cbb3	d68e6457	d4f64250	1496a54f	9162aded	9243f8af
$t = 17$ :	046a193e	b9dcacb1	c3486194	ddl6cbb3	885370b6	d4f64250	1496a54f	9162aded
$t = 18$ :	f402f058	046a193e	b9dcacb1	c3486194	6f433549	885370b6	d4f64250	1496a54f
$t = 19$ :	2139187b	f402f058	046a193e	b9dcacb1	7c304206	6f433549	885370b6	d4f64250
$t = 20$ :	d70ac17d	2139187b	f402f058	046a193e	7cc6b262	7c304206	6f433549	885370b6
$t = 21$ :	1b2b66b8	d70ac17d	2139187b	f402f058	d560b028	7cc6b262	7c304206	6f433549
$t = 22$ :	ae2e2d4f	1b2b66b8	d70ac17d	2139187b	f074fc95	d560b028	7cc6b262	7c304206
$t = 23$ :	59fce6b9	ae2e2d4f	1b2b66b8	d70ac17d	a2c7d51d	f074fc95	d560b028	7cc6b262
$t = 24$ :	4a885065	59fce6b9	ae2e2d4f	1b2b66b8	763597fb	a2c7d51d	f074fc95	d560b028



```

t = 25 : 573221da 4a885065 59fce6b9 ae2e2d4f 36e74eb4 763597fb a2c7d51d f074fc95
t = 26 : 128661da 573221da 4a885065 59fce6b9 1162d575 36e74eb4 763597fb a2c7d51d
t = 27 : 73f858af 128661da 573221da 4a885065 e77c797f 1162d575 36e74eb4 763597fb
t = 28 : 74bcf468 73f858af 128661da 573221da 72abaecd e77c797f 1162d575 36e74eb4
t = 29 : df7151a0 74bcf468 73f858af 128661da 7629c961 72abaecd e77c797f 1162d575
t = 30 : eb43f3ed df7151a0 74bcf468 73f858af 0635d880 7629c961 72abaecd e77c797f
t = 31 : 5581ab07 eb43f3ed df7151a0 74bcf468 df980085 0635d880 7629c961 72abaecd
t = 32 : 9fc905c8 5581ab07 eb43f3ed df7151a0 a94d2af1 df980085 0635d880 7629c961
t = 33 : 9ce5a62f 9fc905c8 5581ab07 eb43f3ed 6ef3b6bd a94d2af1 df980085 0635d880
t = 34 : 1df8e885 9ce5a62f 9fc905c8 5581ab07 2a9e048e 6ef3b6bd a94d2af1 df980085
t = 35 : 0786dce8 1df8e885 9ce5a62f 9fc905c8 de2a21d1 2a9e048e 6ef3b6bd a94d2af1
t = 36 : 2c55d3a6 0786dce8 1df8e885 9ce5a62f b067c1af de2a21d1 2a9e048e 6ef3b6bd
t = 37 : a985b4be 2c55d3a6 0786dce8 1df8e885 f72bf353 b067c1af de2a21d1 2a9e048e
t = 38 : 91ac9d5d a985b4be 2c55d3a6 0786dce8 68d8d590 f72bf353 b067c1af de2a21d1
t = 39 : 7e4d30b8 91ac9d5d a985b4be 2c55d3a6 9f5b9b6d 68d8d590 f72bf353 b067c1af
t = 40 : 7e056794 7e4d30b8 91ac9d5d a985b4be 423b26c0 9f5b9b6d 68d8d590 f72bf353
t = 41 : 508a16ab 7e056794 7e4d30b8 91ac9d5d 45459d97 423b26c0 9f5b9b6d 68d8d590
t = 42 : b62c7013 508a16ab 7e056794 7e4d30b8 80a92a00 45459d97 423b26c0 9f5b9b6d
t = 43 : 167361de b62c7013 508a16ab 7e056794 41dd3844 80a92a00 45459d97 423b26c0
t = 44 : de71e2f2 167361de b62c7013 508a16ab ff61c636 41dd3844 80a92a00 45459d97
t = 45 : 18f0d19d de71e2f2 167361de b62c7013 6b88472c ff61c636 41dd3844 80a92a00
t = 46 : 165be9cd 18f0d19d de71e2f2 167361de a483f080 6b88472c ff61c636 41dd3844
t = 47 : 13d82741 165be9cd 18f0d19d de71e2f2 a7802a4d a483f080 6b88472c ff61c636
t = 48 : 017b9d99 13d82741 165be9cd 18f0d19d aeb10b60 a7802a4d a483f080 6b88472c
t = 49 : 543c99a1 017b9d99 13d82741 165be9cd 16f134b6 aeb10b60 a7802a4d a483f080
t = 50 : 758ca97a 543c99a1 017b9d99 13d82741 100cf2ea 16f134b6 aeb10b60 a7802a4d
t = 51 : 81c1cde0 758ca97a 543c99a1 017b9d99 5c47eb7b 100cf2ea 16f134b6 aeb10b60
t = 52 : b8d55619 81c1cde0 758ca97a 543c99a1 1c806a61 5c47eb7b 100cf2ea 16f134b6
t = 53 : 1d6de87a b8d55619 81c1cde0 758ca97a 3443bed4 1c806a61 5c47eb7b 100cf2ea
t = 54 : f907b313 1d6de87a b8d55619 81c1cde0 61a41711 3443bed4 1c806a61 5c47eb7b
t = 55 : 9e57c4a0 f907b313 1d6de87a b8d55619 eec13548 61a41711 3443bed4 1c806a61
t = 56 : 71629856 9e57c4a0 f907b313 1d6de87a 2f6c8c4e eec13548 61a41711 3443bed4
t = 57 : 7c015a2c 71629856 9e57c4a0 f907b313 cb9d3dd0 2f6c8c4e eec13548 61a41711
t = 58 : 921fccb6 7c015a2c 71629856 9e57c4a0 43d8a034 cb9d3dd0 2f6c8c4e eec13548
t = 59 : e18f259a 921fccb6 7c015a2c 71629856 51e15869 43d8a034 cb9d3dd0 2f6c8c4e
t = 60 : bcfce922 e18f259a 921fccb6 7c015a2c 962d8621 51e15869 43d8a034 cb9d3dd0
t = 61 : f6f443f8 bcfce922 e18f259a 921fccb6 acc75916 962d8621 51e15869 43d8a034
t = 62 : 86126910 f6f443f8 bcfce922 e18f259a 2fc08f85 acc75916 962d8621 51e15869
t = 63 : 1bdc6f6f 86126910 f6f443f8 bcfce922 25d2430a 2fc08f85 acc75916 962d8621

```

That completes the processing of the first message block,  $M^{(1)}$ . The first intermediate hash value,  $H^{(1)}$ , is calculated to be

$$H_0^{(1)} = 6a09e667 + 1bdc6f6f = 85e655d6$$

$$H_1^{(1)} = bb67ae85 + 86126910 = 417a1795$$

$$H_2^{(1)} = 3c6ef372 + f6f443f8 = 3363376a$$

$$H_3^{(1)} = a54fff53a + bcfce922 = 624cde5c$$

$$H_4^{(1)} = 510e527f + 25d2430a = 76e09589$$

$$H_5^{(1)} = 9b05688c + 2fc08f85 = cac5f811$$

$$H_6^{(1)} = 1f83d9ab + acc75916 = cc4b32c1$$

$$H_7^{(1)} = 5be0cd19 + 962d8621 = f20e533a.$$

The words of the *second* padded message block,  $M^{(2)}$ , are then assigned to the words  $W_0, \dots, W_{15}$  of the message schedule:

$W_0$	=	00000000	$W_8$	=	00000000
$W_1$	=	00000000	$W_9$	=	00000000
$W_2$	=	00000000	$W_{10}$	=	00000000
$W_3$	=	00000000	$W_{11}$	=	00000000
$W_4$	=	00000000	$W_{12}$	=	00000000
$W_5$	=	00000000	$W_{13}$	=	00000000
$W_6$	=	00000000	$W_{14}$	=	00000000
$W_7$	=	00000000	$W_{15}$	=	000001c0.

The following schedule shows the hex values for  $a, b, c, d, e, f, g,$  and  $h$  after pass  $t$  of the “for  $t = 0$  to 63” loop described in Sec. 6.2.2, step 4.

	$a$	$b$	$c$	$d$	$e$	$f$	$g$	$h$
$t = 0$ :	7c20c838	85e655d6	417a1795	3363376a	4670ae6e	76e09589	cac5f811	cc4b32c1
$t = 1$ :	7c3c0f86	7c20c838	85e655d6	417a1795	8c51be64	4670ae6e	76e09589	cac5f811
$t = 2$ :	fd1eebdc	7c3c0f86	7c20c838	85e655d6	af71b9ea	8c51be64	4670ae6e	76e09589
$t = 3$ :	f268faa9	fd1eebdc	7c3c0f86	7c20c838	e20362ef	af71b9ea	8c51be64	4670ae6e
$t = 4$ :	185a5d79	f268faa9	fd1eebdc	7c3c0f86	8dff3001	e20362ef	af71b9ea	8c51be64
$t = 5$ :	3eeb6c06	185a5d79	f268faa9	fd1eebdc	fe20cda6	8dff3001	e20362ef	af71b9ea
$t = 6$ :	89bba3f1	3eeb6c06	185a5d79	f268faa9	0a34df03	fe20cda6	8dff3001	e20362ef
$t = 7$ :	bf9a93a0	89bba3f1	3eeb6c06	185a5d79	059abdd1	0a34df03	fe20cda6	8dff3001
$t = 8$ :	2c096744	bf9a93a0	89bba3f1	3eeb6c06	abfa465b	059abdd1	0a34df03	fe20cda6
$t = 9$ :	2d964e86	2c096744	bf9a93a0	89bba3f1	aa27ed82	abfa465b	059abdd1	0a34df03
$t = 10$ :	5b35025b	2d964e86	2c096744	bf9a93a0	10e77723	aa27ed82	abfa465b	059abdd1
$t = 11$ :	5eb4ec40	5b35025b	2d964e86	2c096744	e11b4548	10e77723	aa27ed82	abfa465b
$t = 12$ :	35ee996d	5eb4ec40	5b35025b	2d964e86	5c24e2a2	e11b4548	10e77723	aa27ed82
$t = 13$ :	d74080fa	35ee996d	5eb4ec40	5b35025b	68aa893f	5c24e2a2	e11b4548	10e77723
$t = 14$ :	0cea5cbc	d74080fa	35ee996d	5eb4ec40	60356548	68aa893f	5c24e2a2	e11b4548
$t = 15$ :	16a8cc79	0cea5cbc	d74080fa	35ee996d	0fcb1f6f	60356548	68aa893f	5c24e2a2
$t = 16$ :	f16f634e	16a8cc79	0cea5cbc	d74080fa	8b21cdc1	0fcb1f6f	60356548	68aa893f
$t = 17$ :	23dcb6c2	f16f634e	16a8cc79	0cea5cbc	ca9182d3	8b21cdc1	0fcb1f6f	60356548
$t = 18$ :	dcff40fd	23dcb6c2	f16f634e	16a8cc79	69bf7b95	ca9182d3	8b21cdc1	0fcb1f6f
$t = 19$ :	76f1a2bc	dcff40fd	23dcb6c2	f16f634e	0dc84bb1	69bf7b95	ca9182d3	8b21cdc1
$t = 20$ :	20aad899	76f1a2bc	dcff40fd	23dcb6c2	cc4769f2	0dc84bb1	69bf7b95	ca9182d3
$t = 21$ :	d44dc81a	20aad899	76f1a2bc	dcff40fd	5bace62d	cc4769f2	0dc84bb1	69bf7b95
$t = 22$ :	f13ae55b	d44dc81a	20aad899	76f1a2bc	966aa287	5bace62d	cc4769f2	0dc84bb1
$t = 23$ :	a4195b91	f13ae55b	d44dc81a	20aad899	eddbd6ed	966aa287	5bace62d	cc4769f2
$t = 24$ :	4984fa79	a4195b91	f13ae55b	d44dc81a	a530d939	eddbd6ed	966aa287	5bace62d
$t = 25$ :	aa6cb982	4984fa79	a4195b91	f13ae55b	0b5eeea4	a530d939	eddbd6ed	966aa287
$t = 26$ :	9450fbbc	aa6cb982	4984fa79	a4195b91	09166dda	0b5eeea4	a530d939	eddbd6ed
$t = 27$ :	0d936bab	9450fbbc	aa6cb982	4984fa79	6e495d4b	09166dda	0b5eeea4	a530d939
$t = 28$ :	d958b529	0d936bab	9450fbbc	aa6cb982	c2fa99b1	6e495d4b	09166dda	0b5eeea4
$t = 29$ :	1cfa5eb0	d958b529	0d936bab	9450fbbc	6c49db9f	c2fa99b1	6e495d4b	09166dda
$t = 30$ :	02ef3a5f	1cfa5eb0	d958b529	0d936bab	5da10665	6c49db9f	c2fa99b1	6e495d4b
$t = 31$ :	b0eab1c5	02ef3a5f	1cfa5eb0	d958b529	f6d93952	5da10665	6c49db9f	c2fa99b1

```

t = 32 : 0bfba73c b0eab1c5 02ef3a5f 1cfa5eb0 8b99e3a9 f6d93952 5da10665 6c49db9f
t = 33 : 4bd1df96 0bfba73c b0eab1c5 02ef3a5f 905e44ac 8b99e3a9 f6d93952 5da10665
t = 34 : 9907f1b6 4bd1df96 0bfba73c b0eab1c5 66c3043d 905e44ac 8b99e3a9 f6d93952
t = 35 : ecde4e0d 9907f1b6 4bd1df96 0bfba73c 5dc119e6 66c3043d 905e44ac 8b99e3a9
t = 36 : 2f11c939 ecde4e0d 9907f1b6 4bd1df96 fed4ce1d 5dc119e6 66c3043d 905e44ac
t = 37 : d949682b 2f11c939 ecde4e0d 9907f1b6 32d99008 fed4ce1d 5dc119e6 66c3043d
t = 38 : adca7a96 d949682b 2f11c939 ecde4e0d c6cce4ff 32d99008 fed4ce1d 5dc119e6
t = 39 : 221b8a5a adca7a96 d949682b 2f11c939 0b82c5eb c6cce4ff 32d99008 fed4ce1d
t = 40 : 12d97845 221b8a5a adca7a96 d949682b e4213ca2 0b82c5eb c6cce4ff 32d99008
t = 41 : 2c794876 12d97845 221b8a5a adca7a96 ff6759ba e4213ca2 0b82c5eb c6cce4ff
t = 42 : 8300fca2 2c794876 12d97845 221b8a5a e0e3457c ff6759ba e4213ca2 0b82c5eb
t = 43 : f2ad6322 8300fca2 2c794876 12d97845 cc48c7f3 e0e3457c ff6759ba e4213ca2
t = 44 : 0f154e11 f2ad6322 8300fca2 2c794876 6f9517cb cc48c7f3 e0e3457c ff6759ba
t = 45 : 104a7db4 0f154e11 f2ad6322 8300fca2 5348e8f6 6f9517cb cc48c7f3 e0e3457c
t = 46 : 0b3303a7 104a7db4 0f154e11 f2ad6322 bbe1c39a 5348e8f6 6f9517cb cc48c7f3
t = 47 : d7354d5b 0b3303a7 104a7db4 0f154e11 aad55b6b bbe1c39a 5348e8f6 6f9517cb
t = 48 : b736d7a6 d7354d5b 0b3303a7 104a7db4 68f25260 aad55b6b bbe1c39a 5348e8f6
t = 49 : 2748e5ec b736d7a6 d7354d5b 0b3303a7 d4b58576 68f25260 aad55b6b bbe1c39a
t = 50 : d8aabcf9 2748e5ec b736d7a6 d7354d5b 27844711 d4b58576 68f25260 aad55b6b
t = 51 : 1a6bcf6a d8aabcf9 2748e5ec b736d7a6 ff5e99d0 27844711 d4b58576 68f25260
t = 52 : 4eca6fa0 1a6bcf6a d8aabcf9 2748e5ec 989ed071 ff5e99d0 27844711 d4b58576
t = 53 : ec02560a 4eca6fa0 1a6bcf6a d8aabcf9 7151df8e 989ed071 ff5e99d0 27844711
t = 54 : d9f0c115 ec02560a 4eca6fa0 1a6bcf6a 624150c4 7151df8e 989ed071 ff5e99d0
t = 55 : 92952710 d9f0c115 ec02560a 4eca6fa0 226806d6 624150c4 7151df8e 989ed071
t = 56 : 20d4d0e4 92952710 d9f0c115 ec02560a 4e515a4d 226806d6 624150c4 7151df8e
t = 57 : 4348eb1f 20d4d0e4 92952710 d9f0c115 c21eddf9 4e515a4d 226806d6 624150c4
t = 58 : 286fe5f0 4348eb1f 20d4d0e4 92952710 54076664 c21eddf9 4e515a4d 226806d6
t = 59 : 1c4cddd9 286fe5f0 4348eb1f 20d4d0e4 f487a853 54076664 c21eddf9 4e515a4d
t = 60 : a9f181dd 1c4cddd9 286fe5f0 4348eb1f 27ccb387 f487a853 54076664 c21eddf9
t = 61 : b25cef29 a9f181dd 1c4cddd9 286fe5f0 2aa1bb13 27ccb387 f487a853 54076664
t = 62 : 908c2123 b25cef29 a9f181dd 1c4cddd9 9a392956 2aa1bb13 27ccb387 f487a853
t = 63 : 9ea7148b 908c2123 b25cef29 a9f181dd 2c5c4ed0 9a392956 2aa1bb13 27ccb387

```

That completes the processing of the second and final message block,  $M^{(2)}$ . The final hash value,  $H^{(2)}$ , is calculated to be

$$H_0^{(2)} = 85e655d6 + 9ea7148b = 248d6a61$$

$$H_1^{(2)} = 417a1795 + 908c2123 = d20638b8$$

$$H_2^{(2)} = 3363376a + b25cef29 = e5c02693$$

$$H_3^{(2)} = 624cde5c + a9f181dd = 0c3e6039$$

$$H_4^{(2)} = 76e09589 + 2c5c4ed0 = a33ce459$$

$$H_5^{(2)} = cac5f811 + 9a392956 = 64ff2167$$

$$H_6^{(2)} = cc4b32c1 + 2aa1bb13 = f6ecedd4$$

$$H_7^{(2)} = f20e533a + 27ccb387 = 19db06c1.$$

The resulting 256-bit message digest is

248d6a61 d20638b8 e5c02693 0c3e6039 a33ce459 64ff2167 f6ecedd4 19db06c1.

### **B.3 SHA-256 Example (Long Message)**

Let the message  $M$  be the binary-coded form of the ASCII string which consists of 1,000,000 repetitions of the character ‘**a**’. The resulting SHA-256 message digest is

```
cdc76e5c 9914fb92 81a1c7e2 84d73e67 f1809a48 a497200e 046d39cc c7112cd0.
```

## APPENDIX C: SHA-512 EXAMPLES

This appendix is for informational purposes only and is not required to meet the standard.

### C.1 SHA-512 Example (One-Block Message)

Let the message,  $M$ , be the 24-bit ( $\ell = 24$ ) ASCII string "abc", which is equivalent to the following binary string:

01100001 01100010 01100011.

The message is padded by appending a "1" bit, followed by 871 "0" bits, and ending with the hex value

0000000000000000 0000000000000018

(the two 64-bit word representation of the length, 24). Thus, the final padded message consists of one block ( $N = 1$ ).

For SHA-512, the initial hash value,  $H^{(0)}$ , is

$H_0^{(0)} = 6a09e667f3bcc908$   
 $H_1^{(0)} = bb67ae8584caa73b$   
 $H_2^{(0)} = 3c6ef372fe94f82b$   
 $H_3^{(0)} = a54ff53a5f1d36f1$   
 $H_4^{(0)} = 510e527fade682d1$   
 $H_5^{(0)} = 9b05688c2b3e6c1f$   
 $H_6^{(0)} = 1f83d9abfb41bd6b$   
 $H_7^{(0)} = 5be0cd19137e2179.$

The words of the padded message block are then assigned to the words  $W_0, \dots, W_{15}$  of the message schedule:

$W_0 = 6162638000000000$	$W_8 = 0000000000000000$
$W_1 = 0000000000000000$	$W_9 = 0000000000000000$
$W_2 = 0000000000000000$	$W_{10} = 0000000000000000$
$W_3 = 0000000000000000$	$W_{11} = 0000000000000000$
$W_4 = 0000000000000000$	$W_{12} = 0000000000000000$
$W_5 = 0000000000000000$	$W_{13} = 0000000000000000$
$W_6 = 0000000000000000$	$W_{14} = 0000000000000000$
$W_7 = 0000000000000000$	$W_{15} = 0000000000000018.$

The following schedule shows the hex values for *a*, *b*, *c*, *d*, *e*, *f*, *g*, and *h* after pass *t* of the “for *t* = 0 to 79” loop described in Sec. 6.3.2, step 4.

	<i>a</i> / <i>e</i>	<i>b</i> / <i>f</i>	<i>c</i> / <i>g</i>	<i>d</i> / <i>h</i>
<i>t</i> = 0 :	f6afceb8bcfcddf5 58cb02347ab51f91	6a09e667f3bcc908 510e527fade682d1	bb67ae8584caa73b 9b05688c2b3e6c1f	3c6ef372fe94f82b 1f83d9abfb41bd6b
<i>t</i> = 1 :	1320f8c9fb872cc0 c3d4ebfd48650ffa	f6afceb8bcfcddf5 58cb02347ab51f91	6a09e667f3bcc908 510e527fade682d1	bb67ae8584caa73b 9b05688c2b3e6c1f
<i>t</i> = 2 :	ebcfff07203d91f3 dfa9b239f2697812	1320f8c9fb872cc0 c3d4ebfd48650ffa	f6afceb8bcfcddf5 58cb02347ab51f91	6a09e667f3bcc908 510e527fade682d1
<i>t</i> = 3 :	5a83cb3e80050e82 0b47b4bb1928990e	ebcfff07203d91f3 dfa9b239f2697812	1320f8c9fb872cc0 c3d4ebfd48650ffa	f6afceb8bcfcddf5 58cb02347ab51f91
<i>t</i> = 4 :	b680953951604860 745aca4a342ed2e2	5a83cb3e80050e82 0b47b4bb1928990e	ebcfff07203d91f3 dfa9b239f2697812	1320f8c9fb872cc0 c3d4ebfd48650ffa
<i>t</i> = 5 :	af573b02403e89cd 96f60209b6dc35ba	b680953951604860 745aca4a342ed2e2	5a83cb3e80050e82 0b47b4bb1928990e	ebcfff07203d91f3 dfa9b239f2697812
<i>t</i> = 6 :	c4875b0c7abc076b 5a6c781f54dcc00c	af573b02403e89cd 96f60209b6dc35ba	b680953951604860 745aca4a342ed2e2	5a83cb3e80050e82 0b47b4bb1928990e
<i>t</i> = 7 :	8093d195e0054fa3 86f67263a0f0ec0a	c4875b0c7abc076b 5a6c781f54dcc00c	af573b02403e89cd 96f60209b6dc35ba	b680953951604860 745aca4a342ed2e2
<i>t</i> = 8 :	f1eca5544cb89225 d0403c398fc40002	8093d195e0054fa3 86f67263a0f0ec0a	c4875b0c7abc076b 5a6c781f54dcc00c	af573b02403e89cd 96f60209b6dc35ba
<i>t</i> = 9 :	81782d4a5db48f03 00091f460be46c52	f1eca5544cb89225 d0403c398fc40002	8093d195e0054fa3 86f67263a0f0ec0a	c4875b0c7abc076b 5a6c781f54dcc00c
<i>t</i> = 10 :	69854c4aa0f25b59 d375471bde1ba3f4	81782d4a5db48f03 00091f460be46c52	f1eca5544cb89225 d0403c398fc40002	8093d195e0054fa3 86f67263a0f0ec0a
<i>t</i> = 11 :	db0a9963f80c2eaa 475975b91a7a462c	69854c4aa0f25b59 d375471bde1ba3f4	81782d4a5db48f03 00091f460be46c52	f1eca5544cb89225 d0403c398fc40002
<i>t</i> = 12 :	5e41214388186c14 cdf3bff2883fc9d9	db0a9963f80c2eaa 475975b91a7a462c	69854c4aa0f25b59 d375471bde1ba3f4	81782d4a5db48f03 00091f460be46c52
<i>t</i> = 13 :	44249631255d2ca0 860acf9effba6f61	5e41214388186c14 cdf3bff2883fc9d9	db0a9963f80c2eaa 475975b91a7a462c	69854c4aa0f25b59 d375471bde1ba3f4
<i>t</i> = 14 :	fa967eed85a08028 874bfe5f6aae9f2f	44249631255d2ca0 860acf9effba6f61	5e41214388186c14 cdf3bff2883fc9d9	db0a9963f80c2eaa 475975b91a7a462c
<i>t</i> = 15 :	0ae07c86b1181c75 a77b7c035dd4c161	fa967eed85a08028 874bfe5f6aae9f2f	44249631255d2ca0 860acf9effba6f61	5e41214388186c14 cdf3bff2883fc9d9
<i>t</i> = 16 :	caf81a425d800537 2deecc6b39d64d78	0ae07c86b1181c75 a77b7c035dd4c161	fa967eed85a08028 874bfe5f6aae9f2f	44249631255d2ca0 860acf9effba6f61
<i>t</i> = 17 :	4725be249ad19e6b f47e8353f8047455	caf81a425d800537 2deecc6b39d64d78	0ae07c86b1181c75 a77b7c035dd4c161	fa967eed85a08028 874bfe5f6aae9f2f
<i>t</i> = 18 :	3c4b4104168e3edb 29695fd88d81dbd0	4725be249ad19e6b f47e8353f8047455	caf81a425d800537 2deecc6b39d64d78	0ae07c86b1181c75 a77b7c035dd4c161
<i>t</i> = 19 :	9a3fb4d38ab6cf06 f14998dd5f70767e	3c4b4104168e3edb 29695fd88d81dbd0	4725be249ad19e6b f47e8353f8047455	caf81a425d800537 2deecc6b39d64d78

$t = 20$ :	8dc5ae65569d3855 4bb9e66d1145bfdc	9a3fb4d38ab6cf06 f14998dd5f70767e	3c4b4104168e3edb 29695fd88d81dbd0	4725be249ad19e6b f47e8353f8047455
$t = 21$ :	da34d6673d452dcf 8e30ff09ad488753	8dc5ae65569d3855 4bb9e66d1145bfdc	9a3fb4d38ab6cf06 f14998dd5f70767e	3c4b4104168e3edb 29695fd88d81dbd0
$t = 22$ :	3e2644567b709a78 0ac2b11da8f571c6	da34d6673d452dcf 8e30ff09ad488753	8dc5ae65569d3855 4bb9e66d1145bfdc	9a3fb4d38ab6cf06 f14998dd5f70767e
$t = 23$ :	4f6877b58fe55484 c66005f87db55233	3e2644567b709a78 0ac2b11da8f571c6	da34d6673d452dcf 8e30ff09ad488753	8dc5ae65569d3855 4bb9e66d1145bfdc
$t = 24$ :	9aff71163fa3a940 d3ecf13769180e6f	4f6877b58fe55484 c66005f87db55233	3e2644567b709a78 0ac2b11da8f571c6	da34d6673d452dcf 8e30ff09ad488753
$t = 25$ :	0bc5f791f8e6816b 6ddf1fd7edcce336	9aff71163fa3a940 d3ecf13769180e6f	4f6877b58fe55484 c66005f87db55233	3e2644567b709a78 0ac2b11da8f571c6
$t = 26$ :	884c3bc27bc4f941 e6e48c9a8e948365	0bc5f791f8e6816b 6ddf1fd7edcce336	9aff71163fa3a940 d3ecf13769180e6f	4f6877b58fe55484 c66005f87db55233
$t = 27$ :	eab4a9e5771b8d09 09068a4e255a0dac	884c3bc27bc4f941 e6e48c9a8e948365	0bc5f791f8e6816b 6ddf1fd7edcce336	9aff71163fa3a940 d3ecf13769180e6f
$t = 28$ :	e62349090f47d30a 0fcd99710f21584	eab4a9e5771b8d09 09068a4e255a0dac	884c3bc27bc4f941 e6e48c9a8e948365	0bc5f791f8e6816b 6ddf1fd7edcce336
$t = 29$ :	74bf40f869094c63 f0aec2fe1437f085	e62349090f47d30a 0fcd99710f21584	eab4a9e5771b8d09 09068a4e255a0dac	884c3bc27bc4f941 e6e48c9a8e948365
$t = 30$ :	4c4fbbb75f1873a6 73e025d91b9efea3	74bf40f869094c63 f0aec2fe1437f085	e62349090f47d30a 0fcd99710f21584	eab4a9e5771b8d09 09068a4e255a0dac
$t = 31$ :	ff4d3f1f0d46a736 3cd388e119e8162e	4c4fbbb75f1873a6 73e025d91b9efea3	74bf40f869094c63 f0aec2fe1437f085	e62349090f47d30a 0fcd99710f21584
$t = 32$ :	a0509015ca08c8d4 e1034573654a106f	ff4d3f1f0d46a736 3cd388e119e8162e	4c4fbbb75f1873a6 73e025d91b9efea3	74bf40f869094c63 f0aec2fe1437f085
$t = 33$ :	60d4e6995ed91fe6 efabbd8bf47c041a	a0509015ca08c8d4 e1034573654a106f	ff4d3f1f0d46a736 3cd388e119e8162e	4c4fbbb75f1873a6 73e025d91b9efea3
$t = 34$ :	2c59ec7743632621 0fbae670fa780fd3	60d4e6995ed91fe6 efabbd8bf47c041a	a0509015ca08c8d4 e1034573654a106f	ff4d3f1f0d46a736 3cd388e119e8162e
$t = 35$ :	1a081afc59fdbc2c f098082f502b44cd	2c59ec7743632621 0fbae670fa780fd3	60d4e6995ed91fe6 efabbd8bf47c041a	a0509015ca08c8d4 e1034573654a106f
$t = 36$ :	88df85b0bbe77514 8fbfd0162bbf4675	1a081afc59fdbc2c f098082f502b44cd	2c59ec7743632621 0fbae670fa780fd3	60d4e6995ed91fe6 efabbd8bf47c041a
$t = 37$ :	002bb8e4cd989567 66adcfa249ac7bbd	88df85b0bbe77514 8fbfd0162bbf4675	1a081afc59fdbc2c f098082f502b44cd	2c59ec7743632621 0fbae670fa780fd3
$t = 38$ :	b3bb8542b3376de5 b49596c20feba7de	002bb8e4cd989567 66adcfa249ac7bbd	88df85b0bbe77514 8fbfd0162bbf4675	1a081afc59fdbc2c f098082f502b44cd
$t = 39$ :	8e01e125b855d225 0c710a47ba6a567b	b3bb8542b3376de5 b49596c20feba7de	002bb8e4cd989567 66adcfa249ac7bbd	88df85b0bbe77514 8fbfd0162bbf4675
$t = 40$ :	b01521dd6a6be12c 169008b3a4bb170b	8e01e125b855d225 0c710a47ba6a567b	b3bb8542b3376de5 b49596c20feba7de	002bb8e4cd989567 66adcfa249ac7bbd
$t = 41$ :	e96f89dd48cbd851 f0996439e7b50cb1	b01521dd6a6be12c 169008b3a4bb170b	8e01e125b855d225 0c710a47ba6a567b	b3bb8542b3376de5 b49596c20feba7de
$t = 42$ :	bc05ba8de5d3c480 639cb938e14dc190	e96f89dd48cbd851 f0996439e7b50cb1	b01521dd6a6be12c 169008b3a4bb170b	8e01e125b855d225 0c710a47ba6a567b
$t = 43$ :	35d7e7f41defcbd5	bc05ba8de5d3c480	e96f89dd48cbd851	b01521dd6a6be12c

	cc5100997f5710f2	639cb938e14dc190	f0996439e7b50cb1	169008b3a4bb170b
$t = 44 :$	c47c9d5c7ea8a234 858d832ae0e8911c	35d7e7f41defcbd5 cc5100997f5710f2	bc05ba8de5d3c480 639cb938e14dc190	e96f89dd48cbd851 f0996439e7b50cb1
$t = 45 :$	021fbadbabab5ac6 e95c2a57572d64d9	c47c9d5c7ea8a234 858d832ae0e8911c	35d7e7f41defcbd5 cc5100997f5710f2	bc05ba8de5d3c480 639cb938e14dc190
$t = 46 :$	f61e672694de2d67 c6bc35740d8daa9a	021fbadbabab5ac6 e95c2a57572d64d9	c47c9d5c7ea8a234 858d832ae0e8911c	35d7e7f41defcbd5 cc5100997f5710f2
$t = 47 :$	6b69fc1bb482feac 35264334c03ac8ad	f61e672694de2d67 c6bc35740d8daa9a	021fbadbabab5ac6 e95c2a57572d64d9	c47c9d5c7ea8a234 858d832ae0e8911c
$t = 48 :$	571f323d96b3a047 271580ed6c3e5650	6b69fc1bb482feac 35264334c03ac8ad	f61e672694de2d67 c6bc35740d8daa9a	021fbadbabab5ac6 e95c2a57572d64d9
$t = 49 :$	ca9bd862c5050918 dfe091dab182e645	571f323d96b3a047 271580ed6c3e5650	6b69fc1bb482feac 35264334c03ac8ad	f61e672694de2d67 c6bc35740d8daa9a
$t = 50 :$	813a43dd2c502043 07a0d8ef821c5e1a	ca9bd862c5050918 dfe091dab182e645	571f323d96b3a047 271580ed6c3e5650	6b69fc1bb482feac 35264334c03ac8ad
$t = 51 :$	d43f83727325dd77 483f80a82eaae23e	813a43dd2c502043 07a0d8ef821c5e1a	ca9bd862c5050918 dfe091dab182e645	571f323d96b3a047 271580ed6c3e5650
$t = 52 :$	03df11b32d42e203 504f94e40591cffa	d43f83727325dd77 483f80a82eaae23e	813a43dd2c502043 07a0d8ef821c5e1a	ca9bd862c5050918 dfe091dab182e645
$t = 53 :$	d63f68037ddf06aa a6781efelaa1ce02	03df11b32d42e203 504f94e40591cffa	d43f83727325dd77 483f80a82eaae23e	813a43dd2c502043 07a0d8ef821c5e1a
$t = 54 :$	f650857b5babda4d 9ccfb31a86df0f86	d63f68037ddf06aa a6781efelaa1ce02	03df11b32d42e203 504f94e40591cffa	d43f83727325dd77 483f80a82eaae23e
$t = 55 :$	63b460e42748817e c6b4dd2a9931c509	f650857b5babda4d 9ccfb31a86df0f86	d63f68037ddf06aa a6781efelaa1ce02	03df11b32d42e203 504f94e40591cffa
$t = 56 :$	7a52912943d52b05 d2e89bbd91e00be0	63b460e42748817e c6b4dd2a9931c509	f650857b5babda4d 9ccfb31a86df0f86	d63f68037ddf06aa a6781efelaa1ce02
$t = 57 :$	4b81c3aec976ea4b 70505988124351ac	7a52912943d52b05 d2e89bbd91e00be0	63b460e42748817e c6b4dd2a9931c509	f650857b5babda4d 9ccfb31a86df0f86
$t = 58 :$	581ecb3355dcd9b8 6a3c9b0f71c8bf36	4b81c3aec976ea4b 70505988124351ac	7a52912943d52b05 d2e89bbd91e00be0	63b460e42748817e c6b4dd2a9931c509
$t = 59 :$	2c074484ef1eac8c 4797cde4ed370692	581ecb3355dcd9b8 6a3c9b0f71c8bf36	4b81c3aec976ea4b 70505988124351ac	7a52912943d52b05 d2e89bbd91e00be0
$t = 60 :$	3857dfd2fc37d3ba a6af4e9c9f807e51	2c074484ef1eac8c 4797cde4ed370692	581ecb3355dcd9b8 6a3c9b0f71c8bf36	4b81c3aec976ea4b 70505988124351ac
$t = 61 :$	cfcd928c5424e2b6 09aee5bda1644de5	3857dfd2fc37d3ba a6af4e9c9f807e51	2c074484ef1eac8c 4797cde4ed370692	581ecb3355dcd9b8 6a3c9b0f71c8bf36
$t = 62 :$	a81dedbb9f19e643 84058865d60a05fa	cfcd928c5424e2b6 09aee5bda1644de5	3857dfd2fc37d3ba a6af4e9c9f807e51	2c074484ef1eac8c 4797cde4ed370692
$t = 63 :$	ab44e86276478d85 cd881ee59ca6bc53	a81dedbb9f19e643 84058865d60a05fa	cfcd928c5424e2b6 09aee5bda1644de5	3857dfd2fc37d3ba a6af4e9c9f807e51
$t = 64 :$	5a806d7e9821a501 aa84b086688a5c45	ab44e86276478d85 cd881ee59ca6bc53	a81dedbb9f19e643 84058865d60a05fa	cfcd928c5424e2b6 09aee5bda1644de5
$t = 65 :$	eeb9c21bb0102598 3b5fed0d6a1f96e1	5a806d7e9821a501 aa84b086688a5c45	ab44e86276478d85 cd881ee59ca6bc53	a81dedbb9f19e643 84058865d60a05fa
$t = 66 :$	46c4210ab2cc155d 29fab5a7bff53366	eeb9c21bb0102598 3b5fed0d6a1f96e1	5a806d7e9821a501 aa84b086688a5c45	ab44e86276478d85 cd881ee59ca6bc53



$t = 67$ :	54ba35cf56a0340e 1c66f46d95690bcf	46c4210ab2cc155d 29fab5a7bff53366	eeb9c21bb0102598 3b5fed0d6a1f96e1	5a806d7e9821a501 aa84b086688a5c45
$t = 68$ :	181839d609c79748 0ada78ba2d446140	54ba35cf56a0340e 1c66f46d95690bcf	46c4210ab2cc155d 29fab5a7bff53366	eeb9c21bb0102598 3b5fed0d6a1f96e1
$t = 69$ :	fb6aaaae5d0b6a447 e3711cb6564d112d	181839d609c79748 0ada78ba2d446140	54ba35cf56a0340e 1c66f46d95690bcf	46c4210ab2cc155d 29fab5a7bff53366
$t = 70$ :	7652c579cb60f19c aff62c9665ff80fa	fb6aaaae5d0b6a447 e3711cb6564d112d	181839d609c79748 0ada78ba2d446140	54ba35cf56a0340e 1c66f46d95690bcf
$t = 71$ :	f15e9664b2803575 947c3dfafee570ef	7652c579cb60f19c aff62c9665ff80fa	fb6aaaae5d0b6a447 e3711cb6564d112d	181839d609c79748 0ada78ba2d446140
$t = 72$ :	358406d165aee9ab 8c7b5fd91a794ca0	f15e9664b2803575 947c3dfafee570ef	7652c579cb60f19c aff62c9665ff80fa	fb6aaaae5d0b6a447 e3711cb6564d112d
$t = 73$ :	20878dcd29cdfaf5 054d3536539948d0	358406d165aee9ab 8c7b5fd91a794ca0	f15e9664b2803575 947c3dfafee570ef	7652c579cb60f19c aff62c9665ff80fa
$t = 74$ :	33d48dabb5521de2 2ba18245b50de4cf	20878dcd29cdfaf5 054d3536539948d0	358406d165aee9ab 8c7b5fd91a794ca0	f15e9664b2803575 947c3dfafee570ef
$t = 75$ :	c8960e6be864b916 995019a6ff3ba3de	33d48dabb5521de2 2ba18245b50de4cf	20878dcd29cdfaf5 054d3536539948d0	358406d165aee9ab 8c7b5fd91a794ca0
$t = 76$ :	654ef9abec389ca9 ceb9fc3691ce8326	c8960e6be864b916 995019a6ff3ba3de	33d48dabb5521de2 2ba18245b50de4cf	20878dcd29cdfaf5 054d3536539948d0
$t = 77$ :	d67806db8b148677 25c96a7768fb2aa3	654ef9abec389ca9 ceb9fc3691ce8326	c8960e6be864b916 995019a6ff3ba3de	33d48dabb5521de2 2ba18245b50de4cf
$t = 78$ :	10d9c4c4295599f6 9bb4d39778c07f9e	d67806db8b148677 25c96a7768fb2aa3	654ef9abec389ca9 ceb9fc3691ce8326	c8960e6be864b916 995019a6ff3ba3de
$t = 79$ :	73a54f399fa4b1b2 d08446aa79693ed7	10d9c4c4295599f6 9bb4d39778c07f9e	d67806db8b148677 25c96a7768fb2aa3	654ef9abec389ca9 ceb9fc3691ce8326

That completes the processing of the first and only message block,  $M^{(1)}$ . The final hash value,  $H^{(1)}$ , is calculated to be

$$\begin{aligned}
 H_0^{(1)} &= 6a09e667f3bcc908 + 73a54f399fa4b1b2 = \text{ddaf35a193617aba} \\
 H_1^{(1)} &= \text{bb67ae8584caa73b} + 10d9c4c4295599f6 = \text{cc417349ae204131} \\
 H_2^{(1)} &= 3c6ef372fe94f82b + d67806db8b148677 = \text{12e6fa4e89a97ea2} \\
 H_3^{(1)} &= a54ff53a5f1d36f1 + 654ef9abec389ca9 = \text{0a9eeee64b55d39a} \\
 H_4^{(1)} &= 510e527fade682d1 + d08446aa79693ed7 = \text{2192992a274fcl a8} \\
 H_5^{(1)} &= 9b05688c2b3e6c1f + 9bb4d39778c07f9e = \text{36ba3c23a3feebbd} \\
 H_6^{(1)} &= 1f83d9abfb41bd6b + 25c96a7768fb2aa3 = \text{454d4423643ce80e} \\
 H_7^{(1)} &= 5be0cd19137e2179 + ceb9fc3691ce8326 = \text{2a9ac94fa54ca49f} .
 \end{aligned}$$

The resulting 512-bit message digest is

```
ddaf35a193617aba cc417349ae204131 12e6fa4e89a97ea2 0a9eeee64b55d39a
2192992a274fcl a8 36ba3c23a3feebbd 454d4423643ce80e 2a9ac94fa54ca49f .
```

## C.2 SHA-512 Example (Multi-Block Message)

Let the message,  $M$ , be the 896-bit ( $\ell = 896$ ) ASCII string

**"abcdefghijklmnopghicdefghijdefghijkefghijklfghijklmghijklmn  
hijklmnoijklmnopjklmnopqklmnopqrlmnopqrsmnopqrstnoprstu".**

The message is padded by appending a "1" bit, followed by 1023 "0" bits, and ending with the hex value

000000000000000000 00000000000000380

(the two 64-bit word representation of the length, 896). Thus, the final padded message consists of two blocks ( $N = 2$ ).

For SHA-512, the initial hash value,  $H^{(0)}$ , is

$H_0^{(0)} = 6a09e667f3bcc908$   
 $H_1^{(0)} = bb67ae8584caa73b$   
 $H_2^{(0)} = 3c6ef372fe94f82b$   
 $H_3^{(0)} = a54ff53a5f1d36f1$   
 $H_4^{(0)} = 510e527fade682d1$   
 $H_5^{(0)} = 9b05688c2b3e6c1f$   
 $H_6^{(0)} = 1f83d9abfb41bd6b$   
 $H_7^{(0)} = 5be0cd19137e2179.$

The words of the padded message block are then assigned to the words  $W_0, \dots, W_{15}$  of the message schedule:

$W_0 = 6162636465666768$	$W_8 = 696a6b6c6d6e6f70$
$W_1 = 6263646566676869$	$W_9 = 6a6b6c6d6e6f7071$
$W_2 = 636465666768696a$	$W_{10} = 6b6c6d6e6f707172$
$W_3 = 6465666768696a6b$	$W_{11} = 6c6d6e6f70717273$
$W_4 = 65666768696a6b6c$	$W_{12} = 6d6e6f7071727374$
$W_5 = 666768696a6b6c6d$	$W_{13} = 6e6f707172737475$
$W_6 = 6768696a6b6c6d6e$	$W_{14} = 8000000000000000$
$W_7 = 68696a6b6c6d6e6f$	$W_{15} = 0000000000000000.$

The following schedule shows the hex values for  $a$ ,  $b$ ,  $c$ ,  $d$ ,  $e$ ,  $f$ ,  $g$ , and  $h$  after pass  $t$  of the "for  $t = 0$  to 79" loop described in Sec. 6.3.2, step 4.

	<i>a</i> / <i>e</i>	<i>b</i> / <i>f</i>	<i>c</i> / <i>g</i>	<i>d</i> / <i>h</i>
<i>t</i> = 0 :	f6afce9d2263455d 58cb0218e01b86f9	6a09e667f3bcc908 510e527fade682d1	bb67ae8584caa73b 9b05688c2b3e6c1f	3c6ef372fe94f82b 1f83d9abfb41bd6b
<i>t</i> = 1 :	0b7056a534ae5f62 f8c7198fe39e4c8c	f6afce9d2263455d 58cb0218e01b86f9	6a09e667f3bcc908 510e527fade682d1	bb67ae8584caa73b 9b05688c2b3e6c1f
<i>t</i> = 2 :	2ca82233760c9942 303eccccd65953de	0b7056a534ae5f62 f8c7198fe39e4c8c	f6afce9d2263455d 58cb0218e01b86f9	6a09e667f3bcc908 510e527fade682d1
<i>t</i> = 3 :	a023f17ce52cda7b ffdee5eedcc9ca42	2ca82233760c9942 303eccccd65953de	0b7056a534ae5f62 f8c7198fe39e4c8c	f6afce9d2263455d 58cb0218e01b86f9
<i>t</i> = 4 :	8f0a67d9d591a1a7 cb4cfbb166505f2f	a023f17ce52cda7b ffdee5eedcc9ca42	2ca82233760c9942 303eccccd65953de	0b7056a534ae5f62 f8c7198fe39e4c8c
<i>t</i> = 5 :	b466267371acc493 73d6c84c54d399ee	8f0a67d9d591a1a7 cb4cfbb166505f2f	a023f17ce52cda7b ffdee5eedcc9ca42	2ca82233760c9942 303eccccd65953de
<i>t</i> = 6 :	658269f1a312fccd cdc40314975fb275	b466267371acc493 73d6c84c54d399ee	8f0a67d9d591a1a7 cb4cfbb166505f2f	a023f17ce52cda7b ffdee5eedcc9ca42
<i>t</i> = 7 :	65e3519c5b88181b a657850ab3970c5a	658269f1a312fccd cdc40314975fb275	b466267371acc493 73d6c84c54d399ee	8f0a67d9d591a1a7 cb4cfbb166505f2f
<i>t</i> = 8 :	56604fbb4b6393ec e8b3be22fbe64df7	65e3519c5b88181b a657850ab3970c5a	658269f1a312fccd cdc40314975fb275	b466267371acc493 73d6c84c54d399ee
<i>t</i> = 9 :	c4562769a37d02c0 0062e70a1ef705c1	56604fbb4b6393ec e8b3be22fbe64df7	65e3519c5b88181b a657850ab3970c5a	658269f1a312fccd cdc40314975fb275
<i>t</i> = 10 :	27c0b4c9186e1736 bc9740477a18ae2d	c4562769a37d02c0 0062e70a1ef705c1	56604fbb4b6393ec e8b3be22fbe64df7	65e3519c5b88181b a657850ab3970c5a
<i>t</i> = 11 :	f17f52fb02f4eb74 be58522cb9590ee1	27c0b4c9186e1736 bc9740477a18ae2d	c4562769a37d02c0 0062e70a1ef705c1	56604fbb4b6393ec e8b3be22fbe64df7
<i>t</i> = 12 :	f2c245ac903d4a35 49d5fa3a16dcd502	f17f52fb02f4eb74 be58522cb9590ee1	27c0b4c9186e1736 bc9740477a18ae2d	c4562769a37d02c0 0062e70a1ef705c1
<i>t</i> = 13 :	9b04175ea8090daa ec9c5e98ff98760d	f2c245ac903d4a35 49d5fa3a16dcd502	f17f52fb02f4eb74 be58522cb9590ee1	27c0b4c9186e1736 bc9740477a18ae2d
<i>t</i> = 14 :	481b8a6ee5e07031 e4d35b613a5ac420	9b04175ea8090daa ec9c5e98ff98760d	f2c245ac903d4a35 49d5fa3a16dcd502	f17f52fb02f4eb74 be58522cb9590ee1
<i>t</i> = 15 :	9356ac3ec3e51459 701f17d27582443b	481b8a6ee5e07031 e4d35b613a5ac420	9b04175ea8090daa ec9c5e98ff98760d	f2c245ac903d4a35 49d5fa3a16dcd502
<i>t</i> = 16 :	b889ed34abd7aa37 1d05d9ba779a1a78	9356ac3ec3e51459 701f17d27582443b	481b8a6ee5e07031 e4d35b613a5ac420	9b04175ea8090daa ec9c5e98ff98760d
<i>t</i> = 17 :	bf537b1f3edc7381 c362ff9cf932951d	b889ed34abd7aa37 1d05d9ba779a1a78	9356ac3ec3e51459 701f17d27582443b	481b8a6ee5e07031 e4d35b613a5ac420
<i>t</i> = 18 :	d4e44d54e8242ad8 459e4e6888919f36	bf537b1f3edc7381 c362ff9cf932951d	b889ed34abd7aa37 1d05d9ba779a1a78	9356ac3ec3e51459 701f17d27582443b
<i>t</i> = 19 :	05f3fba454e5de3d caed4b5fa322b984	d4e44d54e8242ad8 459e4e6888919f36	bf537b1f3edc7381 c362ff9cf932951d	b889ed34abd7aa37 1d05d9ba779a1a78
<i>t</i> = 20 :	cdb73772dc0248bf dc8049afa6acd502	05f3fba454e5de3d caed4b5fa322b984	d4e44d54e8242ad8 459e4e6888919f36	bf537b1f3edc7381 c362ff9cf932951d
<i>t</i> = 21 :	1d47a3268ff677ed	cdb73772dc0248bf	05f3fba454e5de3d	d4e44d54e8242ad8

	8407818e9b28cc12	dc8049afa6acd502	caed4b5fa322b984	459e4e6888919f36
$t = 22 :$	af4e23eb622d0df4 64b5ae5424598428	1d47a3268ff677ed 8407818e9b28cc12	cdb73772dc0248bf dc8049afa6acd502	05f3fba454e5de3d caed4b5fa322b984
$t = 23 :$	be50606778de14a6 0a5d727cc92e7adb	af4e23eb622d0df4 64b5ae5424598428	1d47a3268ff677ed 8407818e9b28cc12	cdb73772dc0248bf dc8049afa6acd502
$t = 24 :$	821e44f6678ac478 f367e596d0a038a5	be50606778de14a6 0a5d727cc92e7adb	af4e23eb622d0df4 64b5ae5424598428	1d47a3268ff677ed 8407818e9b28cc12
$t = 25 :$	0c852b1359a77c18 6dec8a3396a80c3f	821e44f6678ac478 f367e596d0a038a5	be50606778de14a6 0a5d727cc92e7adb	af4e23eb622d0df4 64b5ae5424598428
$t = 26 :$	ebb574fad4b7a7e4 a241e7efc1eb6fff9	0c852b1359a77c18 6dec8a3396a80c3f	821e44f6678ac478 f367e596d0a038a5	be50606778de14a6 0a5d727cc92e7adb
$t = 27 :$	a092821c3cdf08da c84e849917a7c08e	ebb574fad4b7a7e4 a241e7efc1eb6fff9	0c852b1359a77c18 6dec8a3396a80c3f	821e44f6678ac478 f367e596d0a038a5
$t = 28 :$	82ba2e1a2df2a4f1 61845f6924789851	a092821c3cdf08da c84e849917a7c08e	ebb574fad4b7a7e4 a241e7efc1eb6fff9	0c852b1359a77c18 6dec8a3396a80c3f
$t = 29 :$	1959ad991c63d06a 231faf24910a891a	82ba2e1a2df2a4f1 61845f6924789851	a092821c3cdf08da c84e849917a7c08e	ebb574fad4b7a7e4 a241e7efc1eb6fff9
$t = 30 :$	9b32d4cacd9a625b 533066919d608799	1959ad991c63d06a 231faf24910a891a	82ba2e1a2df2a4f1 61845f6924789851	a092821c3cdf08da c84e849917a7c08e
$t = 31 :$	dc55339f4d841965 e2517f359998a58d	9b32d4cacd9a625b 533066919d608799	1959ad991c63d06a 231faf24910a891a	82ba2e1a2df2a4f1 61845f6924789851
$t = 32 :$	fdebb1283b12514f b1989170a183c661	dc55339f4d841965 e2517f359998a58d	9b32d4cacd9a625b 533066919d608799	1959ad991c63d06a 231faf24910a891a
$t = 33 :$	b44c7975a83e3334 009ad175b8d588a4	fdebb1283b12514f b1989170a183c661	dc55339f4d841965 e2517f359998a58d	9b32d4cacd9a625b 533066919d608799
$t = 34 :$	0bac61bfc53d18b7 a7d5416d690557b8	b44c7975a83e3334 009ad175b8d588a4	fdebb1283b12514f b1989170a183c661	dc55339f4d841965 e2517f359998a58d
$t = 35 :$	392893c22e75856a 7a7c9eb7bc813248	0bac61bfc53d18b7 a7d5416d690557b8	b44c7975a83e3334 009ad175b8d588a4	fdebb1283b12514f b1989170a183c661
$t = 36 :$	824408631432e09b 5e696a9fda56d6bf	392893c22e75856a 7a7c9eb7bc813248	0bac61bfc53d18b7 a7d5416d690557b8	b44c7975a83e3334 009ad175b8d588a4
$t = 37 :$	a64162f151a8c1cb 0f57062401dc680b	824408631432e09b 5e696a9fda56d6bf	392893c22e75856a 7a7c9eb7bc813248	0bac61bfc53d18b7 a7d5416d690557b8
$t = 38 :$	922537abad1e95a1 4f4c193d435ff721	a64162f151a8c1cb 0f57062401dc680b	824408631432e09b 5e696a9fda56d6bf	392893c22e75856a 7a7c9eb7bc813248
$t = 39 :$	b80591f6fbfadcde 00f4407c0f37237e	922537abad1e95a1 4f4c193d435ff721	a64162f151a8c1cb 0f57062401dc680b	824408631432e09b 5e696a9fda56d6bf
$t = 40 :$	08f151f4b8d0fa2e ec8b96fe402094cd	b80591f6fbfadcde 00f4407c0f37237e	922537abad1e95a1 4f4c193d435ff721	a64162f151a8c1cb 0f57062401dc680b
$t = 41 :$	12b5fcc2b68f65c0 d688101dfd24a148	08f151f4b8d0fa2e ec8b96fe402094cd	b80591f6fbfadcde 00f4407c0f37237e	922537abad1e95a1 4f4c193d435ff721
$t = 42 :$	a71bf5bd64289948 e052bfb7a6945939	12b5fcc2b68f65c0 d688101dfd24a148	08f151f4b8d0fa2e ec8b96fe402094cd	b80591f6fbfadcde 00f4407c0f37237e
$t = 43 :$	890c2cd670c4aea3 dd13e4edeef00e7	a71bf5bd64289948 e052bfb7a6945939	12b5fcc2b68f65c0 d688101dfd24a148	08f151f4b8d0fa2e ec8b96fe402094cd
$t = 44 :$	ca61990b43297ffc	890c2cd670c4aea3	a71bf5bd64289948	12b5fcc2b68f65c0

	139aa55c51d9ee5f	dd13e4edeeff00e7	e052bfb7a6945939	d688101dfd24a148
$t = 45 :$	7196e8fa538ba4bf 046735513cdd14d3	ca61990b43297ffc 139aa55c51d9ee5f	890c2cd670c4aea3 dd13e4edeeff00e7	a71bf5bd64289948 e052bfb7a6945939
$t = 46 :$	1f0720944dbeb6a4 a41eb7e5a27588e3	7196e8fa538ba4bf 046735513cdd14d3	ca61990b43297ffc 139aa55c51d9ee5f	890c2cd670c4aea3 dd13e4edeeff00e7
$t = 47 :$	d6d4f8608b8ab199 24b9c216f915da60	1f0720944dbeb6a4 a41eb7e5a27588e3	7196e8fa538ba4bf 046735513cdd14d3	ca61990b43297ffc 139aa55c51d9ee5f
$t = 48 :$	88761eb67845978e 9fe22e39448d50ed	d6d4f8608b8ab199 24b9c216f915da60	1f0720944dbeb6a4 a41eb7e5a27588e3	7196e8fa538ba4bf 046735513cdd14d3
$t = 49 :$	7d40e6be47d85702 d9c900e01968c33e	88761eb67845978e 9fe22e39448d50ed	d6d4f8608b8ab199 24b9c216f915da60	1f0720944dbeb6a4 a41eb7e5a27588e3
$t = 50 :$	7d0d988df5768598 2ec2e522a7c7d12c	7d40e6be47d85702 d9c900e01968c33e	88761eb67845978e 9fe22e39448d50ed	d6d4f8608b8ab199 24b9c216f915da60
$t = 51 :$	48a8b60575b37f31 7059f9bc8c88a373	7d0d988df5768598 2ec2e522a7c7d12c	7d40e6be47d85702 d9c900e01968c33e	88761eb67845978e 9fe22e39448d50ed
$t = 52 :$	6bc425af294bbf79 6a8143b1716ee33d	48a8b60575b37f31 7059f9bc8c88a373	7d0d988df5768598 2ec2e522a7c7d12c	7d40e6be47d85702 d9c900e01968c33e
$t = 53 :$	307a456158ee8849 4372e85c16ee4440	6bc425af294bbf79 6a8143b1716ee33d	48a8b60575b37f31 7059f9bc8c88a373	7d0d988df5768598 2ec2e522a7c7d12c
$t = 54 :$	af36382c8fd716be a8f8b0033187a916	307a456158ee8849 4372e85c16ee4440	6bc425af294bbf79 6a8143b1716ee33d	48a8b60575b37f31 7059f9bc8c88a373
$t = 55 :$	810eb951c64ca1 16a64f5997b9cca6	af36382c8fd716be a8f8b0033187a916	307a456158ee8849 4372e85c16ee4440	6bc425af294bbf79 6a8143b1716ee33d
$t = 56 :$	2dd7659f1b4d13cd 5da6793bb7286a4b	810eb951c64ca1 16a64f5997b9cca6	af36382c8fd716be a8f8b0033187a916	307a456158ee8849 4372e85c16ee4440
$t = 57 :$	5ac712acff4b98be 91f6395b301adbfcd	2dd7659f1b4d13cd 5da6793bb7286a4b	810eb951c64ca1 16a64f5997b9cca6	af36382c8fd716be a8f8b0033187a916
$t = 58 :$	c1af358833cb03c0 d4883c0c21dda190	5ac712acff4b98be 91f6395b301adbfcd	2dd7659f1b4d13cd 5da6793bb7286a4b	810eb951c64ca1 16a64f5997b9cca6
$t = 59 :$	88a306074d388c7d 9fc52468b897f9c8	c1af358833cb03c0 d4883c0c21dda190	5ac712acff4b98be 91f6395b301adbfcd	2dd7659f1b4d13cd 5da6793bb7286a4b
$t = 60 :$	f11bfd0cf67d3040 47efb6407f74d318	88a306074d388c7d 9fc52468b897f9c8	c1af358833cb03c0 d4883c0c21dda190	5ac712acff4b98be 91f6395b301adbfcd
$t = 61 :$	1f065e7828ed4e1b 7481899904a4ce23	f11bfd0cf67d3040 47efb6407f74d318	88a306074d388c7d 9fc52468b897f9c8	c1af358833cb03c0 d4883c0c21dda190
$t = 62 :$	aebde39f2bc42ec1 62ab526ff177a988	1f065e7828ed4e1b 7481899904a4ce23	f11bfd0cf67d3040 47efb6407f74d318	88a306074d388c7d 9fc52468b897f9c8
$t = 63 :$	d35a94706e3e5df2 53f92b648d5d815c	aebde39f2bc42ec1 62ab526ff177a988	1f065e7828ed4e1b 7481899904a4ce23	f11bfd0cf67d3040 47efb6407f74d318
$t = 64 :$	d72d727c53e09ab9 10746426ba9824f4	d35a94706e3e5df2 53f92b648d5d815c	aebde39f2bc42ec1 62ab526ff177a988	1f065e7828ed4e1b 7481899904a4ce23
$t = 65 :$	3a7235e5a4051d94 afe455daec5c2b00	d72d727c53e09ab9 10746426ba9824f4	d35a94706e3e5df2 53f92b648d5d815c	aebde39f2bc42ec1 62ab526ff177a988
$t = 66 :$	f7f510fe73ef7e76 f1202c0bb7c4583f	3a7235e5a4051d94 afe455daec5c2b00	d72d727c53e09ab9 10746426ba9824f4	d35a94706e3e5df2 53f92b648d5d815c
$t = 67 :$	23c2acfb393523e9 a0bc2a61044ac12e	f7f510fe73ef7e76 f1202c0bb7c4583f	3a7235e5a4051d94 afe455daec5c2b00	d72d727c53e09ab9 10746426ba9824f4

$t = 68$ :	0307d241aled7121 fad5f38f1e0aea12	23c2acfb393523e9 a0bc2a61044ac12e	f7f510fe73ef7e76 f1202c0bb7c4583f	3a7235e5a4051d94 afe455daec5c2b00
$t = 69$ :	191814d82f0a16fb 39d325086e66e200	0307d241aled7121 fad5f38f1e0aea12	23c2acfb393523e9 a0bc2a61044ac12e	f7f510fe73ef7e76 f1202c0bb7c4583f
$t = 70$ :	0aled41b6da18c01 b3d3521e166e5df1	191814d82f0a16fb 39d325086e66e200	0307d241aled7121 fad5f38f1e0aea12	23c2acfb393523e9 a0bc2a61044ac12e
$t = 71$ :	8a3f07db93f6c827 6b370074be040ed7	0aled41b6da18c01 b3d3521e166e5df1	191814d82f0a16fb 39d325086e66e200	0307d241aled7121 fad5f38f1e0aea12
$t = 72$ :	002744d87ef80d28 8c5a245de2d72fe6	8a3f07db93f6c827 6b370074be040ed7	0aled41b6da18c01 b3d3521e166e5df1	191814d82f0a16fb 39d325086e66e200
$t = 73$ :	778dc7880a4a2aa0 45a375b466e5e342	002744d87ef80d28 8c5a245de2d72fe6	8a3f07db93f6c827 6b370074be040ed7	0aled41b6da18c01 b3d3521e166e5df1
$t = 74$ :	a3f11de5ede05b11 f5bbf52f1ab7cc05	778dc7880a4a2aa0 45a375b466e5e342	002744d87ef80d28 8c5a245de2d72fe6	8a3f07db93f6c827 6b370074be040ed7
$t = 75$ :	629c8ae6ecd8af4b 5a8fe5919d3cf136	a3f11de5ede05b11 f5bbf52f1ab7cc05	778dc7880a4a2aa0 45a375b466e5e342	002744d87ef80d28 8c5a245de2d72fe6
$t = 76$ :	c9a8c1e2d063ce94 aacd089bfae8faf9	629c8ae6ecd8af4b 5a8fe5919d3cf136	a3f11de5ede05b11 f5bbf52f1ab7cc05	778dc7880a4a2aa0 45a375b466e5e342
$t = 77$ :	c517cba6a09bb26a e1682bd33c8f8e23	c9a8c1e2d063ce94 aacd089bfae8faf9	629c8ae6ecd8af4b 5a8fe5919d3cf136	a3f11de5ede05b11 f5bbf52f1ab7cc05
$t = 78$ :	11e3570e06e3b74e 075aabbade34fd01	c517cba6a09bb26a e1682bd33c8f8e23	c9a8c1e2d063ce94 aacd089bfae8faf9	629c8ae6ecd8af4b 5a8fe5919d3cf136
$t = 79$ :	d90f1b1237b3a561 867983f69d3a3ad1	11e3570e06e3b74e 075aabbade34fd01	c517cba6a09bb26a e1682bd33c8f8e23	c9a8c1e2d063ce94 aacd089bfae8faf9

That completes the processing of the first message block,  $M^{(1)}$ . The intermediate hash value,  $H^{(1)}$ , is calculated to be

$$\begin{aligned}
 H_0^{(1)} &= 6a09e667f3bcc908 + d90f1b1237b3a561 = 4319017a2b706e69 \\
 H_1^{(1)} &= bb67ae8584caa73b + 11e3570e06e3b74e = cd4b05938bae5e89 \\
 H_2^{(1)} &= 3c6ef372fe94f82b + c517cba6a09bb26a = 0186bf199f30aa95 \\
 H_3^{(1)} &= a54ff53a5f1d36f1 + c9a8c1e2d063ce94 = 6ef8b71d2f810585 \\
 H_4^{(1)} &= 510e527fade682d1 + 867983f69d3a3ad1 = d787d6764b20bda2 \\
 H_5^{(1)} &= 9b05688c2b3e6c1f + 075aabbade34fd01 = a260144709736920 \\
 H_6^{(1)} &= 1f83d9abfb41bd6b + e1682bd33c8f8e23 = 00ec057f37d14b8e \\
 H_7^{(1)} &= 5be0cd19137e2179 + aacd089bfae8faf9 = 06add5b50e671c72.
 \end{aligned}$$

The words of the *second* padded message block,  $M^{(2)}$ , are then assigned to the words  $W_0, \dots, W_{15}$  of the message schedule:

$W_0$	=	0000000000000000	$W_8$	=	0000000000000000
$W_1$	=	0000000000000000	$W_9$	=	0000000000000000
$W_2$	=	0000000000000000	$W_{10}$	=	0000000000000000
$W_3$	=	0000000000000000	$W_{11}$	=	0000000000000000
$W_4$	=	0000000000000000	$W_{12}$	=	0000000000000000
$W_5$	=	0000000000000000	$W_{13}$	=	0000000000000000
$W_6$	=	0000000000000000	$W_{14}$	=	0000000000000000
$W_7$	=	0000000000000000	$W_{15}$	=	0000000000000380.

The following schedule shows the hex values for *a*, *b*, *c*, *d*, *e*, *f*, *g*, and *h* after pass *t* of the “for *t* = 0 to 79” loop described in Sec. 6.1.2, step 4.

	<i>a</i> / <i>e</i>	<i>b</i> / <i>f</i>	<i>c</i> / <i>g</i>	<i>d</i> / <i>h</i>
<i>t</i> = 0 :	b8fdb92bdfb187e8 1d5f4d5ad031b8e6	4319017a2b706e99 d787d6764b20bda2	cd4b05938bae5e89 a260144709736920	0186bf199f30aa95 00ec057f37d14b8e
<i>t</i> = 1 :	6eb90718369c5cd7 4b9b4877d987b0fe	b8fdb92bdfb187e8 1d5f4d5ad031b8e6	4319017a2b706e99 d787d6764b20bda2	cd4b05938bae5e89 a260144709736920
<i>t</i> = 2 :	c83451f2335d5144 d6b67350e0781e99	6eb90718369c5cd7 4b9b4877d987b0fe	b8fdb92bdfb187e8 1d5f4d5ad031b8e6	4319017a2b706e99 d787d6764b20bda2
<i>t</i> = 3 :	28ec1deb2a9ee6e3 25e3136be5999b8c	c83451f2335d5144 d6b67350e0781e99	6eb90718369c5cd7 4b9b4877d987b0fe	b8fdb92bdfb187e8 1d5f4d5ad031b8e6
<i>t</i> = 4 :	806abd86c0479e5b 1b8f7670eab1cf89	28ec1deb2a9ee6e3 25e3136be5999b8c	c83451f2335d5144 d6b67350e0781e99	6eb90718369c5cd7 4b9b4877d987b0fe
<i>t</i> = 5 :	234788f8a54aed38 4fabe51c67d5d156	806abd86c0479e5b 1b8f7670eab1cf89	28ec1deb2a9ee6e3 25e3136be5999b8c	c83451f2335d5144 d6b67350e0781e99
<i>t</i> = 6 :	01264f18257b5e2c 1c3506096b99de50	234788f8a54aed38 4fabe51c67d5d156	806abd86c0479e5b 1b8f7670eab1cf89	28ec1deb2a9ee6e3 25e3136be5999b8c
<i>t</i> = 7 :	5b14f38104dde991 13f8bfdc4001c362	01264f18257b5e2c 1c3506096b99de50	234788f8a54aed38 4fabe51c67d5d156	806abd86c0479e5b 1b8f7670eab1cf89
<i>t</i> = 8 :	f522574a41b2aac6 63a5f09617622ed2	5b14f38104dde991 13f8bfdc4001c362	01264f18257b5e2c 1c3506096b99de50	234788f8a54aed38 4fabe51c67d5d156
<i>t</i> = 9 :	6ec258b855afae5a 211e271d92770b36	f522574a41b2aac6 63a5f09617622ed2	5b14f38104dde991 13f8bfdc4001c362	01264f18257b5e2c 1c3506096b99de50
<i>t</i> = 10 :	9364214ba48b416c d64dcb6ec0fe5bac	6ec258b855afae5a 211e271d92770b36	f522574a41b2aac6 63a5f09617622ed2	5b14f38104dde991 13f8bfdc4001c362
<i>t</i> = 11 :	082ba62147ecbbd5 34fe78473b61266e	9364214ba48b416c d64dcb6ec0fe5bac	6ec258b855afae5a 211e271d92770b36	f522574a41b2aac6 63a5f09617622ed2
<i>t</i> = 12 :	5790f6ba82bba809 d491e309141dcaa3	082ba62147ecbbd5 34fe78473b61266e	9364214ba48b416c d64dcb6ec0fe5bac	6ec258b855afae5a 211e271d92770b36
<i>t</i> = 13 :	a6b8aefd086d33ce 044943c2992cc0f0	5790f6ba82bba809 d491e309141dcaa3	082ba62147ecbbd5 34fe78473b61266e	9364214ba48b416c d64dcb6ec0fe5bac
<i>t</i> = 14 :	bf2324a9a363abe7 0cf5f4bde5977c54	a6b8aefd086d33ce 044943c2992cc0f0	5790f6ba82bba809 d491e309141dcaa3	082ba62147ecbbd5 34fe78473b61266e
<i>t</i> = 15 :	00e8e32076a61aff	bf2324a9a363abe7	a6b8aefd086d33ce	5790f6ba82bba809



	43bf4eb269a2650c	0cf5f4bde5977c54	044943c2992cc0f0	d491e309141dcaa3
$t = 16 :$	f0376dff66fff4a7 69fa5896969e85b8	00e8e32076a61aff 43bf4eb269a2650c	bf2324a9a363abe7 0cf5f4bde5977c54	a6b8aefd086d33ce 044943c2992cc0f0
$t = 17 :$	2fad194272cda857 ddb519d663b7b6ec	f0376dff66fff4a7 69fa5896969e85b8	00e8e32076a61aff 43bf4eb269a2650c	bf2324a9a363abe7 0cf5f4bde5977c54
$t = 18 :$	9ae56936e95325ac 04ceb04676619057	2fad194272cda857 ddb519d663b7b6ec	f0376dff66fff4a7 69fa5896969e85b8	00e8e32076a61aff 43bf4eb269a2650c
$t = 19 :$	d94ccb853f53433b dc0f45813fb5a2	9ae56936e95325ac 04ceb04676619057	2fad194272cda857 ddb519d663b7b6ec	f0376dff66fff4a7 69fa5896969e85b8
$t = 20 :$	837f8075d2945995 272b5f79a91419d8	d94ccb853f53433b dc0f45813fb5a2	9ae56936e95325ac 04ceb04676619057	2fad194272cda857 ddb519d663b7b6ec
$t = 21 :$	786bde689f7aa62d 566586e69ad3f487	837f8075d2945995 272b5f79a91419d8	d94ccb853f53433b dc0f45813fb5a2	9ae56936e95325ac 04ceb04676619057
$t = 22 :$	276457f01812aa6f e78fb8b0dfbbc62f	786bde689f7aa62d 566586e69ad3f487	837f8075d2945995 272b5f79a91419d8	d94ccb853f53433b dc0f45813fb5a2
$t = 23 :$	0de519f5d6c2c298 5ca3e5cd1a30b954	276457f01812aa6f e78fb8b0dfbbc62f	786bde689f7aa62d 566586e69ad3f487	837f8075d2945995 272b5f79a91419d8
$t = 24 :$	54314dff825e2b22 b81a51e0c96ccf77	0de519f5d6c2c298 5ca3e5cd1a30b954	276457f01812aa6f e78fb8b0dfbbc62f	786bde689f7aa62d 566586e69ad3f487
$t = 25 :$	5d3f98dd7b29c363 95d49494f5a0d14a	54314dff825e2b22 b81a51e0c96ccf77	0de519f5d6c2c298 5ca3e5cd1a30b954	276457f01812aa6f e78fb8b0dfbbc62f
$t = 26 :$	5e9da426aa7d4a58 d22cccad2e391cd4	5d3f98dd7b29c363 95d49494f5a0d14a	54314dff825e2b22 b81a51e0c96ccf77	0de519f5d6c2c298 5ca3e5cd1a30b954
$t = 27 :$	3b62dd973298ea43 aceb5d06101e514e	5e9da426aa7d4a58 d22cccad2e391cd4	5d3f98dd7b29c363 95d49494f5a0d14a	54314dff825e2b22 b81a51e0c96ccf77
$t = 28 :$	fd258ff809b2253d 26c991e85352da6f	3b62dd973298ea43 aceb5d06101e514e	5e9da426aa7d4a58 d22cccad2e391cd4	5d3f98dd7b29c363 95d49494f5a0d14a
$t = 29 :$	b462a20846af417d 291eee54c034c326	fd258ff809b2253d 26c991e85352da6f	3b62dd973298ea43 aceb5d06101e514e	5e9da426aa7d4a58 d22cccad2e391cd4
$t = 30 :$	d5471e3dc7171224 0aaf99c59e7fadbd	b462a20846af417d 291eee54c034c326	fd258ff809b2253d 26c991e85352da6f	3b62dd973298ea43 aceb5d06101e514e
$t = 31 :$	9ace856ba1290e6e 658f0bea63804d05	d5471e3dc7171224 0aaf99c59e7fadbd	b462a20846af417d 291eee54c034c326	fd258ff809b2253d 26c991e85352da6f
$t = 32 :$	80a0d154506b37c4 bbe6e3b3bb7fefab	9ace856ba1290e6e 658f0bea63804d05	d5471e3dc7171224 0aaf99c59e7fadbd	b462a20846af417d 291eee54c034c326
$t = 33 :$	fb90a8a76dea1bfe 65234d5b5049e665	80a0d154506b37c4 bbe6e3b3bb7fefab	9ace856ba1290e6e 658f0bea63804d05	d5471e3dc7171224 0aaf99c59e7fadbd
$t = 34 :$	f517b690d940a294 e4dd663f44d313bc	fb90a8a76dea1bfe 65234d5b5049e665	80a0d154506b37c4 bbe6e3b3bb7fefab	9ace856ba1290e6e 658f0bea63804d05
$t = 35 :$	b70883992932880d dc5dd7c12b1cb6e3	f517b690d940a294 e4dd663f44d313bc	fb90a8a76dea1bfe 65234d5b5049e665	80a0d154506b37c4 bbe6e3b3bb7fefab
$t = 36 :$	b2a2be77b0fcf3bf 50fca57291e19874	b70883992932880d dc5dd7c12b1cb6e3	f517b690d940a294 e4dd663f44d313bc	fb90a8a76dea1bfe 65234d5b5049e665
$t = 37 :$	8575839b0f08472b bd7176bd099bb2f2	b2a2be77b0fcf3bf 50fca57291e19874	b70883992932880d dc5dd7c12b1cb6e3	f517b690d940a294 e4dd663f44d313bc
$t = 38 :$	4405d2765de0adfc	8575839b0f08472b	b2a2be77b0fcf3bf	b70883992932880d



	7ca4916f2cd8db10	bd7176bd099bb2f2	50fca57291e19874	dc5dd7c12b1cb6e3
$t = 39$ :	eec6fca5aa657661 7be0b7e70bdabe53	4405d2765de0adfc 7ca4916f2cd8db10	8575839b0f08472b bd7176bd099bb2f2	b2a2be77b0fcf3bf 50fca57291e19874
$t = 40$ :	bb3fcd7585b59e32 2201c7cbd34e31fe	eec6fca5aa657661 7be0b7e70bdabe53	4405d2765de0adfc 7ca4916f2cd8db10	8575839b0f08472b bd7176bd099bb2f2
$t = 41$ :	0e109efc47927341 d43e5686506fa05d	bb3fcd7585b59e32 2201c7cbd34e31fe	eec6fca5aa657661 7be0b7e70bdabe53	4405d2765de0adfc 7ca4916f2cd8db10
$t = 42$ :	55c0dba83bcd6e0 5b634502f1671535	0e109efc47927341 d43e5686506fa05d	bb3fcd7585b59e32 2201c7cbd34e31fe	eec6fca5aa657661 7be0b7e70bdabe53
$t = 43$ :	f5756f847bfaef67 e2d307fd94f4818a	55c0dba83bcd6e0 5b634502f1671535	0e109efc47927341 d43e5686506fa05d	bb3fcd7585b59e32 2201c7cbd34e31fe
$t = 44$ :	f1438c9cf271c06e ad8ac1ed966b2dc6	f5756f847bfaef67 e2d307fd94f4818a	55c0dba83bcd6e0 5b634502f1671535	0e109efc47927341 d43e5686506fa05d
$t = 45$ :	a7dcaffdbefb9d4a 9e46e9f915099c34	f1438c9cf271c06e ad8ac1ed966b2dc6	f5756f847bfaef67 e2d307fd94f4818a	55c0dba83bcd6e0 5b634502f1671535
$t = 46$ :	985ba373680b8e94 7d4c0abc676b1a8b	a7dcaffdbefb9d4a 9e46e9f915099c34	f1438c9cf271c06e ad8ac1ed966b2dc6	f5756f847bfaef67 e2d307fd94f4818a
$t = 47$ :	807f45784852303f 082ee70d3f352aac	985ba373680b8e94 7d4c0abc676b1a8b	a7dcaffdbefb9d4a 9e46e9f915099c34	f1438c9cf271c06e ad8ac1ed966b2dc6
$t = 48$ :	d9c523173b1a1e05 e301dca32c44ca05	807f45784852303f 082ee70d3f352aac	985ba373680b8e94 7d4c0abc676b1a8b	a7dcaffdbefb9d4a 9e46e9f915099c34
$t = 49$ :	b6df019ca515cafb 754b3a461a665640	d9c523173b1a1e05 e301dca32c44ca05	807f45784852303f 082ee70d3f352aac	985ba373680b8e94 7d4c0abc676b1a8b
$t = 50$ :	427a642921b2e645 08a30fefe981f2ec	b6df019ca515cafb 754b3a461a665640	d9c523173b1a1e05 e301dca32c44ca05	807f45784852303f 082ee70d3f352aac
$t = 51$ :	7aab58dbelb9df7b 2749c52d0b3d1225	427a642921b2e645 08a30fefe981f2ec	b6df019ca515cafb 754b3a461a665640	d9c523173b1a1e05 e301dca32c44ca05
$t = 52$ :	974ddd552aec16ce a9e6cbfb416a591f	7aab58dbelb9df7b 2749c52d0b3d1225	427a642921b2e645 08a30fefe981f2ec	b6df019ca515cafb 754b3a461a665640
$t = 53$ :	55e0b99d4404f6ca 6c24ad697b41b1b9	974ddd552aec16ce a9e6cbfb416a591f	7aab58dbelb9df7b 2749c52d0b3d1225	427a642921b2e645 08a30fefe981f2ec
$t = 54$ :	901f632579ee1eee 4ee99476db1bb7a9	55e0b99d4404f6ca 6c24ad697b41b1b9	974ddd552aec16ce a9e6cbfb416a591f	7aab58dbelb9df7b 2749c52d0b3d1225
$t = 55$ :	f90db9f292a60463 5401644992a1f8b8	901f632579ee1eee 4ee99476db1bb7a9	55e0b99d4404f6ca 6c24ad697b41b1b9	974ddd552aec16ce a9e6cbfb416a591f
$t = 56$ :	9b906a7df1007357 f5e402ee21db8915	f90db9f292a60463 5401644992a1f8b8	901f632579ee1eee 4ee99476db1bb7a9	55e0b99d4404f6ca 6c24ad697b41b1b9
$t = 57$ :	71a0a998fb48c0fc 96bece755cd203cb	9b906a7df1007357 f5e402ee21db8915	f90db9f292a60463 5401644992a1f8b8	901f632579ee1eee 4ee99476db1bb7a9
$t = 58$ :	c25e798e50752535 9d548440d8e110f2	71a0a998fb48c0fc 96bece755cd203cb	9b906a7df1007357 f5e402ee21db8915	f90db9f292a60463 5401644992a1f8b8
$t = 59$ :	1ce4f2591812e6ae b27252537a83cf27	c25e798e50752535 9d548440d8e110f2	71a0a998fb48c0fc 96bece755cd203cb	9b906a7df1007357 f5e402ee21db8915
$t = 60$ :	c1700e250dc6ffed 970088839126bda5	1ce4f2591812e6ae b27252537a83cf27	c25e798e50752535 9d548440d8e110f2	71a0a998fb48c0fc 96bece755cd203cb
$t = 61$ :	f8e6924412fd0c64 d50cf4f73910e3ee	c1700e250dc6ffed 970088839126bda5	1ce4f2591812e6ae b27252537a83cf27	c25e798e50752535 9d548440d8e110f2

$t = 62$ :	d53e0a39eee47528 1b6d7234ace15d7d	f8e6924412fd0c64 d50cf4f73910e3ee	c1700e250dc6ffed 970088839126bda5	1ce4f2591812e6ae b27252537a83cf27
$t = 63$ :	3960545ab926c0d5 9eabb5618b4fcd13	d53e0a39eee47528 1b6d7234ace15d7d	f8e6924412fd0c64 d50cf4f73910e3ee	c1700e250dc6ffed 970088839126bda5
$t = 64$ :	b2c164d71abb92fe f1736fbbfb6ebe72	3960545ab926c0d5 9eabb5618b4fcd13	d53e0a39eee47528 1b6d7234ace15d7d	f8e6924412fd0c64 d50cf4f73910e3ee
$t = 65$ :	4d979e985b067e75 d1fb300f35992350	b2c164d71abb92fe f1736fbbfb6ebe72	3960545ab926c0d5 9eabb5618b4fcd13	d53e0a39eee47528 1b6d7234ace15d7d
$t = 66$ :	59d0238ce137abd7 5f3c64b7546e2cec	4d979e985b067e75 d1fb300f35992350	b2c164d71abb92fe f1736fbbfb6ebe72	3960545ab926c0d5 9eabb5618b4fcd13
$t = 67$ :	bf8d9453b9876b0a 6c27893a31b0e07e	59d0238ce137abd7 5f3c64b7546e2cec	4d979e985b067e75 d1fb300f35992350	b2c164d71abb92fe f1736fbbfb6ebe72
$t = 68$ :	c45dd4a2d2fea059 48253e21b26d8cf9	bf8d9453b9876b0a 6c27893a31b0e07e	59d0238ce137abd7 5f3c64b7546e2cec	4d979e985b067e75 d1fb300f35992350
$t = 69$ :	e08471946c17b0b6 714e2adf4e23ff24	c45dd4a2d2fea059 48253e21b26d8cf9	bf8d9453b9876b0a 6c27893a31b0e07e	59d0238ce137abd7 5f3c64b7546e2cec
$t = 70$ :	b4838c1c28fee7bc 371f12f333f7e5b9	e08471946c17b0b6 714e2adf4e23ff24	c45dd4a2d2fea059 48253e21b26d8cf9	bf8d9453b9876b0a 6c27893a31b0e07e
$t = 71$ :	851cf60a77f6e6d1 a2a475deac0e8b42	b4838c1c28fee7bc 371f12f333f7e5b9	e08471946c17b0b6 714e2adf4e23ff24	c45dd4a2d2fea059 48253e21b26d8cf9
$t = 72$ :	f53d23c50249af2d 1e99cae9d4cf0409	851cf60a77f6e6d1 a2a475deac0e8b42	b4838c1c28fee7bc 371f12f333f7e5b9	e08471946c17b0b6 714e2adf4e23ff24
$t = 73$ :	b81e85d427045550 f5794711faa60f63	f53d23c50249af2d 1e99cae9d4cf0409	851cf60a77f6e6d1 a2a475deac0e8b42	b4838c1c28fee7bc 371f12f333f7e5b9
$t = 74$ :	ae70c7d11ea84a83 dc0d633411c289b2	b81e85d427045550 f5794711faa60f63	f53d23c50249af2d 1e99cae9d4cf0409	851cf60a77f6e6d1 a2a475deac0e8b42
$t = 75$ :	5c54592e13c76135 1620dd5479e94b9b	ae70c7d11ea84a83 dc0d633411c289b2	b81e85d427045550 f5794711faa60f63	f53d23c50249af2d 1e99cae9d4cf0409
$t = 76$ :	03a0f79087078a93 57e90fa678e4cc97	5c54592e13c76135 1620dd5479e94b9b	ae70c7d11ea84a83 dc0d633411c289b2	b81e85d427045550 f5794711faa60f63
$t = 77$ :	8df0baad4c6ed50c c6e7246f7f0bdac6	03a0f79087078a93 57e90fa678e4cc97	5c54592e13c76135 1620dd5479e94b9b	ae70c7d11ea84a83 dc0d633411c289b2
$t = 78$ :	bfa9f194894db5b6 90bb8597bb41da1a	8df0baad4c6ed50c c6e7246f7f0bdac6	03a0f79087078a93 57e90fa678e4cc97	5c54592e13c76135 1620dd5479e94b9b
$t = 79$ :	4b7c99fbaf72a571 78955227fde03a42	bfa9f194894db5b6 90bb8597bb41da1a	8df0baad4c6ed50c c6e7246f7f0bdac6	03a0f79087078a93 57e90fa678e4cc97

That completes the processing of the second and final message block,  $M^{(2)}$ . The final hash value,  $H^{(2)}$ , is calculated to be

$$\begin{aligned}
 H_0^{(2)} &= 4319017a2b706e69 + 4b7c99fbaf72a571 = 8e959b75dae313da \\
 H_1^{(2)} &= cd4b05938bae5e89 + bfa9f194894db5b6 = 8cf4f72814fc143f \\
 H_2^{(2)} &= 0186bf199f30aa95 + 8df0baad4c6ed50c = 8f7779c6eb9f7fa1 \\
 H_3^{(2)} &= 6ef8b71d2f810585 + 03a0f79087078a93 = 7299aeadb6889018 \\
 H_4^{(2)} &= d787d6764b20bda2 + 78955227fde03a42 = 501d289e4900f7e4
 \end{aligned}$$

$$\begin{aligned}
 H_5^{(2)} &= \text{a260144709736920} + \text{90bb8597bb41da1a} = \text{331b99dec4b5433a} \\
 H_6^{(2)} &= \text{00ec057f37d14b8e} + \text{c6e7246f7f0bdac6} = \text{c7d329eeb6dd2654} \\
 H_7^{(2)} &= \text{06add5b50e671c72} + \text{57e90fa678e4cc97} = \text{5e96e55b874be909} .
 \end{aligned}$$

The resulting 512-bit message digest is

```

8e959b75dae313da 8cf4f72814fc143f 8f7779c6eb9f7fa1 7299aeadb6889018
501d289e4900f7e4 331b99dec4b5433a c7d329eeb6dd2654 5e96e55b874be909 .

```

### C.3 SHA-512 Example (Long Message)

Let the message  $M$  be the binary-coded form of the ASCII string which consists of 1,000,000 repetitions of the character “a”. The resulting SHA-512 message digest is

```

e718483d0ce76964 4e2e42c7bc15b463 8e1f98b13b204428 5632a803afa973eb
de0ff244877ea60a 4cb0432ce577c31b eb009c5c2c49aa2e 4eadb217ad8cc09b .

```

## APPENDIX D: SHA-384 EXAMPLES

This appendix is for informational purposes only and is not required to meet the standard.

### D.1 SHA-384 Example (One-Block Message)

Let the message,  $M$ , be the 24-bit ( $\ell = 24$ ) ASCII string "abc", which is equivalent to the following binary string:

01100001 01100010 01100011.

The message is padded by appending a "1" bit, followed by 871 "0" bits, and ending with the hex value

0000000000000000 0000000000000018

(the two 64-bit word representation of the length, 24). Thus, the final padded message consists of one block ( $N = 1$ ).

For SHA-384, the initial hash value,  $H^{(0)}$ , is

$H_0^{(0)} = \text{cbbb9d5dc1059ed8}$   
 $H_1^{(0)} = \text{629a292a367cd507}$   
 $H_2^{(0)} = \text{9159015a3070dd17}$   
 $H_3^{(0)} = \text{152fec8d8f70e5939}$   
 $H_4^{(0)} = \text{67332667ffc00b31}$   
 $H_5^{(0)} = \text{8eb44a8768581511}$   
 $H_6^{(0)} = \text{db0c2e0d64f98fa7}$   
 $H_7^{(0)} = \text{47b5481dbefa4fa4}.$

The words of the padded message block are then assigned to the words  $W_0, \dots, W_{15}$  of the message schedule:

$W_0 = \text{6162638000000000}$        $W_8 = \text{0000000000000000}$   
 $W_1 = \text{0000000000000000}$        $W_9 = \text{0000000000000000}$   
 $W_2 = \text{0000000000000000}$        $W_{10} = \text{0000000000000000}$   
 $W_3 = \text{0000000000000000}$        $W_{11} = \text{0000000000000000}$   
 $W_4 = \text{0000000000000000}$        $W_{12} = \text{0000000000000000}$   
 $W_5 = \text{0000000000000000}$        $W_{13} = \text{0000000000000000}$   
 $W_6 = \text{0000000000000000}$        $W_{14} = \text{0000000000000000}$   
 $W_7 = \text{0000000000000000}$        $W_{15} = \text{0000000000000018}.$

The following schedule shows the hex values for *a*, *b*, *c*, *d*, *e*, *f*, *g*, and *h* after pass *t* of the “for *t* = 0 to 79” loop described in Sec. 6.3.2, step 4.

	<i>a</i> / <i>e</i>	<i>b</i> / <i>f</i>	<i>c</i> / <i>g</i>	<i>d</i> / <i>h</i>
<i>t</i> = 0 :	470994ad30873f88 bd03f724be6075f9	cbbb9d5dc1059ed8 67332667ffc00b31	629a292a367cd507 8eb44a8768581511	9159015a3070dd17 db0c2e0d64f98fa7
<i>t</i> = 1 :	2e91230306a12ae0 5e1b4e1695372b9e	470994ad30873f88 bd03f724be6075f9	cbbb9d5dc1059ed8 67332667ffc00b31	629a292a367cd507 8eb44a8768581511
<i>t</i> = 2 :	eebe5d379be707ad 54074a65aef34336	2e91230306a12ae0 5e1b4e1695372b9e	470994ad30873f88 bd03f724be6075f9	cbbb9d5dc1059ed8 67332667ffc00b31
<i>t</i> = 3 :	e308483153e15ad6 086c5b2d36a89178	eebe5d379be707ad 54074a65aef34336	2e91230306a12ae0 5e1b4e1695372b9e	470994ad30873f88 bd03f724be6075f9
<i>t</i> = 4 :	3a7a023c593d8479 8aa1144850633794	e308483153e15ad6 086c5b2d36a89178	eebe5d379be707ad 54074a65aef34336	2e91230306a12ae0 5e1b4e1695372b9e
<i>t</i> = 5 :	333199a85f92b052 7a6316f0ef047ce7	3a7a023c593d8479 8aa1144850633794	e308483153e15ad6 086c5b2d36a89178	eebe5d379be707ad 54074a65aef34336
<i>t</i> = 6 :	76f0741213dd2ef6 74063cba385f0675	333199a85f92b052 7a6316f0ef047ce7	3a7a023c593d8479 8aa1144850633794	e308483153e15ad6 086c5b2d36a89178
<i>t</i> = 7 :	02f2a04d3aab1629 1688b9bf14980fc0	76f0741213dd2ef6 74063cba385f0675	333199a85f92b052 7a6316f0ef047ce7	3a7a023c593d8479 8aa1144850633794
<i>t</i> = 8 :	73e5b2a1704a0349 fd00139f705907d0	02f2a04d3aab1629 1688b9bf14980fc0	76f0741213dd2ef6 74063cba385f0675	333199a85f92b052 7a6316f0ef047ce7
<i>t</i> = 9 :	bf3f67ba12882648 652e311d4f0a4257	73e5b2a1704a0349 fd00139f705907d0	02f2a04d3aab1629 1688b9bf14980fc0	76f0741213dd2ef6 74063cba385f0675
<i>t</i> = 10 :	33254508bb2ea48d 9e18991c4f39f0ba	bf3f67ba12882648 652e311d4f0a4257	73e5b2a1704a0349 fd00139f705907d0	02f2a04d3aab1629 1688b9bf14980fc0
<i>t</i> = 11 :	c1fdb2a0205ea0e5 04732e8bc4044582	33254508bb2ea48d 9e18991c4f39f0ba	bf3f67ba12882648 652e311d4f0a4257	73e5b2a1704a0349 fd00139f705907d0
<i>t</i> = 12 :	185f9ff038a50f39 8b4acfc4d2b8afe6	c1fdb2a0205ea0e5 04732e8bc4044582	33254508bb2ea48d 9e18991c4f39f0ba	bf3f67ba12882648 652e311d4f0a4257
<i>t</i> = 13 :	e5f06744c0d7563a 2fa93d1ce9523015	185f9ff038a50f39 8b4acfc4d2b8afe6	c1fdb2a0205ea0e5 04732e8bc4044582	33254508bb2ea48d 9e18991c4f39f0ba
<i>t</i> = 14 :	7e32dc0e9f414783 3a9950aaa5e75884	e5f06744c0d7563a 2fa93d1ce9523015	185f9ff038a50f39 8b4acfc4d2b8afe6	c1fdb2a0205ea0e5 04732e8bc4044582
<i>t</i> = 15 :	1eab6159ae87ef6d 153b895cfbc436c5	7e32dc0e9f414783 3a9950aaa5e75884	e5f06744c0d7563a 2fa93d1ce9523015	185f9ff038a50f39 8b4acfc4d2b8afe6
<i>t</i> = 16 :	33ef2cebbf1739aa 9d1a64baf1d366aa	1eab6159ae87ef6d 153b895cfbc436c5	7e32dc0e9f414783 3a9950aaa5e75884	e5f06744c0d7563a 2fa93d1ce9523015
<i>t</i> = 17 :	7df1b65f1b87d6ca 5b6e369d36e8e181	33ef2cebbf1739aa 9d1a64baf1d366aa	1eab6159ae87ef6d 153b895cfbc436c5	7e32dc0e9f414783 3a9950aaa5e75884
<i>t</i> = 18 :	63a24014a34bb0f6 e13e610eae680d85	7df1b65f1b87d6ca 5b6e369d36e8e181	33ef2cebbf1739aa 9d1a64baf1d366aa	1eab6159ae87ef6d 153b895cfbc436c5
<i>t</i> = 19 :	f1aabd313309509b 674385f0d87db94f	63a24014a34bb0f6 e13e610eae680d85	7df1b65f1b87d6ca 5b6e369d36e8e181	33ef2cebbf1739aa 9d1a64baf1d366aa

$t = 20$ :	9ba737ae88a72c64 3fc2614c43906c0f	f1aabd313309509b 674385f0d87db94f	63a24014a34bb0f6 e13e610eae680d85	7df1b65f1b87d6ca 5b6e369d36e8e181
$t = 21$ :	042c2dc9a5bf558a 19316bebc88e01f2	9ba737ae88a72c64 3fc2614c43906c0f	f1aabd313309509b 674385f0d87db94f	63a24014a34bb0f6 e13e610eae680d85
$t = 22$ :	7799c75acc748c0f a7bbd65bf64f58c8	042c2dc9a5bf558a 19316bebc88e01f2	9ba737ae88a72c64 3fc2614c43906c0f	f1aabd313309509b 674385f0d87db94f
$t = 23$ :	ccf99a80f92bf002 e52a24fae4e8fc9b	7799c75acc748c0f a7bbd65bf64f58c8	042c2dc9a5bf558a 19316bebc88e01f2	9ba737ae88a72c64 3fc2614c43906c0f
$t = 24$ :	ae993474363efe68 587f308d58681928	ccf99a80f92bf002 e52a24fae4e8fc9b	7799c75acc748c0f a7bbd65bf64f58c8	042c2dc9a5bf558a 19316bebc88e01f2
$t = 25$ :	335063d1a2aec92f c2d6d65e38c6ea79	ae993474363efe68 587f308d58681928	ccf99a80f92bf002 e52a24fae4e8fc9b	7799c75acc748c0f a7bbd65bf64f58c8
$t = 26$ :	53a78b0cca01ba37 3b65a26c3c92c8f3	335063d1a2aec92f c2d6d65e38c6ea79	ae993474363efe68 587f308d58681928	ccf99a80f92bf002 e52a24fae4e8fc9b
$t = 27$ :	ab7ffa529f622930 b9d8a2f2762901ea	53a78b0cca01ba37 3b65a26c3c92c8f3	335063d1a2aec92f c2d6d65e38c6ea79	ae993474363efe68 587f308d58681928
$t = 28$ :	e428bb43afe3d63e 6a8527525f898726	ab7ffa529f622930 b9d8a2f2762901ea	53a78b0cca01ba37 3b65a26c3c92c8f3	335063d1a2aec92f c2d6d65e38c6ea79
$t = 29$ :	bbed541a5128088c 7973aadbde294be9	e428bb43afe3d63e 6a8527525f898726	ab7ffa529f622930 b9d8a2f2762901ea	53a78b0cca01ba37 3b65a26c3c92c8f3
$t = 30$ :	4c5c38df7ec8baf4 422ceea0200e9ee4	bbed541a5128088c 7973aadbde294be9	e428bb43afe3d63e 6a8527525f898726	ab7ffa529f622930 b9d8a2f2762901ea
$t = 31$ :	4ba456ec244033ed 7cf40857056d86b0	4c5c38df7ec8baf4 422ceea0200e9ee4	bbed541a5128088c 7973aadbde294be9	e428bb43afe3d63e 6a8527525f898726
$t = 32$ :	aa4a6ab2ac5f5dd8 ad2b1ecfb5bfc556	4ba456ec244033ed 7cf40857056d86b0	4c5c38df7ec8baf4 422ceea0200e9ee4	bbed541a5128088c 7973aadbde294be9
$t = 33$ :	9cb941f2ced774b3 029f66c7b4569bf0	aa4a6ab2ac5f5dd8 ad2b1ecfb5bfc556	4ba456ec244033ed 7cf40857056d86b0	4c5c38df7ec8baf4 422ceea0200e9ee4
$t = 34$ :	39265f358594de27 3f7b1c260c82e54f	9cb941f2ced774b3 029f66c7b4569bf0	aa4a6ab2ac5f5dd8 ad2b1ecfb5bfc556	4ba456ec244033ed 7cf40857056d86b0
$t = 35$ :	09cca487d39b02a1 4a22b37b58a5b1b0	39265f358594de27 3f7b1c260c82e54f	9cb941f2ced774b3 029f66c7b4569bf0	aa4a6ab2ac5f5dd8 ad2b1ecfb5bfc556
$t = 36$ :	d48d97ce438cf4f0 a239e00b8baa0410	09cca487d39b02a1 4a22b37b58a5b1b0	39265f358594de27 3f7b1c260c82e54f	9cb941f2ced774b3 029f66c7b4569bf0
$t = 37$ :	d6f41e25a8b634d6 25755cb8179dd0b0	d48d97ce438cf4f0 a239e00b8baa0410	09cca487d39b02a1 4a22b37b58a5b1b0	39265f358594de27 3f7b1c260c82e54f
$t = 38$ :	54078334358573b4 0e419fb0802b0efc	d6f41e25a8b634d6 25755cb8179dd0b0	d48d97ce438cf4f0 a239e00b8baa0410	09cca487d39b02a1 4a22b37b58a5b1b0
$t = 39$ :	db24f9a03f4ffff6b d30e99b4b394b090	54078334358573b4 0e419fb0802b0efc	d6f41e25a8b634d6 25755cb8179dd0b0	d48d97ce438cf4f0 a239e00b8baa0410
$t = 40$ :	3604c53a845efc37 791b2b4af7338b99	db24f9a03f4ffff6b d30e99b4b394b090	54078334358573b4 0e419fb0802b0efc	d6f41e25a8b634d6 25755cb8179dd0b0
$t = 41$ :	f41b1c0eee89bdc6 e319b77d9e4e87f9	3604c53a845efc37 791b2b4af7338b99	db24f9a03f4ffff6b d30e99b4b394b090	54078334358573b4 0e419fb0802b0efc
$t = 42$ :	36644ae374632e3a 458250878a3972b2	f41b1c0eee89bdc6 e319b77d9e4e87f9	3604c53a845efc37 791b2b4af7338b99	db24f9a03f4ffff6b d30e99b4b394b090
$t = 43$ :	88806f6ae9fcd65b	36644ae374632e3a	f41b1c0eee89bdc6	3604c53a845efc37

	cfde2e6ea54fa576	458250878a3972b2	e319b77d9e4e87f9	791b2b4af7338b99
$t = 44 :$	51dcaa36995c301d e37f778353998050	88806f6ae9fcd65b cfde2e6ea54fa576	36644ae374632e3a 458250878a3972b2	f41b1c0eee89bdc6 e319b77d9e4e87f9
$t = 45 :$	ef5e3885a2f238df 740e347f24e18fda	51dcaa36995c301d e37f778353998050	88806f6ae9fcd65b cfde2e6ea54fa576	36644ae374632e3a 458250878a3972b2
$t = 46 :$	eb3753f4283f4818 0ae48cf840bb8be9	ef5e3885a2f238df 740e347f24e18fda	51dcaa36995c301d e37f778353998050	88806f6ae9fcd65b cfde2e6ea54fa576
$t = 47 :$	a6998d63a5d09e04 e21095012ee0b72a	eb3753f4283f4818 0ae48cf840bb8be9	ef5e3885a2f238df 740e347f24e18fda	51dcaa36995c301d e37f778353998050
$t = 48 :$	d3698fb64df175b0 c2f0b90ffce80739	a6998d63a5d09e04 e21095012ee0b72a	eb3753f4283f4818 0ae48cf840bb8be9	ef5e3885a2f238df 740e347f24e18fda
$t = 49 :$	317a3b295b991914 1cadff2e6cb5aa4d	d3698fb64df175b0 c2f0b90ffce80739	a6998d63a5d09e04 e21095012ee0b72a	eb3753f4283f4818 0ae48cf840bb8be9
$t = 50 :$	0941da08148ba463 833eb9a4bb5a073e	317a3b295b991914 1cadff2e6cb5aa4d	d3698fb64df175b0 c2f0b90ffce80739	a6998d63a5d09e04 e21095012ee0b72a
$t = 51 :$	494ac238d68c3d0b 80c8fc138e645028	0941da08148ba463 833eb9a4bb5a073e	317a3b295b991914 1cadff2e6cb5aa4d	d3698fb64df175b0 c2f0b90ffce80739
$t = 52 :$	c87e9168db9e97de 65cf7f6a829aca04	494ac238d68c3d0b 80c8fc138e645028	0941da08148ba463 833eb9a4bb5a073e	317a3b295b991914 1cadff2e6cb5aa4d
$t = 53 :$	edb4448879391dbb 7729c85475dd318f	c87e9168db9e97de 65cf7f6a829aca04	494ac238d68c3d0b 80c8fc138e645028	0941da08148ba463 833eb9a4bb5a073e
$t = 54 :$	073775c2456dc7db a9cca0b6266b1d77	edb4448879391dbb 7729c85475dd318f	c87e9168db9e97de 65cf7f6a829aca04	494ac238d68c3d0b 80c8fc138e645028
$t = 55 :$	54de8857b24afaf7 8de51cff2ae4b068	073775c2456dc7db a9cca0b6266b1d77	edb4448879391dbb 7729c85475dd318f	c87e9168db9e97de 65cf7f6a829aca04
$t = 56 :$	8a9cdd80f7f09c05 a60ba5e9ebaeb96a	54de8857b24afaf7 8de51cff2ae4b068	073775c2456dc7db a9cca0b6266b1d77	edb4448879391dbb 7729c85475dd318f
$t = 57 :$	3eeb22a7524d8d7f e2e6830b139df58f	8a9cdd80f7f09c05 a60ba5e9ebaeb96a	54de8857b24afaf7 8de51cff2ae4b068	073775c2456dc7db a9cca0b6266b1d77
$t = 58 :$	0ed77c9cde8883d3 38413a2052387a9e	3eeb22a7524d8d7f e2e6830b139df58f	8a9cdd80f7f09c05 a60ba5e9ebaeb96a	54de8857b24afaf7 8de51cff2ae4b068
$t = 59 :$	e64e4135f9d30dbc 45b640454c75c349	0ed77c9cde8883d3 38413a2052387a9e	3eeb22a7524d8d7f e2e6830b139df58f	8a9cdd80f7f09c05 a60ba5e9ebaeb96a
$t = 60 :$	1ca93a293d544328 efbef83a35c0319e	e64e4135f9d30dbc 45b640454c75c349	0ed77c9cde8883d3 38413a2052387a9e	3eeb22a7524d8d7f e2e6830b139df58f
$t = 61 :$	3dc764f89e54043a a57784945550cf94	1ca93a293d544328 efbef83a35c0319e	e64e4135f9d30dbc 45b640454c75c349	0ed77c9cde8883d3 38413a2052387a9e
$t = 62 :$	56fb5883f1c87a05 f5198a41eb80e022	3dc764f89e54043a a57784945550cf94	1ca93a293d544328 efbef83a35c0319e	e64e4135f9d30dbc 45b640454c75c349
$t = 63 :$	24a1124262a331c7 06edacae6e7b54ad	56fb5883f1c87a05 f5198a41eb80e022	3dc764f89e54043a a57784945550cf94	1ca93a293d544328 efbef83a35c0319e
$t = 64 :$	eb85d19201c89694 9ced24983eec8723	24a1124262a331c7 06edacae6e7b54ad	56fb5883f1c87a05 f5198a41eb80e022	3dc764f89e54043a a57784945550cf94
$t = 65 :$	cc981ab3a59c1db4 eac5516336bc8882	eb85d19201c89694 9ced24983eec8723	24a1124262a331c7 06edacae6e7b54ad	56fb5883f1c87a05 f5198a41eb80e022
$t = 66 :$	ceef5d997e148b44 617bbf70bb165212	cc981ab3a59c1db4 eac5516336bc8882	eb85d19201c89694 9ced24983eec8723	24a1124262a331c7 06edacae6e7b54ad



$t = 67$ :	689edf608a8e3f14 3280d88472c100fd	ceef5d997e148b44 617bbf70bb165212	cc981ab3a59c1db4 eac5516336bc8882	eb85d19201c89694 9ced24983eec8723
$t = 68$ :	1e6e0255ab88079f f2001138439902b1	689edf608a8e3f14 3280d88472c100fd	ceef5d997e148b44 617bbf70bb165212	cc981ab3a59c1db4 eac5516336bc8882
$t = 69$ :	8c5d3b7fdad66e70 90d18ec8b69f0345	1e6e0255ab88079f f2001138439902b1	689edf608a8e3f14 3280d88472c100fd	ceef5d997e148b44 617bbf70bb165212
$t = 70$ :	32e5ed8655871e9b 51105f6241313777	8c5d3b7fdad66e70 90d18ec8b69f0345	1e6e0255ab88079f f2001138439902b1	689edf608a8e3f14 3280d88472c100fd
$t = 71$ :	bcd5061679be7336 454b99f654443ad0	32e5ed8655871e9b 51105f6241313777	8c5d3b7fdad66e70 90d18ec8b69f0345	1e6e0255ab88079f f2001138439902b1
$t = 72$ :	e7d913b6678e78ef 1ff613b5aa63776e	bcd5061679be7336 454b99f654443ad0	32e5ed8655871e9b 51105f6241313777	8c5d3b7fdad66e70 90d18ec8b69f0345
$t = 73$ :	e6b8cb8dfa3475ab 2e75f34303d39bb0	e7d913b6678e78ef 1ff613b5aa63776e	bcd5061679be7336 454b99f654443ad0	32e5ed8655871e9b 51105f6241313777
$t = 74$ :	fdd4a30e168c4ae5 83a35dbe2a64fc26	e6b8cb8dfa3475ab 2e75f34303d39bb0	e7d913b6678e78ef 1ff613b5aa63776e	bcd5061679be7336 454b99f654443ad0
$t = 75$ :	12aeb6268dfa3e14 f660943b276786f7	fdd4a30e168c4ae5 83a35dbe2a64fc26	e6b8cb8dfa3475ab 2e75f34303d39bb0	e7d913b6678e78ef 1ff613b5aa63776e
$t = 76$ :	055b73814cf102b4 c4b149710f5d6a71	12aeb6268dfa3e14 f660943b276786f7	fdd4a30e168c4ae5 83a35dbe2a64fc26	e6b8cb8dfa3475ab 2e75f34303d39bb0
$t = 77$ :	95d33150de6df44c c7f7bff08ebf0d30	055b73814cf102b4 c4b149710f5d6a71	12aeb6268dfa3e14 f660943b276786f7	fdd4a30e168c4ae5 83a35dbe2a64fc26
$t = 78$ :	5306143f64497b00 ca06a219cc701096	95d33150de6df44c c7f7bff08ebf0d30	055b73814cf102b4 c4b149710f5d6a71	12aeb6268dfa3e14 f660943b276786f7
$t = 79$ :	ff44d7e1849dbfb3 1952e0c3a227c0f2	5306143f64497b00 ca06a219cc701096	95d33150de6df44c c7f7bff08ebf0d30	055b73814cf102b4 c4b149710f5d6a71

That completes the processing of the first and only message block,  $M^{(1)}$ . The final hash value,  $H^{(1)}$ , is calculated to be

$$\begin{aligned}
 H_0^{(1)} &= \text{cbbb9d5dc1059ed8} + \text{ff44d7e1849dbfb3} = \text{cb00753f45a35e8b} \\
 H_1^{(1)} &= \text{629a292a367cd507} + \text{5306143f64497b00} = \text{b5a03d699ac65007} \\
 H_2^{(1)} &= \text{9159015a3070dd17} + \text{95d33150de6df44c} = \text{272c32ab0eded163} \\
 H_3^{(1)} &= \text{152fec8df70e5939} + \text{055b73814cf102b4} = \text{1a8b605a43ff5bed} \\
 H_4^{(1)} &= \text{67332667ffc00b31} + \text{1952e0c3a227c0f2} = \text{8086072ba1e7cc23} \\
 H_5^{(1)} &= \text{8eb44a8768581511} + \text{ca06a219cc701096} = \text{58baeca134c825a7} \\
 H_6^{(1)} &= \text{db0c2e0d64f98fa7} + \text{c7f7bff08ebf0d30} = \text{a303edfdf3b89cd7} \\
 H_7^{(1)} &= \text{47b5481dbefa4fa4} + \text{c4b149710f5d6a71} = \text{0c66918ece57ba15} .
 \end{aligned}$$

The final hash value is truncated to its left-most 384 bits (i.e.,  $H_0^{(1)}, \dots, H_5^{(1)}$ ), resulting in the 384-bit message digest

```
cb00753f45a35e8b b5a03d699ac65007 272c32ab0eded163 1a8b605a43ff5bed
8086072ba1e7cc23 58baeca134c825a7.
```



## D.2 SHA-384 Example (Multi-Block Message)

Let the message,  $M$ , be the 896-bit ( $\ell = 896$ ) ASCII string

**"abcdefghijklmghijklmn  
hijklmnoijklmnopqklmnopqrlmnopqrsmnopqrstu".**

The message is padded by appending a '1' bit, followed by 1023 '0' bits, and ending with the hex value

0000000000000000 0000000000000380

(the two 64-bit word representation of the length, 896). Thus, the final padded message consists of two blocks ( $N = 2$ ).

For SHA-384, the initial hash value,  $H^{(0)}$ , is

$H_0^{(0)} = \text{cbbb9d5dc1059ed8}$   
 $H_1^{(0)} = \text{629a292a367cd507}$   
 $H_2^{(0)} = \text{9159015a3070dd17}$   
 $H_3^{(0)} = \text{152fec8d8f70e5939}$   
 $H_4^{(0)} = \text{67332667ffc00b31}$   
 $H_5^{(0)} = \text{8eb44a8768581511}$   
 $H_6^{(0)} = \text{db0c2e0d64f98fa7}$   
 $H_7^{(0)} = \text{47b5481dbefa4fa4}.$

The words of the padded message block are then assigned to the words  $W_0, \dots, W_{15}$  of the message schedule:

$W_0 = \text{6162636465666768}$	$W_8 = \text{696a6b6c6d6e6f70}$
$W_1 = \text{6263646566676869}$	$W_9 = \text{6a6b6c6d6e6f7071}$
$W_2 = \text{636465666768696a}$	$W_{10} = \text{6b6c6d6e6f707172}$
$W_3 = \text{6465666768696a6b}$	$W_{11} = \text{6c6d6e6f70717273}$
$W_4 = \text{65666768696a6b6c}$	$W_{12} = \text{6d6e6f7071727374}$
$W_5 = \text{666768696a6b6c6d}$	$W_{13} = \text{6e6f707172737475}$
$W_6 = \text{6768696a6b6c6d6e}$	$W_{14} = \text{8000000000000000}$
$W_7 = \text{68696a6b6c6d6e6f}$	$W_{15} = \text{0000000000000000}.$

The following schedule shows the hex values for  $a, b, c, d, e, f, g,$  and  $h$  after pass  $t$  of the “for  $t = 0$  to 79” loop described in Sec. 6.3.2, step 4.

	<b>a</b> / <b>e</b>	<b>b</b> / <b>f</b>	<b>c</b> / <b>g</b>	<b>d</b> / <b>h</b>
$t = 0 :$	4709949195eda6f0 bd03f70923c6dd61	cbbb9d5dc1059ed8 67332667ffc00b31	629a292a367cd507 8eb44a8768581511	9159015a3070dd17 db0c2e0d64f98fa7
$t = 1 :$	78d3f8bc03a38303 ae067f071cd18a36	4709949195eda6f0 bd03f70923c6dd61	cbbb9d5dc1059ed8 67332667ffc00b31	629a292a367cd507 8eb44a8768581511
$t = 2 :$	ed59d30beff95306 c180c7a74ed5cf1f	78d3f8bc03a38303 ae067f071cd18a36	4709949195eda6f0 bd03f70923c6dd61	cbbb9d5dc1059ed8 67332667ffc00b31
$t = 3 :$	8e7fe2aba3168f2b d92d19667920b327	ed59d30beff95306 c180c7a74ed5cf1f	78d3f8bc03a38303 ae067f071cd18a36	4709949195eda6f0 bd03f70923c6dd61
$t = 4 :$	1174f9b374a9263a dd371f2d13661c52	8e7fe2aba3168f2b d92d19667920b327	ed59d30beff95306 c180c7a74ed5cf1f	78d3f8bc03a38303 ae067f071cd18a36
$t = 5 :$	27aaafb7fbef806b 21af3c6430a9af9c	1174f9b374a9263a dd371f2d13661c52	8e7fe2aba3168f2b d92d19667920b327	ed59d30beff95306 c180c7a74ed5cf1f
$t = 6 :$	b352d03a0bd34d65 69397de9a30e1473	27aaafb7fbef806b 21af3c6430a9af9c	1174f9b374a9263a dd371f2d13661c52	8e7fe2aba3168f2b d92d19667920b327
$t = 7 :$	412db7f990563d7c 5062fd5924e2b62e	b352d03a0bd34d65 69397de9a30e1473	27aaafb7fbef806b 21af3c6430a9af9c	1174f9b374a9263a dd371f2d13661c52
$t = 8 :$	0f79040546e6edf7 6b6c511b25a6bdbc	412db7f990563d7c 5062fd5924e2b62e	b352d03a0bd34d65 69397de9a30e1473	27aaafb7fbef806b 21af3c6430a9af9c
$t = 9 :$	ebf02410f67b8ee7 dac695b91543ae80	0f79040546e6edf7 6b6c511b25a6bdbc	412db7f990563d7c 5062fd5924e2b62e	b352d03a0bd34d65 69397de9a30e1473
$t = 10 :$	97aa05d89b8dbe6d 83b8b72646c0b598	ebf02410f67b8ee7 dac695b91543ae80	0f79040546e6edf7 6b6c511b25a6bdbc	412db7f990563d7c 5062fd5924e2b62e
$t = 11 :$	23d0a36b692118eb a5f6c5155e221e8c	97aa05d89b8dbe6d 83b8b72646c0b598	ebf02410f67b8ee7 dac695b91543ae80	0f79040546e6edf7 6b6c511b25a6bdbc
$t = 12 :$	e1041368d2fca1a2 ae01675bfb003180	23d0a36b692118eb a5f6c5155e221e8c	97aa05d89b8dbe6d 83b8b72646c0b598	ebf02410f67b8ee7 dac695b91543ae80
$t = 13 :$	45bd6f69efec540d c35cc50c1cf7ef98	e1041368d2fca1a2 ae01675bfb003180	23d0a36b692118eb a5f6c5155e221e8c	97aa05d89b8dbe6d 83b8b72646c0b598
$t = 14 :$	c237fa23abb9bc16 a16c4f134b28923e	45bd6f69efec540d c35cc50c1cf7ef98	e1041368d2fca1a2 ae01675bfb003180	23d0a36b692118eb a5f6c5155e221e8c
$t = 15 :$	b4092df1c0f81853 008178e17fa649f2	c237fa23abb9bc16 a16c4f134b28923e	45bd6f69efec540d c35cc50c1cf7ef98	e1041368d2fca1a2 ae01675bfb003180
$t = 16 :$	21e5c91d11809c13 a26dfa04ed8c9b63	b4092df1c0f81853 008178e17fa649f2	c237fa23abb9bc16 a16c4f134b28923e	45bd6f69efec540d c35cc50c1cf7ef98
$t = 17 :$	2c957137cd4304a5 6be210614b10949b	21e5c91d11809c13 a26dfa04ed8c9b63	b4092df1c0f81853 008178e17fa649f2	c237fa23abb9bc16 a16c4f134b28923e
$t = 18 :$	2180e61afe322bc7 76396996200065f7	2c957137cd4304a5 6be210614b10949b	21e5c91d11809c13 a26dfa04ed8c9b63	b4092df1c0f81853 008178e17fa649f2
$t = 19 :$	f2911c11c96e5ff5 1bc2160f4f3711dc	2180e61afe322bc7 76396996200065f7	2c957137cd4304a5 6be210614b10949b	21e5c91d11809c13 a26dfa04ed8c9b63
$t = 20 :$	5eab10b19a5143a8 98d2b19d201f2bb6	f2911c11c96e5ff5 1bc2160f4f3711dc	2180e61afe322bc7 76396996200065f7	2c957137cd4304a5 6be210614b10949b
$t = 21 :$	29c5348d87cd5590	5eab10b19a5143a8	f2911c11c96e5ff5	2180e61afe322bc7

	4324c8caccf7753c	98d2b19d201f2bb6	1bc2160f4f3711dc	76396996200065f7
$t = 22 :$	33c6b4a0166b7c9c d49cef5bd2dec121	29c5348d87cd5590 4324c8caccf7753c	5eab10b19a5143a8 98d2b19d201f2bb6	f2911c11c96e5ff5 1bc2160f4f3711dc
$t = 23 :$	1db4ee606d2a7a96 b17d15b397521ab3	33c6b4a0166b7c9c d49cef5bd2dec121	29c5348d87cd5590 4324c8caccf7753c	5eab10b19a5143a8 98d2b19d201f2bb6
$t = 24 :$	5cef5b2f00142660 789e540f22e13932	1db4ee606d2a7a96 b17d15b397521ab3	33c6b4a0166b7c9c d49cef5bd2dec121	29c5348d87cd5590 4324c8caccf7753c
$t = 25 :$	ff74f4a162435903 6c0be33dcc6e7572	5cef5b2f00142660 789e540f22e13932	1db4ee606d2a7a96 b17d15b397521ab3	33c6b4a0166b7c9c d49cef5bd2dec121
$t = 26 :$	41740b736e9676a9 d8e401251592da6c	ff74f4a162435903 6c0be33dcc6e7572	5cef5b2f00142660 789e540f22e13932	1db4ee606d2a7a96 b17d15b397521ab3
$t = 27 :$	931059fe9279ff1d 7f31116887eea596	41740b736e9676a9 d8e401251592da6c	ff74f4a162435903 6c0be33dcc6e7572	5cef5b2f00142660 789e540f22e13932
$t = 28 :$	356d08d982e2ead4 40c28c34b1bbe906	931059fe9279ff1d 7f31116887eea596	41740b736e9676a9 d8e401251592da6c	ff74f4a162435903 6c0be33dcc6e7572
$t = 29 :$	89dc825e7235c74b 7a499ae05da50bf2	356d08d982e2ead4 40c28c34b1bbe906	931059fe9279ff1d 7f31116887eea596	41740b736e9676a9 d8e401251592da6c
$t = 30 :$	97901f333e662fdc 4472b2e331ddf4b4	89dc825e7235c74b 7a499ae05da50bf2	356d08d982e2ead4 40c28c34b1bbe906	931059fe9279ff1d 7f31116887eea596
$t = 31 :$	69c8f40eb38b6022 177589502dd39aa2	97901f333e662fdc 4472b2e331ddf4b4	89dc825e7235c74b 7a499ae05da50bf2	356d08d982e2ead4 40c28c34b1bbe906
$t = 32 :$	4920943ffe52b207 6b813a0d0cdf4991	69c8f40eb38b6022 177589502dd39aa2	97901f333e662fdc 4472b2e331ddf4b4	89dc825e7235c74b 7a499ae05da50bf2
$t = 33 :$	b4cb0df332d108ab 8fe3d28097f18618	4920943ffe52b207 6b813a0d0cdf4991	69c8f40eb38b6022 177589502dd39aa2	97901f333e662fdc 4472b2e331ddf4b4
$t = 34 :$	e7748fbf744a5240 0d7ab03208f1d7a5	b4cb0df332d108ab 8fe3d28097f18618	4920943ffe52b207 6b813a0d0cdf4991	69c8f40eb38b6022 177589502dd39aa2
$t = 35 :$	7416ca18d9e265e0 11200c2d47c082f8	e7748fbf744a5240 0d7ab03208f1d7a5	b4cb0df332d108ab 8fe3d28097f18618	4920943ffe52b207 6b813a0d0cdf4991
$t = 36 :$	75476f5456e82f9c 3024702447f76224	7416ca18d9e265e0 11200c2d47c082f8	e7748fbf744a5240 0d7ab03208f1d7a5	b4cb0df332d108ab 8fe3d28097f18618
$t = 37 :$	f638a568b53a2f8f 6217c1c02153302c	75476f5456e82f9c 3024702447f76224	7416ca18d9e265e0 11200c2d47c082f8	e7748fbf744a5240 0d7ab03208f1d7a5
$t = 38 :$	c418f6f90602c79a 87f0901c227adbb3	f638a568b53a2f8f 6217c1c02153302c	75476f5456e82f9c 3024702447f76224	7416ca18d9e265e0 11200c2d47c082f8
$t = 39 :$	4f1f4f21df3dcf43 fb7c63fcddf4a1c2	c418f6f90602c79a 87f0901c227adbb3	f638a568b53a2f8f 6217c1c02153302c	75476f5456e82f9c 3024702447f76224
$t = 40 :$	13eb82e4b98d0e67 fb6c0e54d48d4f2d	4f1f4f21df3dcf43 fb7c63fcddf4a1c2	c418f6f90602c79a 87f0901c227adbb3	f638a568b53a2f8f 6217c1c02153302c
$t = 41 :$	820e75046567bace b16a9397472f0123	13eb82e4b98d0e67 fb6c0e54d48d4f2d	4f1f4f21df3dcf43 fb7c63fcddf4a1c2	c418f6f90602c79a 87f0901c227adbb3
$t = 42 :$	741fa5dc290dd02c ed40c88214823792	820e75046567bace b16a9397472f0123	13eb82e4b98d0e67 fb6c0e54d48d4f2d	4f1f4f21df3dcf43 fb7c63fcddf4a1c2
$t = 43 :$	a4809bf6da6aa8bd bec3d7e88c855194	741fa5dc290dd02c ed40c88214823792	820e75046567bace b16a9397472f0123	13eb82e4b98d0e67 fb6c0e54d48d4f2d
$t = 44 :$	d70b1aa4c800979c	a4809bf6da6aa8bd	741fa5dc290dd02c	820e75046567bace

	4962f310bdbd54b0	bec3d7e88c855194	ed40c88214823792	b16a9397472f0123
$t = 45 :$	9a195492cfdb4745 2c82d09cf05cf687	d70b1aa4c800979c 4962f310bdbd54b0	a4809bf6da6aa8bd bec3d7e88c855194	741fa5dc290dd02c ed40c88214823792
$t = 46 :$	b7e68364f07f017e 2a1ffb84031b1b6c	9a195492cfdb4745 2c82d09cf05cf687	d70b1aa4c800979c 4962f310bdbd54b0	a4809bf6da6aa8bd bec3d7e88c855194
$t = 47 :$	0e574b8e0b35e452 29bdab29ee472a23	b7e68364f07f017e 2a1ffb84031b1b6c	9a195492cfdb4745 2c82d09cf05cf687	d70b1aa4c800979c 4962f310bdbd54b0
$t = 48 :$	c176009cf82fa842 cca47fbe31b335f4	0e574b8e0b35e452 29bdab29ee472a23	b7e68364f07f017e 2a1ffb84031b1b6c	9a195492cfdb4745 2c82d09cf05cf687
$t = 49 :$	5d4f78c7a9bbed2 eaf198615e99ffdc	c176009cf82fa842 cca47fbe31b335f4	0e574b8e0b35e452 29bdab29ee472a23	b7e68364f07f017e 2a1ffb84031b1b6c
$t = 50 :$	51ab3be828d8d13c bd527cd188fb59ae	5d4f78c7a9bbed2 eaf198615e99ffdc	c176009cf82fa842 cca47fbe31b335f4	0e574b8e0b35e452 29bdab29ee472a23
$t = 51 :$	4d639ef80d0f6d3e b2611b90f90d732f	51ab3be828d8d13c bd527cd188fb59ae	5d4f78c7a9bbed2 eaf198615e99ffdc	c176009cf82fa842 cca47fbe31b335f4
$t = 52 :$	bba9c9efe0fbc6c8 fc0579337591a2c9	4d639ef80d0f6d3e b2611b90f90d732f	51ab3be828d8d13c bd527cd188fb59ae	5d4f78c7a9bbed2 eaf198615e99ffdc
$t = 53 :$	3405d7cad2e8a689 0f6649f64ec8e109	bba9c9efe0fbc6c8 fc0579337591a2c9	4d639ef80d0f6d3e b2611b90f90d732f	51ab3be828d8d13c bd527cd188fb59ae
$t = 54 :$	ea54d908505798b3 ef48a48999108077	3405d7cad2e8a689 0f6649f64ec8e109	bba9c9efe0fbc6c8 fc0579337591a2c9	4d639ef80d0f6d3e b2611b90f90d732f
$t = 55 :$	be31d1c0ccc143bc 4fc2d4cad0c91afc	ea54d908505798b3 ef48a48999108077	3405d7cad2e8a689 0f6649f64ec8e109	bba9c9efe0fbc6c8 fc0579337591a2c9
$t = 56 :$	285a76d23f6a0073 a730855599b738a3	be31d1c0ccc143bc 4fc2d4cad0c91afc	ea54d908505798b3 ef48a48999108077	3405d7cad2e8a689 0f6649f64ec8e109
$t = 57 :$	a714cefff14bebc24 53c581dae1831d80	285a76d23f6a0073 a730855599b738a3	be31d1c0ccc143bc 4fc2d4cad0c91afc	ea54d908505798b3 ef48a48999108077
$t = 58 :$	697ca14913a50a26 34d39344354aacd2	a714cefff14bebc24 53c581dae1831d80	285a76d23f6a0073 a730855599b738a3	be31d1c0ccc143bc 4fc2d4cad0c91afc
$t = 59 :$	3a38fa3775d7007c e26f3a21e9a27691	697ca14913a50a26 34d39344354aacd2	a714cefff14bebc24 53c581dae1831d80	285a76d23f6a0073 a730855599b738a3
$t = 60 :$	44ea14d8e450c844 5319374fb88dd485	3a38fa3775d7007c e26f3a21e9a27691	697ca14913a50a26 34d39344354aacd2	a714cefff14bebc24 53c581dae1831d80
$t = 61 :$	0928b75c925f91e2 79f4be3c5a372911	44ea14d8e450c844 5319374fb88dd485	3a38fa3775d7007c e26f3a21e9a27691	697ca14913a50a26 34d39344354aacd2
$t = 62 :$	6db5469fa19c0e27 16beec0fec168e79	0928b75c925f91e2 79f4be3c5a372911	44ea14d8e450c844 5319374fb88dd485	3a38fa3775d7007c e26f3a21e9a27691
$t = 63 :$	384e3159898a7362 55fa3ad1102298a8	6db5469fa19c0e27 16beec0fec168e79	0928b75c925f91e2 79f4be3c5a372911	44ea14d8e450c844 5319374fb88dd485
$t = 64 :$	483c64d3fdebfb828 1a238431921ea75e	384e3159898a7362 55fa3ad1102298a8	6db5469fa19c0e27 16beec0fec168e79	0928b75c925f91e2 79f4be3c5a372911
$t = 65 :$	c9464988a1939bcf e3f3f08ac90f86cd	483c64d3fdebfb828 1a238431921ea75e	384e3159898a7362 55fa3ad1102298a8	6db5469fa19c0e27 16beec0fec168e79
$t = 66 :$	98bc93bca795059c 9e04fb49a5fd91de	c9464988a1939bcf e3f3f08ac90f86cd	483c64d3fdebfb828 1a238431921ea75e	384e3159898a7362 55fa3ad1102298a8
$t = 67 :$	b6fc101ad1d74e20 fd13cd3620f6c1f4	98bc93bca795059c 9e04fb49a5fd91de	c9464988a1939bcf e3f3f08ac90f86cd	483c64d3fdebfb828 1a238431921ea75e

$t = 68 :$	fac26e6e4da4705d 0d60228aa6e55b6e	b6fc101ad1d74e20 fd13cd3620f6c1f4	98bc93bca795059c 9e04fb49a5fd91de	c9464988a1939bcf e3f3f08ac90f86cd
$t = 69 :$	2a630c58cc27fcaa a2f7f27a3ec25aba	fac26e6e4da4705d 0d60228aa6e55b6e	b6fc101ad1d74e20 fd13cd3620f6c1f4	98bc93bca795059c 9e04fb49a5fd91de
$t = 70 :$	159a02d4faee11b4 b2860fc55bdedaa6	2a630c58cc27fcaa a2f7f27a3ec25aba	fac26e6e4da4705d 0d60228aa6e55b6e	b6fc101ad1d74e20 fd13cd3620f6c1f4
$t = 71 :$	9d38bdb9df22b557 dfc37c68af65f8bc	159a02d4faee11b4 b2860fc55bdedaa6	2a630c58cc27fcaa a2f7f27a3ec25aba	fac26e6e4da4705d 0d60228aa6e55b6e
$t = 72 :$	d42c3a57cfa78513 bb56dea6a325ba32	9d38bdb9df22b557 dfc37c68af65f8bc	159a02d4faee11b4 b2860fc55bdedaa6	2a630c58cc27fcaa a2f7f27a3ec25aba
$t = 73 :$	abab4b0ca75a17c7 9ac71d1c037a8bbd	d42c3a57cfa78513 bb56dea6a325ba32	9d38bdb9df22b557 dfc37c68af65f8bc	159a02d4faee11b4 b2860fc55bdedaa6
$t = 74 :$	500f7b61186f6c2e 8347f5736531b3ec	abab4b0ca75a17c7 9ac71d1c037a8bbd	d42c3a57cfa78513 bb56dea6a325ba32	9d38bdb9df22b557 dfc37c68af65f8bc
$t = 75 :$	4abe0af6a67db2fe 14e986342ddced0f	500f7b61186f6c2e 8347f5736531b3ec	abab4b0ca75a17c7 9ac71d1c037a8bbd	d42c3a57cfa78513 bb56dea6a325ba32
$t = 76 :$	e1053fc85f9e56be 4779767cc2ec5321	4abe0af6a67db2fe 14e986342ddced0f	500f7b61186f6c2e 8347f5736531b3ec	abab4b0ca75a17c7 9ac71d1c037a8bbd
$t = 77 :$	7001201948fb3d71 5cdf6c58fc052572	e1053fc85f9e56be 4779767cc2ec5321	4abe0af6a67db2fe 14e986342ddced0f	500f7b61186f6c2e 8347f5736531b3ec
$t = 78 :$	88146da76ff6f23a 8901cffe7a74db98	7001201948fb3d71 5cdf6c58fc052572	e1053fc85f9e56be 4779767cc2ec5321	4abe0af6a67db2fe 14e986342ddced0f
$t = 79 :$	5ec3802b9ecfef33 5f2eead69efb4233	88146da76ff6f23a 8901cffe7a74db98	7001201948fb3d71 5cdf6c58fc052572	e1053fc85f9e56be 4779767cc2ec5321

That completes the processing of the first message block,  $M^{(1)}$ . The intermediate hash value,  $H^{(1)}$ , is calculated to be

$$\begin{aligned}
 H_0^{(1)} &= \text{cbbb9d5dc1059ed8} + 5\text{ec3802b9ecfef33} = 2\text{a7f1d895fd58e0b} \\
 H_1^{(1)} &= 6\text{29a292a367cd507} + 8\text{8146da76ff6f23a} = \text{eaae96d1a673c741} \\
 H_2^{(1)} &= 9\text{159015a3070dd17} + 7\text{001201948fb3d71} = 0\text{15a2173796c1a88} \\
 H_3^{(1)} &= 1\text{52fec8d8f70e5939} + \text{e1053fc85f9e56be} = \text{f6352ca156acaff7} \\
 H_4^{(1)} &= 6\text{7332667ffc00b31} + 5\text{f2eead69efb4233} = \text{c662113e9ebb4d64} \\
 H_5^{(1)} &= 8\text{eb44a8768581511} + 8\text{901cffe7a74db98} = 1\text{7b61a85e2ccf0a9} \\
 H_6^{(1)} &= \text{db0c2e0d64f98fa7} + 5\text{cdf6c58fc052572} = 3\text{7eb9a6660feb519} \\
 H_7^{(1)} &= 4\text{7b5481dbefa4fa4} + 4\text{779767cc2ec5321} = 8\text{f2ebe9a81e6a2c5}.
 \end{aligned}$$

The words of the *second* padded message block,  $M^{(2)}$ , are then assigned to the words  $W_0, \dots, W_{15}$  of the message schedule:

$W_0$	= 0000000000000000	$W_8$	= 0000000000000000
$W_1$	= 0000000000000000	$W_9$	= 0000000000000000
$W_2$	= 0000000000000000	$W_{10}$	= 0000000000000000
$W_3$	= 0000000000000000	$W_{11}$	= 0000000000000000
$W_4$	= 0000000000000000	$W_{12}$	= 0000000000000000
$W_5$	= 0000000000000000	$W_{13}$	= 0000000000000000
$W_6$	= 0000000000000000	$W_{14}$	= 0000000000000000
$W_7$	= 0000000000000000	$W_{15}$	= 0000000000000380.

The following schedule shows the hex values for  $a$ ,  $b$ ,  $c$ ,  $d$ ,  $e$ ,  $f$ ,  $g$ , and  $h$  after pass  $t$  of the “for  $t = 0$  to 79” loop described in Sec. 6.3.2, step 4.

	$a$ / $e$	$b$ / $f$	$c$ / $g$	$d$ / $h$
$t = 0 :$	657a3c2ca9639d40 791f2ad0055fdd62	2a7f1d895fd58e0b c662113e9ebb4d64	eaae96d1a673c741 17b61a85e2ccf0a9	015a2173796c1a88 37eb9a6660feb519
$t = 1 :$	2a4ad5d9b9fd6d86 dbf2e656b5be3f14	657a3c2ca9639d40 791f2ad0055fdd62	2a7f1d895fd58e0b c662113e9ebb4d64	eaae96d1a673c741 17b61a85e2ccf0a9
$t = 2 :$	f0aa6758653d1664 6e0466c82f4fd35d	2a4ad5d9b9fd6d86 dbf2e656b5be3f14	657a3c2ca9639d40 791f2ad0055fdd62	2a7f1d895fd58e0b c662113e9ebb4d64
$t = 3 :$	43a76f011a73d317 1367bd36d15e8b40	f0aa6758653d1664 6e0466c82f4fd35d	2a4ad5d9b9fd6d86 dbf2e656b5be3f14	657a3c2ca9639d40 791f2ad0055fdd62
$t = 4 :$	d802c2dfd7cc48f6 f73d759b839a2a21	43a76f011a73d317 1367bd36d15e8b40	f0aa6758653d1664 6e0466c82f4fd35d	2a4ad5d9b9fd6d86 dbf2e656b5be3f14
$t = 5 :$	481208e5e8314602 6b2271a46f14c843	d802c2dfd7cc48f6 f73d759b839a2a21	43a76f011a73d317 1367bd36d15e8b40	f0aa6758653d1664 6e0466c82f4fd35d
$t = 6 :$	af9f8112df35cf33 257f4a7d524d7b0b	481208e5e8314602 6b2271a46f14c843	d802c2dfd7cc48f6 f73d759b839a2a21	43a76f011a73d317 1367bd36d15e8b40
$t = 7 :$	6730781342d1131b 81957ad408cec995	af9f8112df35cf33 257f4a7d524d7b0b	481208e5e8314602 6b2271a46f14c843	d802c2dfd7cc48f6 f73d759b839a2a21
$t = 8 :$	82e64c677356a82e 10b62fdce4ebaa51	6730781342d1131b 81957ad408cec995	af9f8112df35cf33 257f4a7d524d7b0b	481208e5e8314602 6b2271a46f14c843
$t = 9 :$	203578820a8f27d0 9937b3a0cb9248a1	82e64c677356a82e 10b62fdce4ebaa51	6730781342d1131b 81957ad408cec995	af9f8112df35cf33 257f4a7d524d7b0b
$t = 10 :$	0bac2a84c29a1e2b 6ad288dab3de0d53	203578820a8f27d0 9937b3a0cb9248a1	82e64c677356a82e 10b62fdce4ebaa51	6730781342d1131b 81957ad408cec995
$t = 11 :$	dd3ff8a140485c25 3149b728123c465e	0bac2a84c29a1e2b 6ad288dab3de0d53	203578820a8f27d0 9937b3a0cb9248a1	82e64c677356a82e 10b62fdce4ebaa51
$t = 12 :$	e826239f830c5346 4bb7b199c4ced186	dd3ff8a140485c25 3149b728123c465e	0bac2a84c29a1e2b 6ad288dab3de0d53	203578820a8f27d0 9937b3a0cb9248a1
$t = 13 :$	32215ce49aae40f8 9a2872c72d790d49	e826239f830c5346 4bb7b199c4ced186	dd3ff8a140485c25 3149b728123c465e	0bac2a84c29a1e2b 6ad288dab3de0d53
$t = 14 :$	859533bac457f94e 539f225d25ebeb4c	32215ce49aae40f8 9a2872c72d790d49	e826239f830c5346 4bb7b199c4ced186	dd3ff8a140485c25 3149b728123c465e
$t = 15 :$	a88704d9962849f3	859533bac457f94e	32215ce49aae40f8	e826239f830c5346

	63bf0472ef24f7a5	539f225d25eb4c	9a2872c72d790d49	4bb7b199c4ced186
$t = 16 :$	3aa5c566a6cfad1c ce23f6380ead33c2	a88704d9962849f3 63bf0472ef24f7a5	859533bac457f94e 539f225d25eb4c	32215ce49aae40f8 9a2872c72d790d49
$t = 17 :$	2e9c483a7c08c9c1 b033f945f3e6b4a2	3aa5c566a6cfad1c ce23f6380ead33c2	a88704d9962849f3 63bf0472ef24f7a5	859533bac457f94e 539f225d25eb4c
$t = 18 :$	5a68585ae0835231 8a0187a9ce93d875	2e9c483a7c08c9c1 b033f945f3e6b4a2	3aa5c566a6cfad1c ce23f6380ead33c2	a88704d9962849f3 63bf0472ef24f7a5
$t = 19 :$	cf9cd481e6407ced 37a29fa30531bac7	5a68585ae0835231 8a0187a9ce93d875	2e9c483a7c08c9c1 b033f945f3e6b4a2	3aa5c566a6cfad1c ce23f6380ead33c2
$t = 20 :$	3f463f864f6474d9 0cf45bb3c07e847d	cf9cd481e6407ced 37a29fa30531bac7	5a68585ae0835231 8a0187a9ce93d875	2e9c483a7c08c9c1 b033f945f3e6b4a2
$t = 21 :$	cea26288dff931a5 34f1b5f46bf48a73	3f463f864f6474d9 0cf45bb3c07e847d	cf9cd481e6407ced 37a29fa30531bac7	5a68585ae0835231 8a0187a9ce93d875
$t = 22 :$	89634cd0f4f6c08a 3a728a543405a8e4	cea26288dff931a5 34f1b5f46bf48a73	3f463f864f6474d9 0cf45bb3c07e847d	cf9cd481e6407ced 37a29fa30531bac7
$t = 23 :$	625fa38464e5c880 cee1b47a49b2fc42	89634cd0f4f6c08a 3a728a543405a8e4	cea26288dff931a5 34f1b5f46bf48a73	3f463f864f6474d9 0cf45bb3c07e847d
$t = 24 :$	7dd21453a15a3b92 9308bfa1be1f800b	625fa38464e5c880 cee1b47a49b2fc42	89634cd0f4f6c08a 3a728a543405a8e4	cea26288dff931a5 34f1b5f46bf48a73
$t = 25 :$	3d76277bc8cb0601 480e017f5d1f0b1e	7dd21453a15a3b92 9308bfa1be1f800b	625fa38464e5c880 cee1b47a49b2fc42	89634cd0f4f6c08a 3a728a543405a8e4
$t = 26 :$	c8d904196f5a1f54 4bd2f1f6e940c332	3d76277bc8cb0601 480e017f5d1f0b1e	7dd21453a15a3b92 9308bfa1be1f800b	625fa38464e5c880 cee1b47a49b2fc42
$t = 27 :$	b033139b58b6e423 f816ec1cbe0adafb	c8d904196f5a1f54 4bd2f1f6e940c332	3d76277bc8cb0601 480e017f5d1f0b1e	7dd21453a15a3b92 9308bfa1be1f800b
$t = 28 :$	097768182cb65f57 62e3de54dcd8f974	b033139b58b6e423 f816ec1cbe0adafb	c8d904196f5a1f54 4bd2f1f6e940c332	3d76277bc8cb0601 480e017f5d1f0b1e
$t = 29 :$	3196649ab5f5cc39 f6887de116d0bd8f	097768182cb65f57 62e3de54dcd8f974	b033139b58b6e423 f816ec1cbe0adafb	c8d904196f5a1f54 4bd2f1f6e940c332
$t = 30 :$	f78d3d221d16965f c7e4859c2858ed3c	3196649ab5f5cc39 f6887de116d0bd8f	097768182cb65f57 62e3de54dcd8f974	b033139b58b6e423 f816ec1cbe0adafb
$t = 31 :$	f58e9876b4984b51 621352b394b8ca02	f78d3d221d16965f c7e4859c2858ed3c	3196649ab5f5cc39 f6887de116d0bd8f	097768182cb65f57 62e3de54dcd8f974
$t = 32 :$	38fbf0e726e04f78 4319856f17a0a430	f58e9876b4984b51 621352b394b8ca02	f78d3d221d16965f c7e4859c2858ed3c	3196649ab5f5cc39 f6887de116d0bd8f
$t = 33 :$	f4be0b32a57597a2 c6d392a3b4eb0ed8	38fbf0e726e04f78 4319856f17a0a430	f58e9876b4984b51 621352b394b8ca02	f78d3d221d16965f c7e4859c2858ed3c
$t = 34 :$	f8a6b3fe2e4f0634 602663c0f34eff33	f4be0b32a57597a2 c6d392a3b4eb0ed8	38fbf0e726e04f78 4319856f17a0a430	f58e9876b4984b51 621352b394b8ca02
$t = 35 :$	9bc3871be8046113 05542ecd9883c6ba	f8a6b3fe2e4f0634 602663c0f34eff33	f4be0b32a57597a2 c6d392a3b4eb0ed8	38fbf0e726e04f78 4319856f17a0a430
$t = 36 :$	f1bd2d46be619585 e47b9933bafdc655	9bc3871be8046113 05542ecd9883c6ba	f8a6b3fe2e4f0634 602663c0f34eff33	f4be0b32a57597a2 c6d392a3b4eb0ed8
$t = 37 :$	24c84b58d119affe 5ae0b1175beb5d2b	f1bd2d46be619585 e47b9933bafdc655	9bc3871be8046113 05542ecd9883c6ba	f8a6b3fe2e4f0634 602663c0f34eff33
$t = 38 :$	ec6d3abc2b291fd3	24c84b58d119affe	f1bd2d46be619585	9bc3871be8046113



	9ecc381d277748a3	5ae0b1175beb5d2b	e47b9933bafdc655	05542ecd9883c6ba
$t = 39$ :	e266c1f77d5ee90e d92f34c110296b32	ec6d3abc2b291fd3 9ecc381d277748a3	24c84b58d119affe 5ae0b1175beb5d2b	f1bd2d46be619585 e47b9933bafdc655
$t = 40$ :	5adbaa463642b570 83e8f410f859388e	e266c1f77d5ee90e d92f34c110296b32	ec6d3abc2b291fd3 9ecc381d277748a3	24c84b58d119affe 5ae0b1175beb5d2b
$t = 41$ :	50fdb7bb2e499a34 257ed8ea645e933a	5adbaa463642b570 83e8f410f859388e	e266c1f77d5ee90e d92f34c110296b32	ec6d3abc2b291fd3 9ecc381d277748a3
$t = 42$ :	06514212bb7fa152 466781db35181abe	50fdb7bb2e499a34 257ed8ea645e933a	5adbaa463642b570 83e8f410f859388e	e266c1f77d5ee90e d92f34c110296b32
$t = 43$ :	673ed5a55ff2b07d ba78f3545e7914f0	06514212bb7fa152 466781db35181abe	50fdb7bb2e499a34 257ed8ea645e933a	5adbaa463642b570 83e8f410f859388e
$t = 44$ :	125e2e5118393e2b 4453b23a3e13b090	673ed5a55ff2b07d ba78f3545e7914f0	06514212bb7fa152 466781db35181abe	50fdb7bb2e499a34 257ed8ea645e933a
$t = 45$ :	07ee813df5910cec eae013a0510d23cc	125e2e5118393e2b 4453b23a3e13b090	673ed5a55ff2b07d ba78f3545e7914f0	06514212bb7fa152 466781db35181abe
$t = 46$ :	0a0508f0a1d719c3 a93815eb58891016	07ee813df5910cec eae013a0510d23cc	125e2e5118393e2b 4453b23a3e13b090	673ed5a55ff2b07d ba78f3545e7914f0
$t = 47$ :	0fc8f3b3efcb1b96 a071cc73b966e801	0a0508f0a1d719c3 a93815eb58891016	07ee813df5910cec eae013a0510d23cc	125e2e5118393e2b 4453b23a3e13b090
$t = 48$ :	02aa5b28199f304a a49f1e14f8a2be7a	0fc8f3b3efcb1b96 a071cc73b966e801	0a0508f0a1d719c3 a93815eb58891016	07ee813df5910cec eae013a0510d23cc
$t = 49$ :	9223e1b34382f104 bfe2106e512a7331	02aa5b28199f304a a49f1e14f8a2be7a	0fc8f3b3efcb1b96 a071cc73b966e801	0a0508f0a1d719c3 a93815eb58891016
$t = 50$ :	e01a1e47ee8d5656 592b899b35469a78	9223e1b34382f104 bfe2106e512a7331	02aa5b28199f304a a49f1e14f8a2be7a	0fc8f3b3efcb1b96 a071cc73b966e801
$t = 51$ :	fa7b17aad857c2f4 eb6e85e4682c1671	e01a1e47ee8d5656 592b899b35469a78	9223e1b34382f104 bfe2106e512a7331	02aa5b28199f304a a49f1e14f8a2be7a
$t = 52$ :	0c523b7a3c84ab77 b5e80e871ac0c005	fa7b17aad857c2f4 eb6e85e4682c1671	e01a1e47ee8d5656 592b899b35469a78	9223e1b34382f104 bfe2106e512a7331
$t = 53$ :	c773d8b69da1fde2 be2b0602fc6f8f65	0c523b7a3c84ab77 b5e80e871ac0c005	fa7b17aad857c2f4 eb6e85e4682c1671	e01a1e47ee8d5656 592b899b35469a78
$t = 54$ :	c6b1bc79a4f23679 c80bdc57f38a05e4	c773d8b69da1fde2 be2b0602fc6f8f65	0c523b7a3c84ab77 b5e80e871ac0c005	fa7b17aad857c2f4 eb6e85e4682c1671
$t = 55$ :	bef9bb0fe467fd60 1dab0bd116e434e5	c6b1bc79a4f23679 c80bdc57f38a05e4	c773d8b69da1fde2 be2b0602fc6f8f65	0c523b7a3c84ab77 b5e80e871ac0c005
$t = 56$ :	8e3db3e380ec7f22 32ef50751734ffee	bef9bb0fe467fd60 1dab0bd116e434e5	c6b1bc79a4f23679 c80bdc57f38a05e4	c773d8b69da1fde2 be2b0602fc6f8f65
$t = 57$ :	1003ec42412c7b7d 1ec0d46f349fd058	8e3db3e380ec7f22 32ef50751734ffee	bef9bb0fe467fd60 1dab0bd116e434e5	c6b1bc79a4f23679 c80bdc57f38a05e4
$t = 58$ :	375facc76291f85e 59c8bc0488f9768b	1003ec42412c7b7d 1ec0d46f349fd058	8e3db3e380ec7f22 32ef50751734ffee	bef9bb0fe467fd60 1dab0bd116e434e5
$t = 59$ :	bd113d92e0354fb9 e66c73db3fad397d	375facc76291f85e 59c8bc0488f9768b	1003ec42412c7b7d 1ec0d46f349fd058	8e3db3e380ec7f22 32ef50751734ffee
$t = 60$ :	2f61d4fd8e36d9d4 e9f21933e1c02948	bd113d92e0354fb9 e66c73db3fad397d	375facc76291f85e 59c8bc0488f9768b	1003ec42412c7b7d 1ec0d46f349fd058
$t = 61$ :	1blad88b92701ae2 6fd0c1719bcac335	2f61d4fd8e36d9d4 e9f21933e1c02948	bd113d92e0354fb9 e66c73db3fad397d	375facc76291f85e 59c8bc0488f9768b



$t = 62$ :	93d09fc06a19c5da b765273f571a571e	1blad88b92701ae2 6fd0c1719bcac335	2f61d4fd8e36d9d4 e9f21933e1c02948	bd113d92e0354fb9 e66c73db3fad397d
$t = 63$ :	04bea2ce99cc3bf6 6ab0e443c2f63714	93d09fc06a19c5da b765273f571a571e	1blad88b92701ae2 6fd0c1719bcac335	2f61d4fd8e36d9d4 e9f21933e1c02948
$t = 64$ :	02ebfc0a13492f52 77300c52e05af415	04bea2ce99cc3bf6 6ab0e443c2f63714	93d09fc06a19c5da b765273f571a571e	1blad88b92701ae2 6fd0c1719bcac335
$t = 65$ :	1bf525abce8d6f04 8faf12c33bb371b9	02ebfc0a13492f52 77300c52e05af415	04bea2ce99cc3bf6 6ab0e443c2f63714	93d09fc06a19c5da b765273f571a571e
$t = 66$ :	b6a36a3431547328 fa8bb40b4e08100f	1bf525abce8d6f04 8faf12c33bb371b9	02ebfc0a13492f52 77300c52e05af415	04bea2ce99cc3bf6 6ab0e443c2f63714
$t = 67$ :	ffdaf83202af0d72 8045a82f723a9b4e	b6a36a3431547328 fa8bb40b4e08100f	1bf525abce8d6f04 8faf12c33bb371b9	02ebfc0a13492f52 77300c52e05af415
$t = 68$ :	12737373d2985232 870dbce23bad8988	ffdaf83202af0d72 8045a82f723a9b4e	b6a36a3431547328 fa8bb40b4e08100f	1bf525abce8d6f04 8faf12c33bb371b9
$t = 69$ :	6189f68162b256b5 8c059af157146580	12737373d2985232 870dbce23bad8988	ffdaf83202af0d72 8045a82f723a9b4e	b6a36a3431547328 fa8bb40b4e08100f
$t = 70$ :	20b0a9a1d21c482d f22b874c96785ec8	6189f68162b256b5 8c059af157146580	12737373d2985232 870dbce23bad8988	ffdaf83202af0d72 8045a82f723a9b4e
$t = 71$ :	ef6d863c2127b394 b7aee28337d69dab	20b0a9a1d21c482d f22b874c96785ec8	6189f68162b256b5 8c059af157146580	12737373d2985232 870dbce23bad8988
$t = 72$ :	d3efe8b442689074 22491ab9cdec6b0	ef6d863c2127b394 b7aee28337d69dab	20b0a9a1d21c482d f22b874c96785ec8	6189f68162b256b5 8c059af157146580
$t = 73$ :	4694354944a9f487 659890a5818d0c50	d3efe8b442689074 22491ab9cdec6b0	ef6d863c2127b394 b7aee28337d69dab	20b0a9a1d21c482d f22b874c96785ec8
$t = 74$ :	b93c2403773dd08c 88c2c2ac52c4f679	4694354944a9f487 659890a5818d0c50	d3efe8b442689074 22491ab9cdec6b0	ef6d863c2127b394 b7aee28337d69dab
$t = 75$ :	025848e3ab6b69d3 750da3d4e16a1b64	b93c2403773dd08c 88c2c2ac52c4f679	4694354944a9f487 659890a5818d0c50	d3efe8b442689074 22491ab9cdec6b0
$t = 76$ :	396b53e58d04471b 700486bf252cba75	025848e3ab6b69d3 750da3d4e16a1b64	b93c2403773dd08c 88c2c2ac52c4f679	4694354944a9f487 659890a5818d0c50
$t = 77$ :	51b6f9a3c1ceeb4a e6b3850de8ae6230	396b53e58d04471b 700486bf252cba75	025848e3ab6b69d3 750da3d4e16a1b64	b93c2403773dd08c 88c2c2ac52c4f679
$t = 78$ :	526a98f5dc595406 4f0dcf74aea76f90	51b6f9a3c1ceeb4a e6b3850de8ae6230	396b53e58d04471b 700486bf252cba75	025848e3ab6b69d3 750da3d4e16a1b64
$t = 79$ :	deb3eaaa973bb9dd 3665b5dbb6c2e055	526a98f5dc595406 4f0dcf74aea76f90	51b6f9a3c1ceeb4a e6b3850de8ae6230	396b53e58d04471b 700486bf252cba75

That completes the processing of the second and final message block,  $M^{(2)}$ . The final hash value,  $H^{(2)}$ , is calculated to be

$$\begin{aligned}
 H_0^{(2)} &= 2a7f1d895fd58e0b + deb3eaaa973bb9dd = 09330c33f71147e8 \\
 H_1^{(2)} &= eaae96d1a673c741 + 526a98f5dc595406 = 3d192fc782cd1b47 \\
 H_2^{(2)} &= 015a2173796c1a88 + 51b6f9a3c1ceeb4a = 53111b173b3b05d2 \\
 H_3^{(2)} &= f6352ca156acaff7 + 396b53e58d04471b = 2fa08086e3b0f712 \\
 H_4^{(2)} &= c662113e9ebb4d64 + 3665b5dbb6c2e055 = fcc7c71a557e2db9
 \end{aligned}$$

$$\begin{aligned}
 H_5^{(2)} &= 17b61a85e2ccf0a9 + 4f0dcf74aea76f90 = 66c3e9fa91746039 \\
 H_6^{(2)} &= 37eb9a6660feb519 + e6b3850de8ae6230 = 1e9f1f7449ad1749 \\
 H_7^{(2)} &= 8f2ebe9a81e6a2c5 + 700486bf252cba75 = ff334559a7135d3a .
 \end{aligned}$$

The final hash value is truncated to its left-most 384 bits (i.e.,  $H_0^{(1)}, \dots, H_5^{(1)}$ ), resulting in the 384-bit message digest

```
09330c33f71147e8 3d192fc782cd1b47 53111b173b3b05d2 2fa08086e3b0f712
fcc7c71a557e2db9 66c3e9fa91746039 .
```

### D.3 SHA-384 Example (Long Message)

Let the message  $M$  be the binary-coded form of the ASCII string which consists of 1,000,000 repetitions of the character ‘a’. The resulting SHA-384 message digest is

```
9d0e1809716474cb 086e834e310a4a1c ed149e9c00f24852 7972cec5704c2a5b
07b8b3dc38ecc4eb ae97ddd87f3d8985 .
```

## APPENDIX E: REFERENCES

- [180-1] Federal Information Processing Standards (FIPS) Publication 180-1, *Secure Hash Standard (SHS)*, U.S. DoC/NIST, April 17, 1995.
- [HAC] A. Menezes, P. van Oorschot, and S. Vanstone. *Handbook of Applied Cryptography*, CRC Press, Inc., October 1997.